

FASP 0

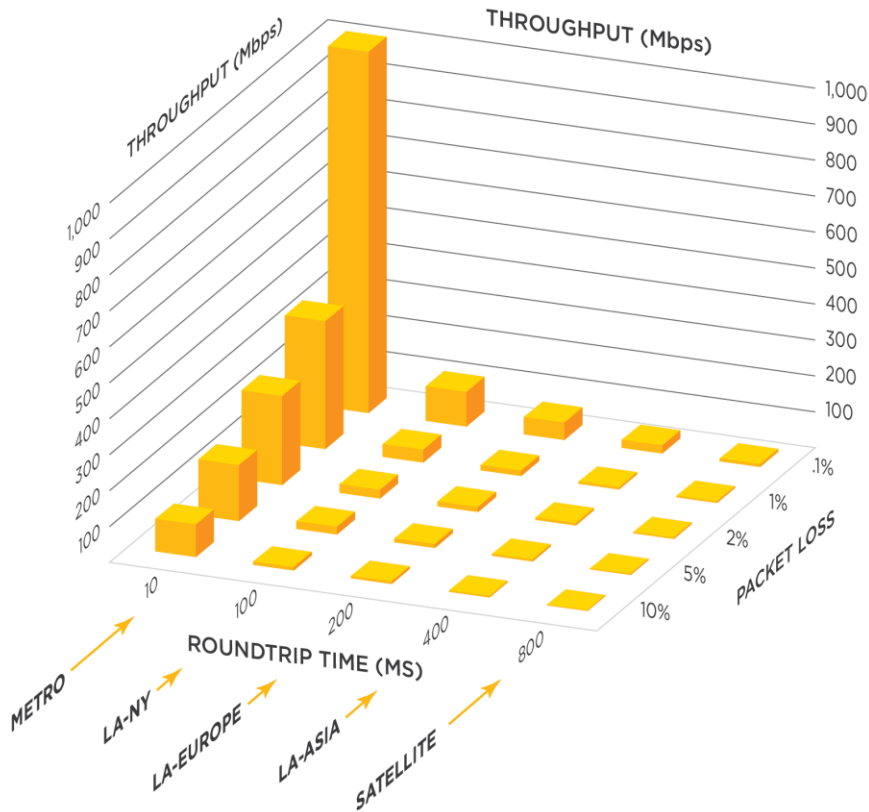
Aspera's FASPeSt FASP

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Creating next-generation transport technologies
that move the world's digital assets at maximum speed,
regardless of file size, transfer distance and network conditions

- Founded in 2004 with a focus on solving data transmission across the WAN
- Real networks, Any platform, Any data
- With development branch of FASP 0 performance shown at 100gbit/s using a Single node w/ encryption, regular packets, DPDK acceleration.
- Performance expected to improve as PCIe interconnects improve



Distance degrades conditions on all networks

- Latency (or Round Trip Times) increases
- Packet loss increases
- Fast networks are just as prone to degradation

TCP performance degrades severely with distance

- TCP was designed for LANs and does not perform well over distance
- Throughput bottlenecks are severe as latency & packet loss increase

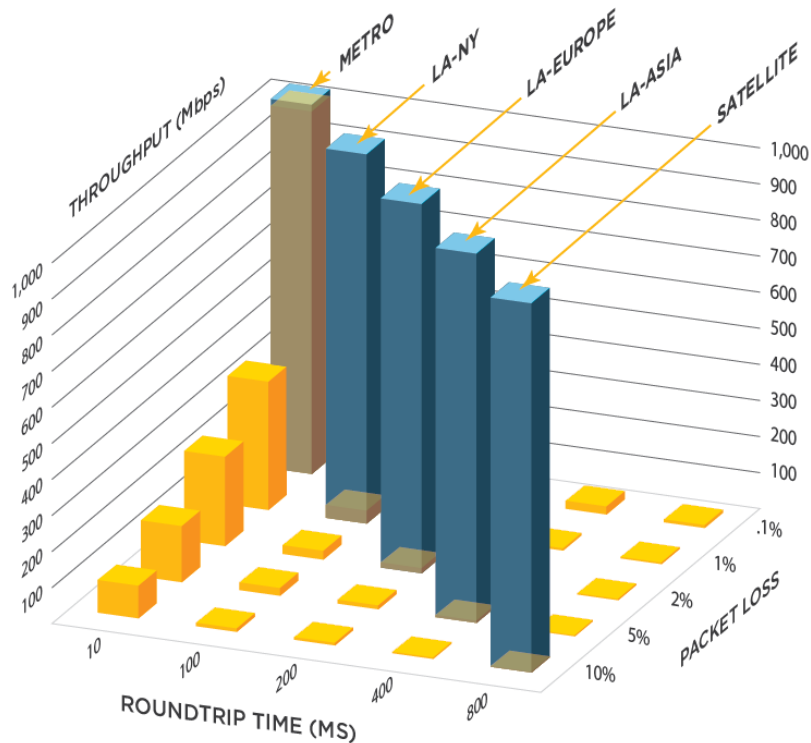
TCP does not scale with bandwidth

- TCP designed for low bandwidth
- Adding more bandwidth does not improve throughput

Alternative technologies

- TCP-based - Network latency & packet loss must be low to work well
- UDP blasters – Inefficient use of bandwidth leads to congestion
- Modified TCP – Does not scale well on high-speed networks
- Data caching - Inappropriate for many large file transfer workflows
- Data compression - Time consuming & impractical for some file types

Note: Table displays throughput degradation of TCP transfers on a 1Gbps network as estimated round trip time and packet loss increases with distance.



Maximum transfer speed

- Optimal end-to-end throughput efficiency
- Transfer performance scales with bandwidth independent of transfer distance and resilient to packet loss

Congestion avoidance and policy control

- Automatic, full utilization of available bandwidth (fair play)
- On-the-fly prioritization of transfers
- Set caps on bandwidth allocation for transfers

Uncompromising security and reliability

- Secure, SSH user/endpoint authentication
- AES-128 to 256 cryptography of every packet in transit
- Encryption at rest (EAR) requires second password
- FIPS 140-2 compliant, built on the open SSL libraries
- Automatic resume of partial or failed transfers

Scalable management, monitoring and control

- Support highly concurrent transfers
- Real-time progress, performance and bandwidth utilization
- Detailed transfer history, logging, and manifest

Note: The relative bandwidth utilization for FASP transfers over a 1 Gbps network are immune to latency (distance) with very little effect from packet loss.



MOVING A 10GB FILE		Across US	US - Europe	US - Asia
Legacy Transport	100 Mbps	10-20 Hours	15-20 Hours	Impractical
	1 Gbps			
	10 Gbps			
Aspera FASP®	100 Mbps	14 Min	14 Min	14 Min
	1 Gbps	1.4 Min	1.4 Min	1.4 Min
	10 Gbps	8.4 Sec	8.4 Sec	8.4 Sec

Location Agnostic

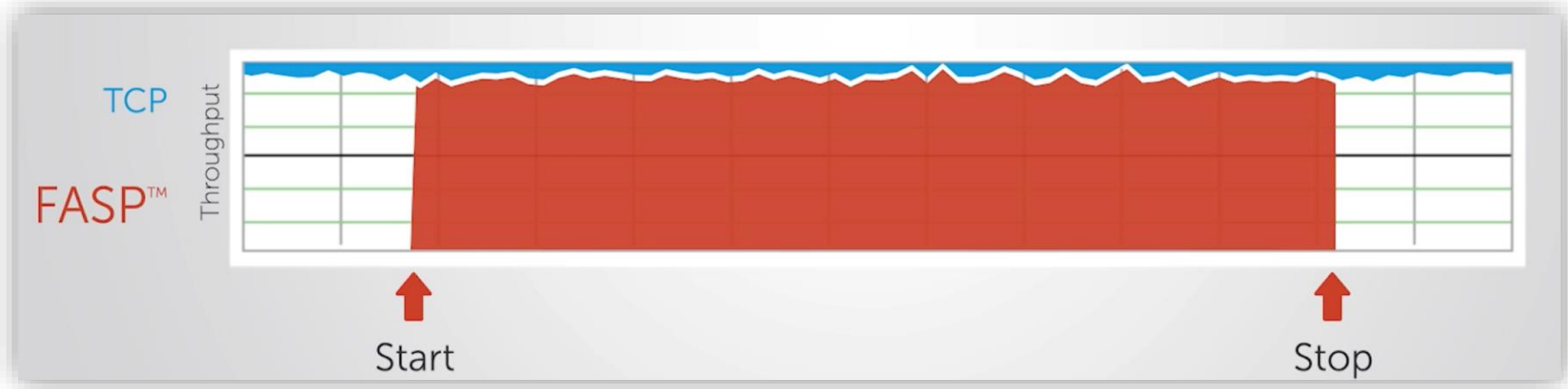
FASP transfer speeds remain virtually constant as transfer distances increase while FTP speeds dramatically decrease

Predictable & Reliable

FASP transfer times decrease linearly as bandwidth increases. However, FTP transfer times don't improve with bandwidth

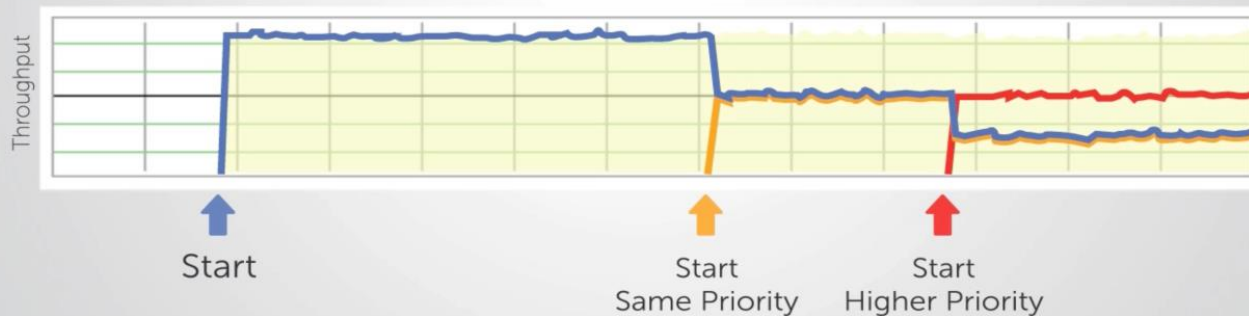
Versatile

Supports massive file sizes (500 GB+) as easily as very large sets (millions) of small files



Extraordinary bandwidth control that doesn't saturate the network

- Automatic detection & full utilization of available bandwidth with “fair” policy protection of other network traffic
- Allows “bursts” in TCP traffic and reclaims unused bandwidth as it becomes available



Real-time prioritization of transfers

- On-the-fly, per flow, user and job prioritization of transfers
- Concurrent transfers adjust bandwidth on the fly, allocating available bandwidth based on transfer priority

System-wide monitoring and reporting

- Real-time progress and performance analysis along with detailed transfer history, logging and manifest

Extraordinary bandwidth control

- Automatic, full utilization of available bandwidth with protection of other network traffic with “fair” policy
- Allows “bursts” in TCP traffic and reclaims unused bandwidth as it becomes available

Complete Built-in Security

- Secure endpoint authentication, data encryption on-the-fly and at rest, and per-packet integrity verification
- FIPS 140-2 compliant, built on the openssl libraries

Secure User/Endpoint Authentication

- Authentication via secure SSH mechanisms: interactive password or public key
- LDAP, Active Directory user authentication
- Native File System Access Control support across all operating systems

AES-128 Cryptography

- On-the-fly data encryption
- Data encryption in transit and (optionally) at rest (secured storage of transferred content), client and server options

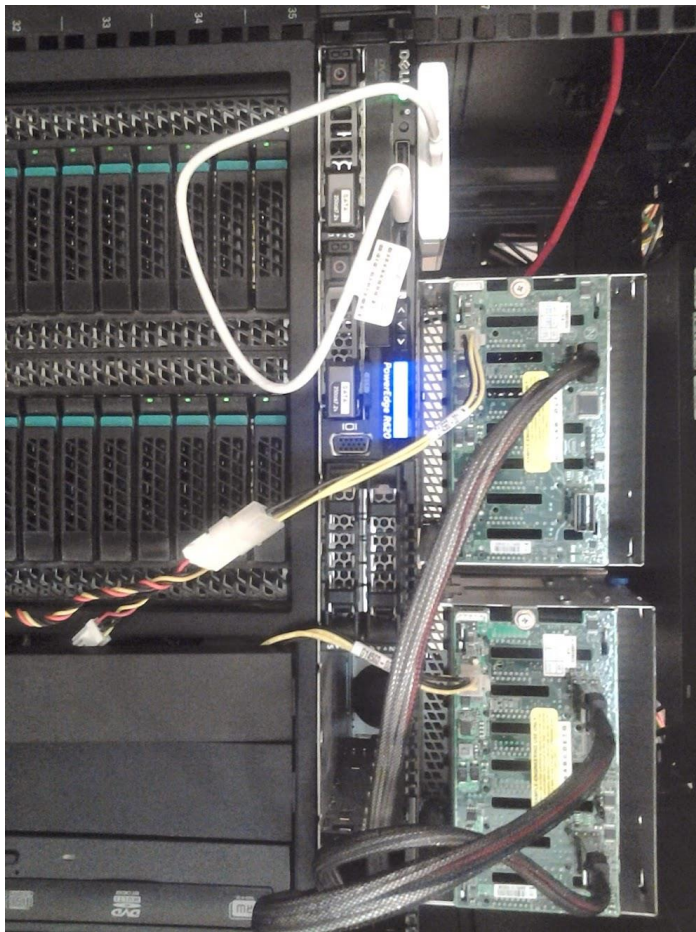
Data Integrity Verification

- Each transmitted data block is verified with a cryptographic hash function
- Protects against man-in-the-middle, re-play, and UDP denial-of-service attacks

100% Reliable Data Transmission

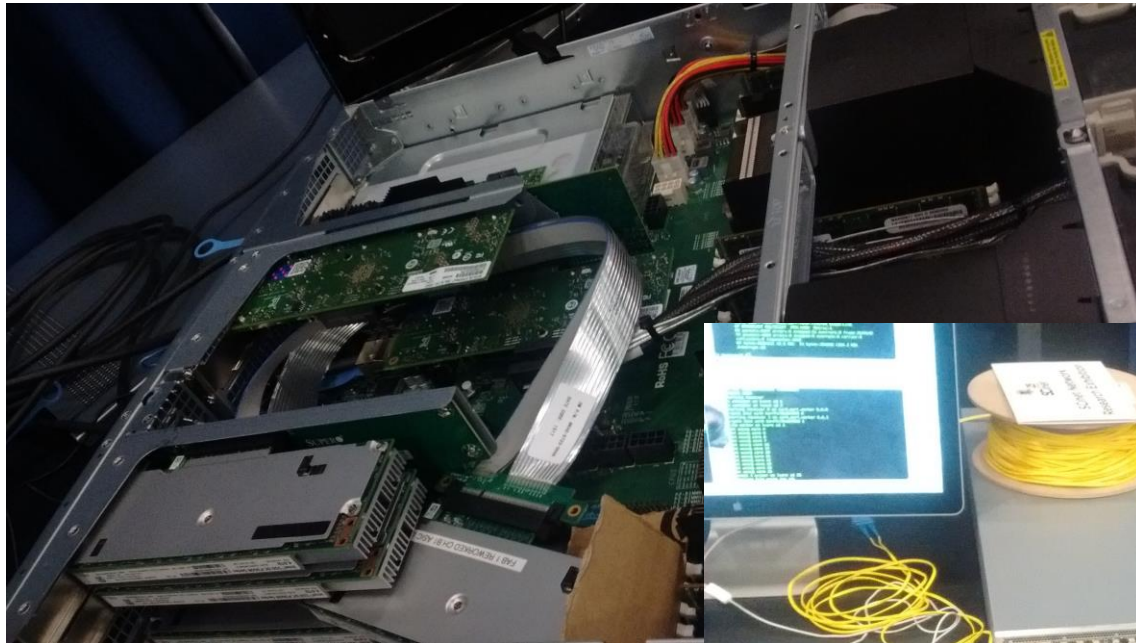
- Session semantics guarantee 100% bit-for-bit identical data copy at the destination
- Automatic resume of partial or failed transfers
- Automatic HTTP fallback in highly restrictive networks

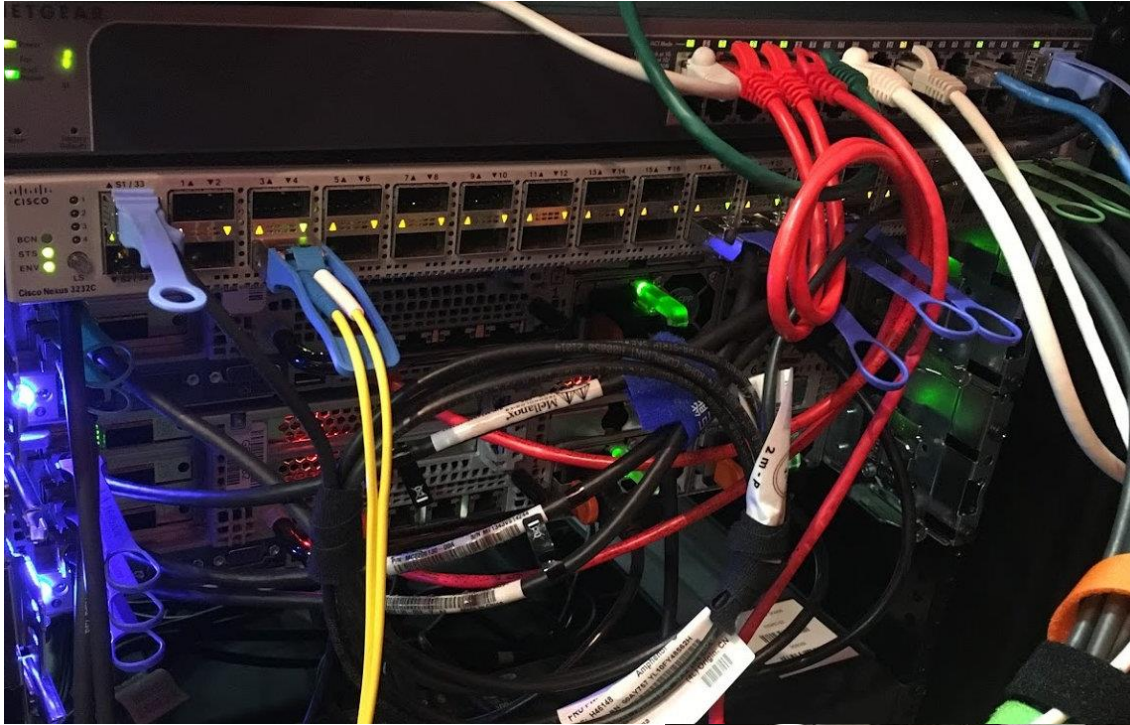
FASP 0



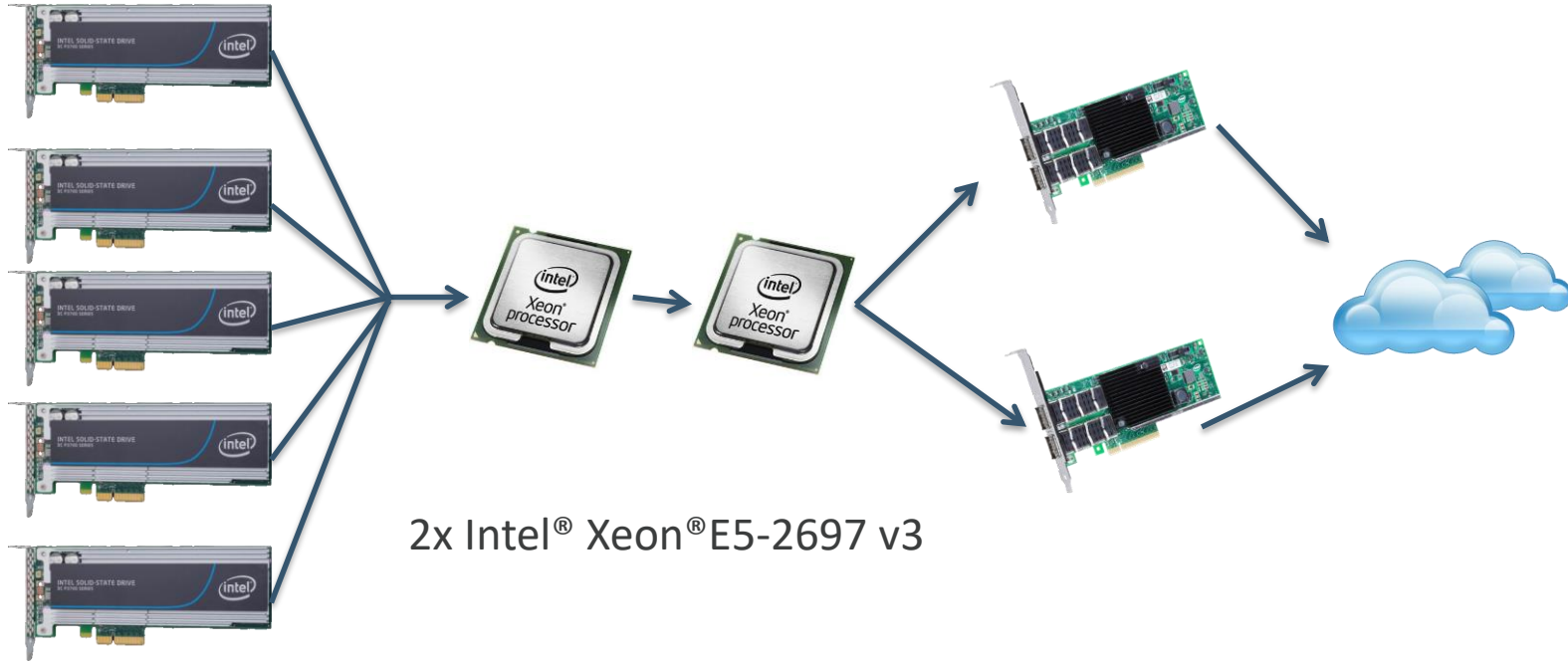
- Initial Xeon Ivy Bridge based DTN
- Used Hardware RAID and 4x10 gbit Intel NICs
- Single PCIe v2 x16 Hardware RAID card
- 12 Intel SATA SSDs
- Storage was bottleneck
- Hacked together using a combination of Intel Grizzly Pass systems and Dell Server. Note the internal SAS cables routing outside the Dell chassis along with power fed from other system.
- Whitepaper:
http://asperasoft.com/fileadmin/media/Asperasoft.com/Resources/White_Papers/Big_Data_Transfer_Phase_2_WP_FINAL.pdf









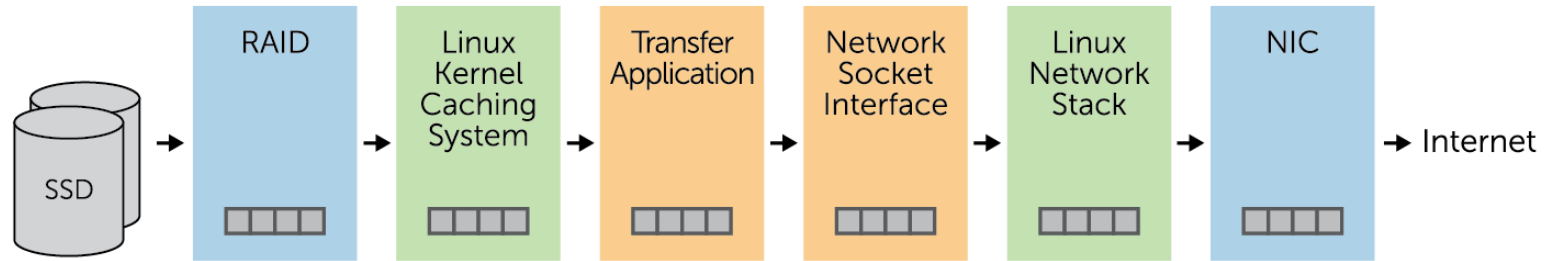


2x Intel® Xeon® E5-2697 v3

5x Intel® DC P3700 NVMe SSD

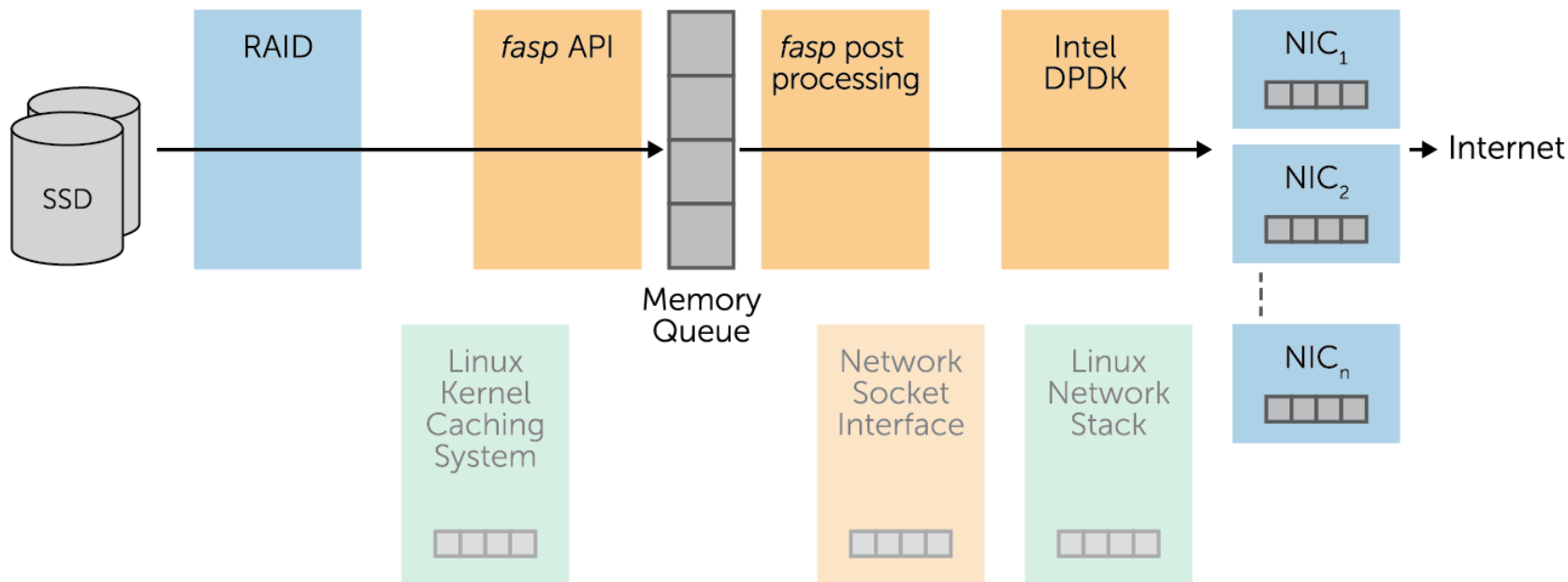
2x Intel® XL710 40 GbE Ethernet QSFP+

BEFORE



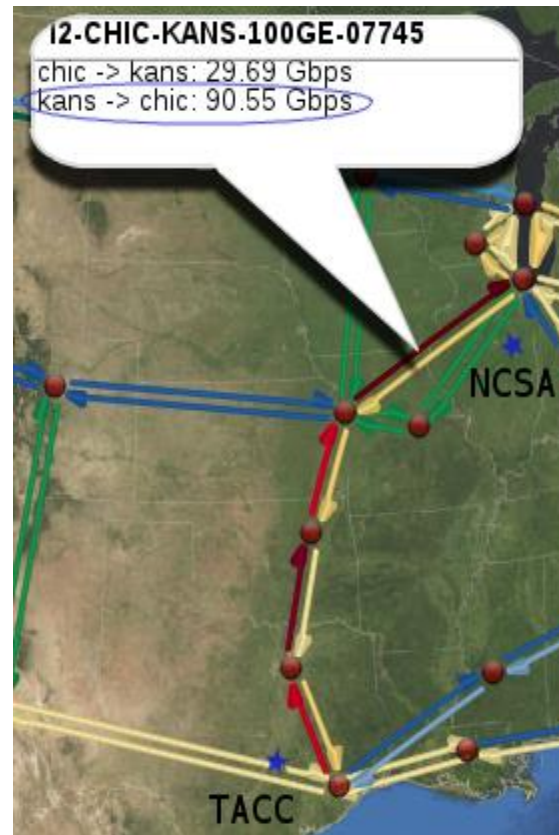
■ Hardware ■ Kernel ■ Transfer Application

AFTER



■ Hardware ■ Kernel ■ Transfer Application

- Performance of in Development FASP 0 Protocol
 - 65 Gbit/s Wire Rate transfer, TACC to NCSA
 - 61 Gbit/s Effective rate, over 90% of available bandwidth utilized for transfer
- Eliminate traditional bottlenecks which impede the efficient transmission of data
- Single Stream Single Node Transport Solution
- 91 Gbit/s effective throughput within LAN environment with single Mellanox ConnectX®-4
 - 1 PB of data transferred every day
 - 675 GB per minute



Version 1: Optimized around major bottleneck areas

- Minimize memory usage (Zero Copy Transfer Solution)
- Optimize Network (DPDK Based Network Stack)
- Optimize IO (Direct I/O to disk, memory aligned around Page Boundaries)

Version 2: Optimize memory layout

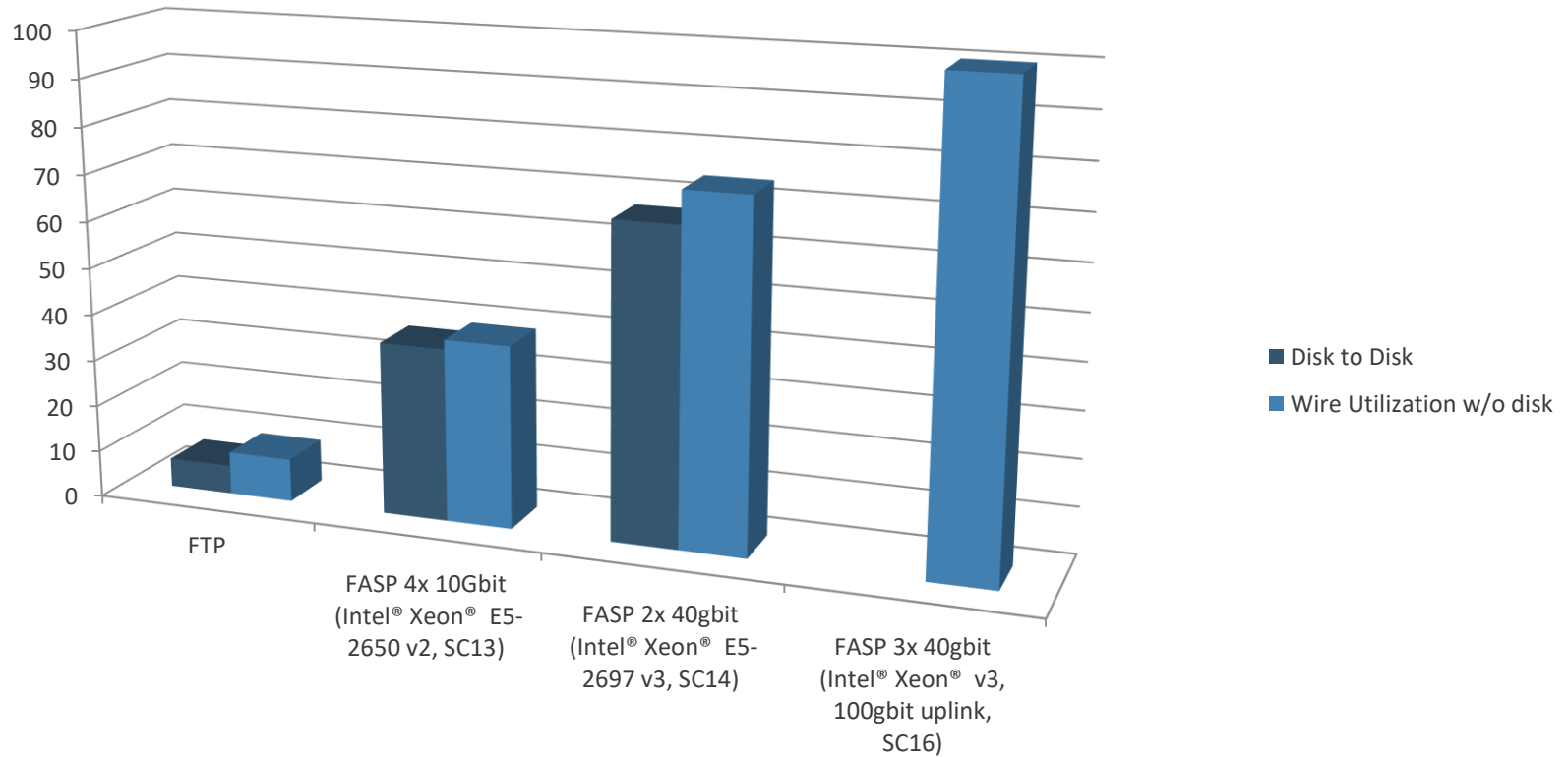
- Threading performance strongly correlated to how processors share data
- Improved throughput from ~40 to ~70 GBPS

Version 3: Improve Memory Locality

- Reduce how often threads need to synchronize
- Refine “Block” structure and how memory pages are owned by threads
- Improved Performance (120 GBPS) & Stability (> 1.2 PB transferred single session)

- Support for regular sockets
- Better negotiation of traffic flows
 - Dynamically change ports to better take advantage of Receive Side Scaling (RSS), a way of distributing network load across cores
 - Simplified session negotiation
- Integration into ASCP 4
(Two GA's of ASCP; *ASCP standard* and *ASCP 4*)
- Support stand alone DTLS authentication

Performance & Security



Remove Bottlenecks Typically Associated with Network Transfers

pcap_0_0.pcap [Wireshark 2.4.3 (v2.4.3)] [intel1]

File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help

Filter: Expression... Clear Apply Save

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	IntelCor_a1:05:69	Broadcast	ARP	64	Who has 10.211.0.1? Tell 10.211.0.2 [ETHERNET FRAME CHECK SE
2	0.398405016	10.211.0.2	10.212.0.2	DTLSv1.2	99	Application Data
3	0.399672909	10.211.0.2	10.212.0.2	DTLSv1.2	121	Application Data
4	0.406209269	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
5	0.406212321	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
6	0.406212531	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
7	0.406212774	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
8	0.406214765	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
9	0.406215003	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
10	0.406215198	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
11	0.406217163	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
12	0.406217345	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
13	0.406219305	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data
14	0.406219483	10.211.0.2	10.212.0.2	DTLSv1.2	1479	Application Data

▶ Frame 5: 1479 bytes on wire (11832 bits), 1479 bytes captured (11832 bits)

- ▶ Ethernet II, Src: IntelCor_a1:05:69 (3c:fd:fe:a1:05:69), Dst: Cisco_7c:ea:73 (00:3a:7d:7c:ea:73)
- ▶ Internet Protocol Version 4, Src: 10.211.0.2, Dst: 10.212.0.2
- ▶ User Datagram Protocol, Src Port: 11265, Dst Port: 777
- ▼ Datagram Transport Layer Security
 - ▼ DTLSv1.2 Record Layer: Application Data Protocol: Application Data
 - Content Type: Application Data (23)
 - Version: DTLS 1.2 (0xfefd)
 - Epoch: 6
 - Sequence Number: 1
 - Length: 1424
 - Encrypted Application Data: 6316f2d71d376e0d529c10f37edb5be914e16aebd2179453...

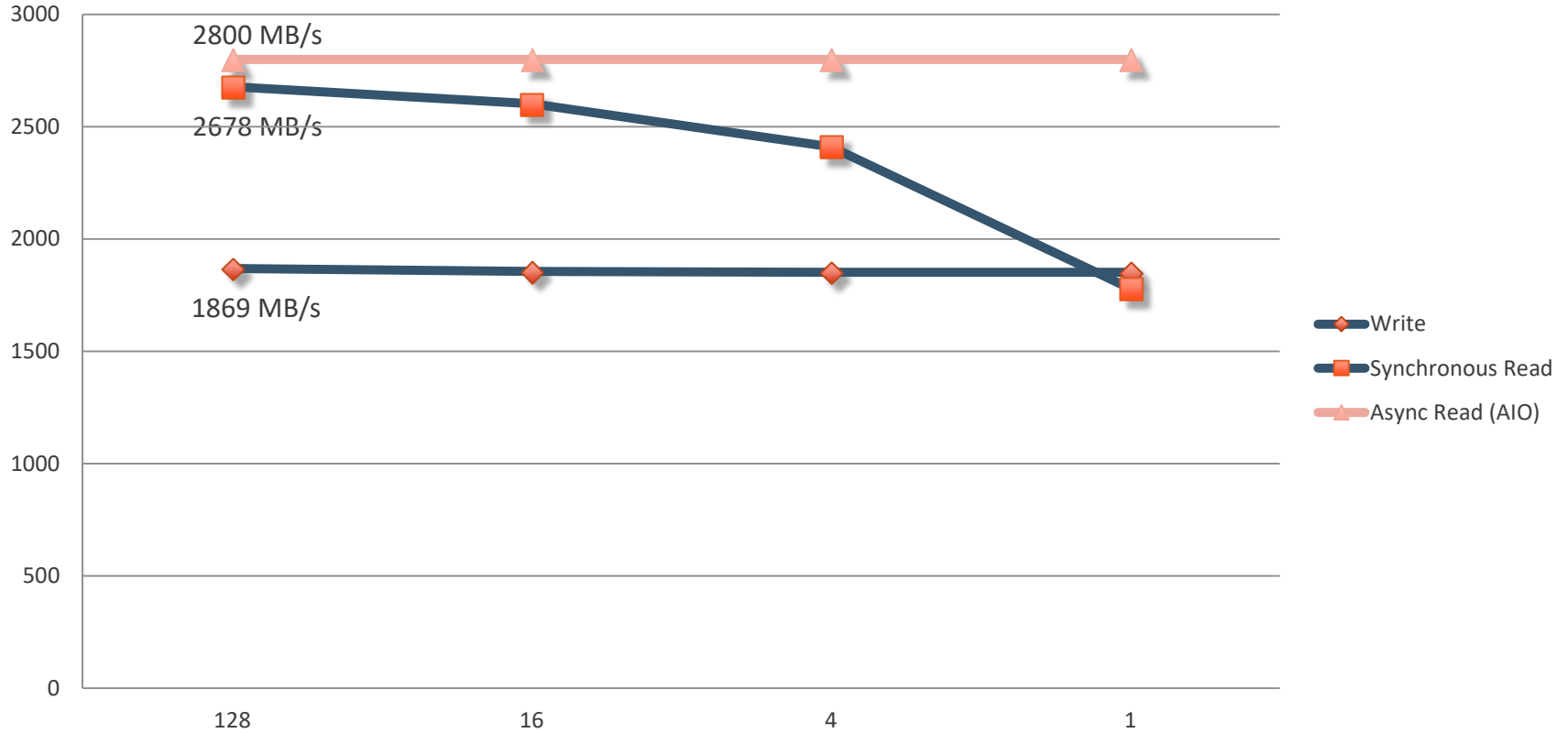
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0000 00 3a 7d 7c ea 73 3c fd fe a1 05 69 08 00 45 00  .:}] .s<. ...i..E.
0010 05 b9 be a7 00 00 40 11 00 00 0a d3 00 02 0a d4  .....@. ....
0020 00 02 2c 01 03 09 05 a5 a3 b4 17 fe fd 00 06 00  .,.....
0030 00 00 00 00 01 05 90 63 16 f2 d7 1d 37 6e 0d 52  .....c ....7n.R
0040 9c 10 f3 7e db 5b e9 14 e1 6a eb d2 17 94 53 84  .~[...].S...S.
0050 42 83 09 e2 e1 f6 07 bb c4 14 fc b5 5e 2f c1 2d  B.....^/..
0060 22 db c7 9b 92 89 53 66 4f 1f fc ef 30 a5 e0 04  "%....Sf 0...0...
0070 25 1f 8e 2b 4b 49 fa 9c cb 58 88 27 57 b2 fa 90  %..+KI...X.'W...
0080 66 3c 92 4c 69 4d e7 2e bb 79 99 99 93 1b 9d 14  f<.LiM...y.....
0090 3f 08 bb ba 84 ad 57 ed 95 f7 1b 65 1c a9 17 be  ?....W...e....
00a0 10 3d 18 a5 b6 87 d1 7e 59 69 bh 06 25 25 c5 97  -H...>Y!...%
  
```

File: "/tmp/pcap_o.o.pcap" 757 M... Packets: 507245 - Displayed: 507245 (100.0%) · Load time: 0:06.641 Profile: Default

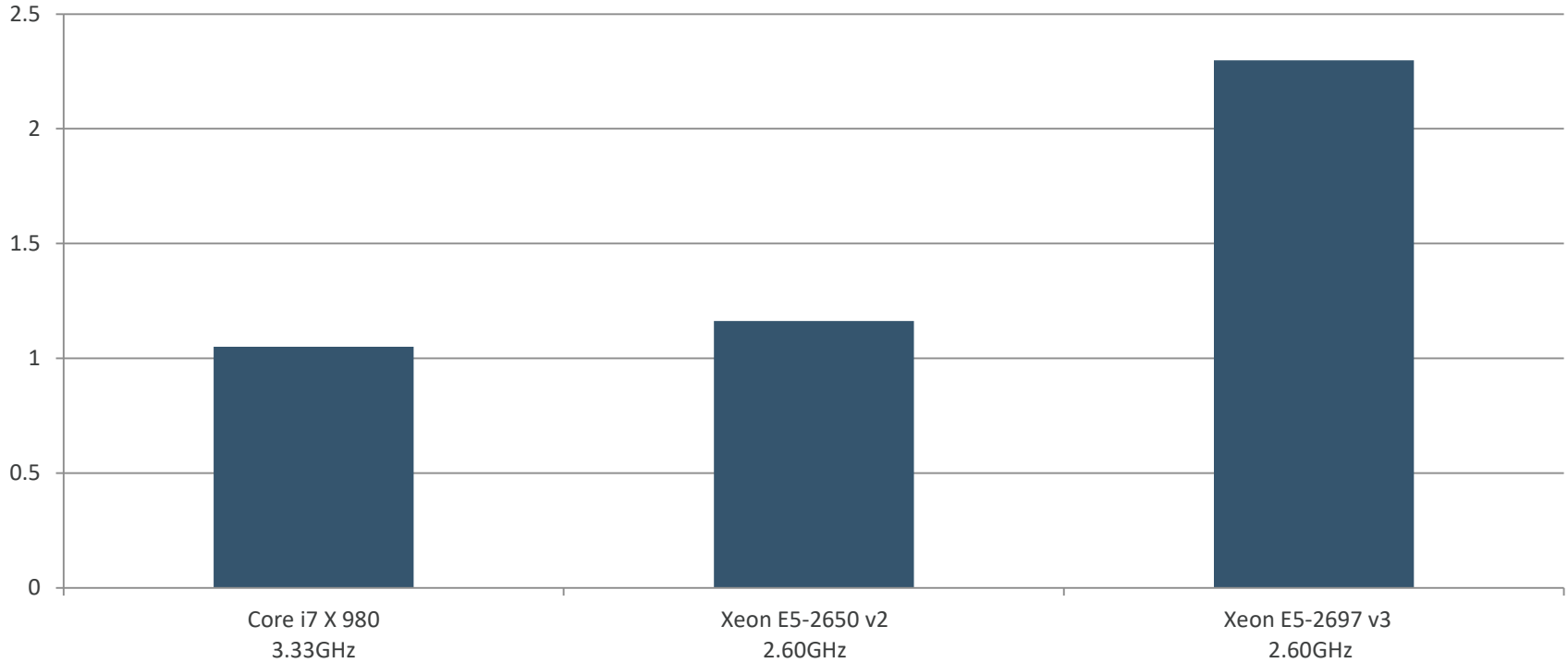
The screenshot displays the Aspera Enterprise Server interface. The main window shows a completed transfer of 1TB of data. The status is 'Complete', finished at 1:27 PM. The average speed was 69.5 Gbps, with an elapsed time of 2m 12s. The transfer information shows the source as 'This Computer' and the destination as 'aspi@10.0.200.253'. The speed graph shows a consistent rate of approximately 69.5 Gbps. The table below shows the transfer details.

Name	Source	Destination	Status	Speed	Size	Files	Remaining
1T	This Computer - 1T	aspi@10.0.200.253 - tmp	Complete	69.5 Gbps	1 TB	1 / 1	-



Performance relative to block size (SI Units in MB)

AES 128 GCM Encryption Rate in GB/s per Core



- Originally developed as a faster version of SCP.
 - Uses SSH for Authentication and Authorization.
 - Command line interface originally derived from SCP, but significantly different feature set to better support customer use cases
 - Default configuration tries to provide a good mix of security and usability.
- Language bindings for most languages (C, C++, Java, Python, .Net, Ruby, Javascript, Go, ...)
 - Persistent Sessions
 - Possible to bypass SSH Authentication, useful for propagating a session over TLS
 - Data transfer still over encrypted UDP. Internally encrypted UDP changing to DTLS 1.2 as FASP engine is upgraded across Aspera products.
 - Fully transparent; Language bindings fully expose transfer functionality
 - Streaming Capable
 - Cloud enabled (Direct to object store, SAAS, Cloud VMs)
- Web based collaboration Suite
 - AoC, Formerly Files, Aspera's SAAS offering
 - Faspex & Shares (Web based file collaboration), Console (Management), Orchestrator

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Aspera's FASPeSt FASP

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Thank you for your interest!