### A CUDA IMPLEMENTATION OF THE HPCG BENCHMARK

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# OUTLINE

High Performance Conjugate Gradient Benchmark

- Motivation
- Overview
- Optimization
- Performance Results
  - Single GPU
  - GPU Supercomputers
- Conclusion

# WHY HPCG ?

### HPL (Linpack) Top500 benchmark

Supercomputer Ranking / Evaluation

- Dense Linear Algebra (Ax = b)
  - **Compute intensive** 
    - DGEMM (Matrix-Matrix Multiply)
    - O(N3)FLOPS / O(N2) Data
    - 10-100 Flop/Byte



Workload does not correlate with many modern applications

# WHY HPCG?

New Benchmark to Supplement HPL

- Common Computation Patterns not addressed by HPL
- Numerical Solution of PDEs
- Memory Intensive
- Network



# HPCG BENCHMARK

Preconditioned Conjugate Gradient Algorithm
 Sparse Linear Algebra (Ax = b), Iterative solver
 Bandwidth Intensive: 1/6 Flop/Byte
 Simple Problem (sparsity pattern of Matrix A)

- Simplifies matrix generation/solution validation
- Regular 3D grid, 27-point stencil
- Nx x Ny x Nz local domain / Px x Py x Pz Processors
- Communications: boundary + global reduction



# HPCG ALGORITHM

#### Multi-Grid Preconditioner

Symmetric-Gauss-Seidel Smoother (SYMGS)



Sparse Matrix Vector Multiply (SPMV)
 Dot Product - MPI\_Allreduce()

Algorithm 1 Preconditioned Conjugate Gradient 1: k = 02: Compute the residual  $r_0 = b - Ax_0$ 3: while  $(||r_k|| < \epsilon)$  do  $z_k = M^{-1} r_k$ 4: 5:  $\overline{k} = k + 1$ 6: if k = 1 then 7:  $p_1 = z_0$ 8: else  $\beta_{k} = \frac{r_{k-1}^{T} z_{k-1}}{p_{k} = z_{k-1} + \beta_{k} p_{k-1}}$ 9: 10: 11: end if  $\alpha_k = r_{k-1}^T z_{k-1} / p_k^T A p_k$ 12:13: $x_k = x_{k-1} + \alpha_k p_k$ 14:  $r_k = r_{k-1} - \alpha_k A p_k$ 15: end while 16:  $x = x_k$ 

# HPCG BENCHMARK

- Problem Setup initialize data structures
- Optimization (required to expose parallelism in SYMGS smoother)
  - Matrix analysis / reordering / data layout
  - Time counted against final performance result
- Reference Run 50 iterations with reference code Record Residual
- Optimized Run converge to Reference Residual
  - Matrix re-ordering slows convergence (55-60 iterations)
  - Additional iterations counted against final performance result
  - Repeat to fill target execution time (few minutes typical, 1 hour for official run )

# HPCG

SPMV (y = Ax)

```
Exchange_Halo(x) //neighbor communications
for row = 0 to nrows
sum \leftarrow 0
for j = 0 to nonzeros_in_row[ row ]
col \leftarrow A_col[ j ]
val \leftarrow A_val[ j ]
sum \leftarrow sum + val * x[ col ]
y[ row ] \leftarrow sum
```

No dependencies between rows, safe to process rows in parallel

# HPCG

### SYMGS (Ax = y, smooth x)

```
Exchange_Halo(x) //neighbor communications
for row = 0 to nrows (Fwd Sweep, then Backward Sweep for row = nrows to 0)
sum ← b[ row ]
for j = 0 to nonzeros_in_row[ row ]
col ← A_col[ j ]
val ← A_val[ j ]
if( col != row ) sum ← sum - val * x[ col ]
x[ row ] ← sum / A_diag[ row ]
```

if col < row, must wait for x[col] to be updated

# MATRIX REORDERING (COLORING)

### SYMGS - order requirement

- Previous rows must have new value
- reorder by color (independent rows)
- 2D example: 5-point stencil -> red-black
- 3D 27-point stencil = 8 colors



# MATRIX REORDERING (COLORING)

- Coloring to extract parallelism
- Assignment of "color" (integer) to vertices (rows), with no two adjacent vertices the same color

"Efficient Graph Matching and Coloring on the GPU" - (Jon Cohen)

- Luby / Jones-Plassman based algorithm
- Compare hash of row index with neighbors
- Assign color if local extrema
- Optional: recolor to reduce # of colors



# MORE OPTIMIZATIONS

- > Overlap Computation with neighbor communication
- > Overlap 1/3 MPI\_Allreduce with Computation
  - \_\_\_LDG loads for irregular access patterns (SPMV + SYMGS)

# OPTIMIZATIONS

SPMV Overlap Computation with communications

Gather to GPU send\_buffer
 Copy send\_buffer to CPU
 MPI\_send / MPI\_recv
 Copy recv\_buffer to GPU
 Launch SPMV Kernel

Time

GPU

CPU

# OPTIMIZATIONS

SPMV Overlap Computation with communications

 Gather to GPU send\_buffer Copy send\_buffer to CPU Launch SPMV interior Kernel MPI\_send / MPI\_recv Copy recv\_buffer to GPU Launch SPMV boundary Kernel

GPU Stream A GPU Stream B CPU

Time





#### Single GPU HPCG GFLOPS

SPMV MG TOTAL



#### **HPCG vs STREAM Memory Bandwidth**



■ K20X ■ K40 ■ IVB

#### HPCG GF vs STREAM BW



# **RESULTS - GPU SUPERCOMPUTERS**

### Titan @ ORNL

- Cray XK7, 18688 Nodes
- 16-core AMD Interlagos + K20X
- Gemini Network 3D Torus Topology
- Piz Daint @ CSCS
  - Cray XC30, 5272 Nodes
  - 8-core Xeon E5 + K20X
  - Aries Network Dragonfly Topology





# **RESULTS - GPU SUPERCOMPUTERS**

- I GPU = 20.8 GFLOPS (ECC ON)
- ~7% iteration overhead at scale
- Titan @ ORNL
  - 322 TFLOPS (18648 K20X)
  - 89% efficiency (17.3 GF per GPU)
  - Piz Daint @ CSCS
    - > 97 TFLOPS (5265 K20X)
    - 97% efficiency (19.0 GF per GPU)



# **RESULTS - GPU SUPERCOMPUTERS**

DDOT (-10%)

- MPI\_Allreduce()
- Scales as Log(#nodes)

> MG (-2%)

Exchange Halo (neighbor)
SPMV (-0%)

Overlapped w/Compute



# SUPERCOMPUTER COMPARISON

HPCG	System	HPCG	Itera-	# Procs	Processor	HPCG	Bandwidth	Efficiency
Rank		GFLOPS	tions		Type	Per Proc	Per Proc	FLOP/BYTE
1	Tianhe-2	$580,\!109$	57	$46,\!080$	Xeon-Phi-31S1P	$12.59 \ \mathrm{GF}$	320  GB/s	0.039
2	К	$426,\!972$	51	$82,\!944$	Sparc64-viiifx	$5.15~\mathrm{GF}$	64  GB/s	0.080
3	Titan	$322,\!321$	55	$18,\!648$	Tesla-K20X+ECC	$17.28 \ \mathrm{GF}$	$250 \mathrm{~GB/s}$	0.069
5	Piz-Daint	$98,\!979$	55	$5,\!208$	Tesla-K20X+ECC	$19.01 \ \mathrm{GF}$	$250 \mathrm{~GB/s}$	0.076
8	HPC2	$49,\!145$	54	$2,\!610$	Tesla-K20X+ECC	$18.83 \ \mathrm{GF}$	$250 \mathrm{~GB/s}$	0.075
	HPC2	$60,\!642$	54	2,600	Tesla-K20X	23.32  GF	$250~\mathrm{GB/s}$	0.093

# CONCLUSIONS

GPUs proven effective for HPL, especially for power efficiency

- High flop rate
- GPUs also very effective for HPCG
  - High memory bandwidth
  - Stacked memory will give a huge boost
- Future work will add CPU + GPU

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