## Bioinformatics

## Phylogenetic trees

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## Overview

- Phylogentics
- Trees
- Definitions
- Properties
- Molecular clock[s]
- Tree [re]construction
- Distance methods
- UPGMA
- Neighbour-joining
- Character methods
- Maximum parsimony
- Maximum likelihood
- Assessing trees
- Bootstrapping trees


## Phylogenetic Trees

- Phylogeny : The evolutionary history and line of descent of a species
- Phylogenetic Tree : A diagram setting out the genealogy of a species
- Purpose
- to reconstruct the correct genealogical ties between related objects
- To estimate the time of divergence between them since they last shared a common ancestor
- Objects typically are protein or nucleic acid sequences


## What is phylogeny good for?

- Evolutionary history ("tree of life")
- Population history
- Rates of evolutionary change
- Origins of diseases
- Prediction of sequence function
- Can be applied to organisms, sequences, viruses, languages, etc.

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## Our place in Nature


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Phylogenetic Trees

## HIV: where did it come from and how is it transmitted?


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Yersinia pestis


Caenorhabitis elegans


Helicobacter pylori

rat
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Arabidopsis thaliana


Campylobacter jejuni


Mycobacterium leprae


Rickettsia prowazekii


Buchnerasp. APS


Chlamydia pneumoniae

mouse


Saccharomyces cerevisiae


Aquifex aeolicus


Vibrio cholerae


Neisseria meningitidis Z2491



Archaeoglobus
fulgidus


Drosophila melanogaster


Plasmodium falciparum


Bacillus subtilis

Borrelia burgorferi


Escherichia coli


Pseudomonas aeruginosa


Thermotoga maritima


Mycobacterium tuberculosis


Thermoplasma acidophilum


Ureaplasma urealyticum


Xylella fastidiosa

## Tree terminology



## Tree as text

## ((A, (B,C)), D)



## Definition of a tree

- A tree is either

A leaf
or
( LeftTree, RightTree)
where both LeftTree and RightTree are trees

## Trees with labels

- We can add data to
- Leaves
- Branches
- Nodes


(A:6, (B:2,C:2):4)


6: ( A , 4:( B, C ) )

## Isomorphic trees

## ((A, (B,C)), D)

( (A, (C,B)), D)

## Etc...

Derive [all] the isomorphic trees and draw them!

## Isomorphic trees

- Isotrees(T1,T2) /Two trees T1 and T2 are isomorphic/ if

T 1 is a leaf and T 2 is a leaf and $\mathrm{T} 1==\mathrm{T} 2$
OR

$$
\begin{aligned}
\mathrm{T} 1= & (\mathrm{L} 1, \mathrm{R} 1), \mathrm{T} 2=(\mathrm{L} 2, \mathrm{R} 2) \text { AND } \\
& \text { Isotrees(L1, L2) AND Isotrees(R1, R2) } \\
& \mathrm{OR} \\
& \operatorname{Isotrees(L1,~R2)~AND~Isotrees(R1,~L2)~}
\end{aligned}
$$

## Trees can be unrooted

A


Are there alternative rootings? Draw them...

## Evolution - basic concepts

- Mutation in DNA a natural evolutionary process
- DNA replication errors: (nucleotide)
- substitutions
- insertions
- deletions


## $\}$ indels

- Similarity between sequences
- clue to common evolutionary origin, or
- clue to common function
- This is a simplistic story: in fact the altered function of the expressed protein will determine if the organism will survive to reproduce, and hence pass on [transmit] the altered gene


## Evolution - related sequences



Back mutations - mutate to a previous sequence

## Definitions

- Phylum (phyla pl): A primary division of a kingdom, as of the animal kingdom, ranking next above a class in size.
- Phylogeny:
- the sequence of events involved in the evolutionary development of a species or taxonomic group of organisms
- Evolutionary relationships within and between taxonomic levels, particularly the patterns of lines of descent. Phylogenetics -The taxonomical classification of organisms based on their degree of evolutionary relatedness. Phylogenetic tree - A variety of dendrogram (diagram) in which organisms are shown arranged on branches that link them according to their relatedness and evolutionary descent.
- Phylogentics: study of evolutionary relationships
- Phylogenetic analysis: the means of inferring or estimating these relationships
- Taxonomy: The science of naming and classifying organisms 1. [n] practice of classifying plants and animals according to their presumed natural relationships
- 2. [n] (biology) study of the general principles of scientific classification
- 3. [n] a classification of organisms into groups based on similarities of structure or origin etc
- Taxon: any named group of organisms
- Species: taxonomic group whose members can interbreed (but more or less able to...)


## Definitions (cont)

- Clade: a group of biological taxa or species that share features inherited from a common ancestor; A monophyletic taxon; a group of organisms which includes the most recent common ancestor of all of its members and all of the descendants of that most recent common ancestor. From the Greek word "klados", meaning branch or twig.
- Tree : A data structure consisting of nodes which may contain other nodes via its branches. Unlike a tree in nature, the root node is usually represented at the top of the structure and does not have a parent node. All other nodes have a single parent. Nodes having no child nodes are called leaf nodes
- Dendrogram: Any branching diagram (or tree) (cf. cladogram, phylogram, phenogram); A dendrogram is a 'tree-like' diagram that summaries the process of clustering. Similar cases are joined by links whose position in the diagram is determined by the level of similarity between the cases. A treelike figure used to represent graphically a hierarchy. (dendron - greek 'tree')


## Homologues

- Homologues: sequences that have common origins but may or may not share common activity
- Orthologues: homologues produced by speciation. Tend to have similar function
- Paralogues: homologues produced by gene duplication. Tend to have different functions
- Xenologues: homologues resulting from horizontal gene transfer between 2 organisms.


## Tree of orthologues based on a set of $\alpha$-haemoglobins



## Tree of paralogues: human haemoglobins and myoglobin



## Orthologues \& paralogues



## Horizontal (lateral) gene transfer!

Bacteria can take up and spread genes in different ways

- Uptake of naked DNA directly from surroundings
- Obtain genes from infecting viruses
- Take up genes through cross-species mating


## Horizontal gene transfer



## Phylogenetic Trees

Approaches to reconstructing phylogenetic trees

- Distance based methods
- Maximum parsimony methods
- Maximum likelihood methods


## Phylogenetic analysis - 4 steps

1. Alignment / distance computation
2. Determine the data model
3. Tree building
4. Tree evaluation

## Tree building methods

1. Distance-based: Transform the data into pairwise distances (dissimilarities), and then use a matrix during tree building. E.g. data from from immunology, nucleic acid hybridization, and breeding experiments, are automatically expressed as pair-wise distances but most character data need to be mathematically transformed into distances.
2. Character-based: Use the aligned characters, such as DNA or protein sequences, directly during tree inference - based on substitutions.

## Tree building algorithms

- Distance-based:
- input = evolutionary distance data (sequence edit dist; melting temp DNA hybridisations, strength of antibody cross reactions..)
- Goal: construct weighted tree whose pairwise distances agree with evolutionary distances
- Ultrametric distance data - elegant solution: UPGMA
- Additive data - efficient solution: NJ
- FM - Fitch-Margoliash
- ME - minimum evolution
- Not additive - reconstruction not guaranteed; find trees whose distances "best approximate" given data
- Rooting trees


## Metric, ultrametric

- A metric on a set of objects $O$ is given by the assignment of a real number $d(x, y)$ for every pair of objects $x, y$ in $O$ such that

```
d(x,y)>0 for }x\not=
d(x,y)=0 for }x=
d(x,y)=d(y,x) for all x,y
d(x,y) \leqd(x,z)+d(y,z) for all }x,y,z\mathrm{ (triangle inequality)
```

- In addition, for an ultrametric
$d(x, y) \leq \max (d(x, z), d(y, z))$
- An ultrametric tree is characterised by the condition
$d(x, y) \leq d(x, z)=d(y, z)$ (for any 3 points their pairwise distances are all equal, or 2 are equal and 1 is smaller)

Ages \{Fred:20, Mary:22, Jane:24\} e.g. add, subtract - are these metric / ultrametric / ...?

## How many pairwise comparisons?

|  | A | B | C | D | E |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A |  |  |  |  |  |
| B |  |  |  |  |  |
| C |  |  |  |  |  |
| D |  |  |  |  |  |
| E |  |  |  |  |  |

- How many for whole table?
- Identity: $(x, y)=0$
for $x=y$
- Symmetry: $d(x, y)=d(y, x)$ for all $x, y$


## Metric, Ultrametric trees



Distance matrix

|  | A | B | C |
| :---: | :---: | :---: | :---: |
| A |  |  |  |
| $B$ |  |  |  |
| C |  |  |  |

Tree
(c) David Gilbert 2008


Phylogenetic Trees

## Ultrametric trees

Def: Given D a symmetric matrix $n$ by $n$ of real numbers; an ultrametric tree for D is a rooted tree T with the following properties:

1. T contains $n$ leaves, each labelled by a unique row $D$
2. Each internal node of T is labelled by one entry from D and has at least 2 children
3. Along any path from the root to the leaf, the numbers labelling the internal nodes strictly decrease
4. For any two leaves $\mathrm{i}, \mathrm{j}$ of $\mathrm{T}, \mathrm{D}(\mathrm{i}, \mathrm{j})$ is the label of the least common ancestor of i and j in T

Def: A min-ultrametric tree for D is a rooted tree T with all the properties of an ultrametric tree except that (3) is changed to:
3. Along any path from the root to the leaf, the numbers labelling the internal nodes strictly increase

## Symmetric matrix \& Ultrametric tree

|  | A | B | C | D | E |
| :--- | :--- | :--- | :--- | :--- | :--- |
| A | 0 | 10 | 10 | 6 | 4 |
| B |  | 0 | 4 | 10 | 10 |
| C |  |  | 0 | 10 | 10 |
| D |  |  |  | 0 | 6 |
| E |  |  |  |  | 0 |



## Obtaining pairwise distances

- Melting temp DNA hybridisations,
- Strength of antibody cross reactions
- Protein structure comparison
- Sequences: edit distance and variations
- $\mathrm{d}(\mathrm{i}, \mathrm{j}):$ fraction $f$ of sites $u$ where residues $\mathrm{x}_{\mathrm{iu}} \mathrm{x}_{\mathrm{ju}}$ differ (assume an alignment)
- For 2 unrelated sequences : random substitutions cause $f$ to approach fraction of differences expected by chance, hence:
- Jukes-Cantor distance $\mathrm{d}(\mathrm{i}, \mathrm{j})=-3 / 4 * \log (1-4 f / 3)$
- Tends to $\infty$ as equilibrium value of $f(75 \%$ residues differ) reached


## Gene sequence distance programs

- Blast (!) [but 'score' increases with similarity...]
- Clustal[w] - fasta format, but outputs a tree
- Phylip - its own format - sequences must be of the same length...hence use clustalw to align the sequences first [output in Phylip format, no tree]


## Clustering methods: UPGMA

## Unweighted Pair Group Method using Arithmetic Averages

Define the distance $\mathrm{d}(\mathrm{i}, \mathrm{j})$ between two clusters Ci and Cj to be the average distance between pairs of sequences from each cluster:

$$
d(i, j)=\frac{1}{\left|C_{i}\right| *\left|C_{j}\right|} \sum_{p \in C_{i}, q \in C_{j}} d(p, q)
$$

where $|\mathrm{Ci}|$ and $|\mathrm{Cj}|$ denote the number of sequences in clusters i and j respectively Note that if $\mathrm{Ck}=\mathrm{Ci} \cup \mathrm{Cj}$ and Cl is any other cluster then

$$
d(k, l)=\frac{d(i, l) *\left|C_{i}\right|+d(j, l) *\left|C_{j}\right|}{\left|C_{i}\right|+\left|C_{j}\right|}
$$

## UPGMA algorithm

## Initialisation:

- Assign each sequence i to its own cluster Ci
- Define one leaf T for each sequence, and place it height zero


## Iterate

- Determine the two clusters $\mathrm{i}, \mathrm{j}$ for which $\mathrm{d}(\mathrm{i}, \mathrm{j})$ [distance $\mathrm{i}-\mathrm{j}]$ is minimal.
(If there are more than two such pairs, pick one at random)
- Define a new cluster k by $\mathrm{Ck}=\mathrm{Ci} \cup \mathrm{Cj}$, and define $\mathrm{d}(\mathrm{k}, \mathrm{l})$ for all 1 by

$$
d(k, l)=\frac{d(i, l) *\left|C_{i}\right|+d(j, l) *\left|C_{j}\right|}{\left|C_{i}\right|+\left|C_{j}\right|}
$$

- Define a node k with daughter nodes i and j , and place it at height $\mathrm{d}(\mathrm{i}, \mathrm{j}) / 2$
- Add k to the current clusters and remove i and j


## Termination

- When only two clusters $\mathrm{i}, \mathrm{j}$ remain, place the root at height $\mathrm{d}(\mathrm{i}, \mathrm{j}) / 2$


## UPGMA example

12


## UPGMA example



## UPGMA example



## UPGMA example



## UPGMA example



## UPGMA tree construction example

|  | $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{D}$ | $\mathbf{E}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{A}$ | 0 | 8 | 4 | 6 | 8 |
| $\mathbf{B}$ |  | 0 | 8 | 8 | 4 |
| $\mathbf{C}$ |  |  | 0 | 6 | 8 |
| $\mathbf{D}$ |  |  |  | 0 | 8 |
| $\mathbf{E}$ |  |  |  |  | 0 |

## Like this?

Given: set of sequences

## Tree (re)construction

## Assume that

1 the pairwise sequence alignments provide a measure for the evolutionary distance between the sequences
2. the the resulting distance matrix constitutes an ultrameric (the ideal case),

Reconstruct the phylogenetic, ultrametric tree by the general clustering procedure:
Always pick the closest pair from the distance matrix and merge these two objects into one.

Schemes differ in how the distance between a newly formed object and the other objects is defined.
If object $x$ has been formed by merging $y$ and $z$, and $u$ is another object:

- Single linkage clustering: $\quad d(x, u)=\min (d(y, u), d(z, u))$
- maximal linkage clustering: $d(x, u)=\max (d(y, u), d(z, u))$
- average linkage clustering: $\quad d(x, u)=(d(y, u)+d(z, u)) / 2$


## Assumptions underlying UPGMA

- Molecular clock with constant rate
- Edge lengths correspond to times measured by clock
- Divergences of (sequences) assumed to occur at constant rate at all points in the tree, I.e.
- Sum of times down a path to the leaves from any node is the same, whatever choice of path
- Additivity of edge lengths
- Distance between any pair of leaves = sum of the lengths of the edges on the paths connecting them
- (automatic when UPGMA tree constructed)


## Incorrect tree construction by UPGMA???

|  | A | B | C |
| :---: | :---: | :---: | :---: |
| A | 0 | 3 | 6 |
| B |  | 0 | 5 |
| C |  |  | 0 |



## Incorrect tree construction by UPGMA



What will UPGMA give?

## Incorrect tree reconstruction by UPGMA

|  | $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{D}$ | $\mathbf{E}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{A}$ | 0 | 16 | 16 | 10 | 6 |
| $\mathbf{B}$ |  | 0 | 6 | 16 | 16 |
| $\mathbf{C}$ |  |  | 0 | 16 | 16 |
| $\mathbf{D}$ |  |  |  | 0 | 10 |
| $\mathbf{E}$ |  |  |  |  | 0 |

Construct the tree by UPGMA

## Addivity \& neighbour-joining

- Possible to have molecular clock property fail but addivity hold
- Use e.g. neighbour-joining algorithm
- Does not produce rooted trees...


## Tree whose closest leaves are not neighbours.



## Distances

Hence do not pick closest leaves, but subtract the averaged distances to all other leaves to compensate for long edges

$$
\begin{aligned}
& D(i, j)=d(i, j)-\left(r_{i}+r_{j}\right) \\
& \text { where } \\
& r_{i}=\frac{1}{|L|-2} \sum_{k \in L} d(i, k)
\end{aligned}
$$

## Test for additivity

Four point condition:
For every set of 4 leaves $i, j, k, l$
two of the distances
$\mathbf{d}(\mathbf{I}, \mathbf{j})+\mathbf{d}(\mathbf{k}, \mathbf{l}) \mathbf{d}(\mathbf{I}, \mathbf{k})+\mathbf{d}(\mathbf{j}, \mathbf{l}), \mathbf{d}(\mathbf{I}, \mathbf{l})+\mathbf{d}(\mathbf{j}, \mathbf{k})$ must be equal, and also larger to the third

Can use NJ even if lengths are not additive, but correct reconstruction not guaranteed.

## Rooting trees

- Add an outgroup
- E.g. add axolotl
- If no outgroup, then ad-hoc strategies
- See also TreeView by Rod Page (Glasgow)


## Quick demo

http://www.dina.dk/~sestoft/bsa/Match7Applet.html

## Parsimony

- Character-based method
- Finds the tree which can explain the observed sequences with a minimal number of substitutions

1. Assigns a cost to a given tree
2. Searches through all trees to find overall minimum of this cost


## Parsimony method

- Treats each site independently
- Given a topology \& assignment of residues $\rightarrow$ leaves:
- Basic step $=$ count minimal number of changes needed at 1 site


## Weighted parsimony

- Count number of subsitutions and
- Add costs $S(a, b)$ for each substitution of $a \operatorname{by} b$
- Aim : minimise this cost
- Weighted parsimony = traditional parsimony when $S(a, b)=1$ for all $a \neq b$
- Algorithm starts at the leaves and works up to the root : post-order traversal


## Traditional parsimony

- Just count the number of substitutions :
- Keep a list of minimal cost residues at each node, plus current cost


## Parsimony : rooted / unrooted

- Parsimony forumlated for rooted trees
- Minimal cost in trad.parsimony independent of location of root
- Root can be removed - number of trees to be searched over is reduced
- But easier to count costs in rooted tree


## Search stratgeies

- Improve over simple enumeration
- Stochastic methods: randomly swap branches on tree and chose altered tree if better than current
- Not guaranteed to find overall best tree; adding sequences in different orders can give different trees
- Build tree by adding edges 1 at time [ditto]
- Branch and bound:
- begin systematicall building trees with increasing number of leaves
- Abandon avenue whenever current incomplete tree has cost > smallest cost obtained so far for complete tree


## Assessing trees - bootstrap

- How much trees should be trusted
- Given dataset of alignment of sequences
- Generate artificial dataset (same size) by picking columns from alignment at random with replacement (given column in original can appear several times)
- Apply treebuilding to new dataset
- Repeat many (1000) times
- Frequency with which phylogenetic feature appears $\Rightarrow$ measure of confidence in feature


## Maximum likelihood

- Rank trees according to their
- likelihood P(data|tree)
- Posterior probability (Bayesian view) P(tree|data)
- Define \& compute the probability of a set of data given a tree
- Need a model of evolution - mutation and selection events that change sequences along the edges of a tree


## Maximum likelihood

- Maximum Likelihood
- Given a probabilistic model for nucleotide or aminoacid substitution, select the tree that has the highest probability of generating the observed data - sequences
- Give data $D$ and model $M$, find tree $T$ such that $P(D \mid T, M)$ maximised
- Actually what we really want is $P(T \mid D, M)$
- Not so easy to obtain as summation over all possible tree topologies required to obtain posterior

$$
P(T \mid D, M)=\frac{P(D \mid T, M) P(T, M)}{\sum_{T} P(D \mid T, M) P(T, M)}
$$

## Phylogenetic Trees

- Assumptions
- Different sites evolve independently
- Diverged sequences evolve independently after diverging
- Consider following tree
- Likelihood corresponds to $P(D \mid T, M)=P(X 1, X 2, X 3 \mid T, M)$



## Phylogenetic Trees

- Note that

$$
\begin{aligned}
&- P(D \mid T, M) \\
& \quad=P(X 1, X 2, X 3 \mid T, M) \\
& \quad= \sum_{\mathrm{X} 4, \mathrm{X} 5} P(X 1, X 2, X 3, X 4, X 5 \mid T, M) \\
& \quad=\sum_{\mathrm{X} 4, \mathrm{X} 5} P(X 1 \mid X 4) P(X 2 \mid X 4) P(X 3 \mid X 5) P(X 4 \mid X 5) P(X 5) \mid T, M
\end{aligned}
$$

- Efficient methods to obtain required summation for likelihood (Felsenstein Algorithm)
- Have now only obtained likelihood we require maximum of likelihood over all possible topologies and models (branch lengths)
- Approximate (stochastic )search methods required


## Maximum likelihood tree

- Search over tree topolgies
- For each topology: search over all possible lengths of edges
- ( $2 \mathrm{n}-3$ )!! rooted binary trees with n leaves
- $\mathrm{N}=10:: 2$ million, $\mathrm{N}=20:: 2.2 \times 10^{20}$
- Thus need optimisation techniques


## Probabilistic models of evolution

- Residue subsitution
- Deletion \& insertion of groups of residues
- Simple model:
- every site independent
- only substitutions


## Maximising likelihood

- Small number of sequences (2..5) : can enumerate all trees
- Write down likelihood as function of edge lengths
- Maximise likelihood by numerical technique (Kishano et al 1990, protein sequences - use PAM matrices)
- Larger number of sequences: Felsenstein's algorithm
- For protein sequences computationally demanding 20x20 substitution matrix
- Can use sampling methods


## More realistic evolutionary models

- Different rates at different sites
- Models with gaps


## Internet resources and programs

- Compilation of available programs:
http://evolution.genetics.washington.edu/phylip/software.html
- http://bioweb.pasteur.fr/seqanal/phylogeny/phylip-uk.html -- phylip , e.g. for nj
- http://evolution.genetics.washington.edu/phylip.html Phylip home page [widely used, poor interface...]
- http://www.hiv.lanl.gov/content/hiv-db/TREE_TUTORIAL/Tree-tutorial.html


## Phylogenetic Trees

Differences between methods

- UPGMA does not employ models of evolution
- Maximum Likelihood \& Maximum Parsimony employ models of evolution
- Maximum Likelihood claimed to be best for large evolutionary distances
- Very time consuming to build trees


## Summary

- Phylogenies
- Trees, representations, rooted, unrooted
- Metrics, ultrametrics, additivity, molecular clock
- Distance methods: UPGMA (ultrametric), NJ (additive; no root)...
- Tree rooting
- Some Programs

