# Chapter (3) Data Modeling Using the Entity-Relationship Model

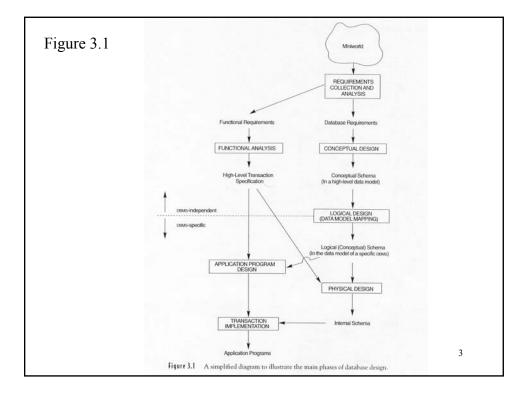
### **Objectives**

- Presenting the role of high-level conceptual data models in database design.
- Understanding the traditional approach of concentrating on the database structures and constraints during database design.
- Understanding the modeling concepts of the **Entity-Relationship (ER)** model.
- Understanding the basic data-structuring concepts and constraints of the ER model.
- Understanding the concepts of entities and attributes.

Conceptual modeling is an important phase in designing a successful database application. A **database application** refers to a particular database – for example, a BANK database – and the associated *programs* that implement the database queries and updates.

# **Main Topics**

- Using High-Level Conceptual Data Models for Database Design
- An Example Database Application
- Entity Types, Entity Sets, Attributes, and Keys
- Relationships, Relationship Types, Roles, and Structural Constraints
- Weak Entity Types
- Refining the ER Design for the COMPANY Database
- ER Diagrams, Naming Conventions, and Design Issues
- Summary



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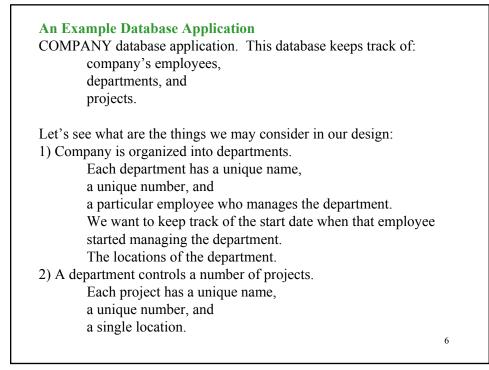
### **Database Design Process – cont.**

### Step (3) – Logical design (data model mapping)

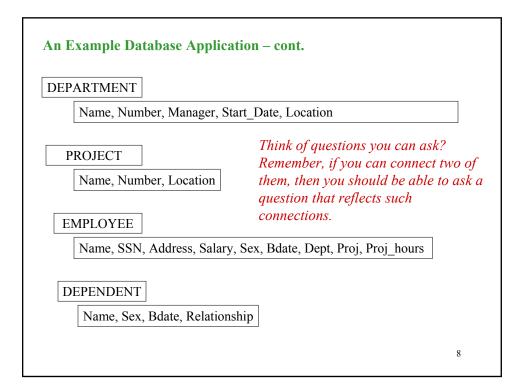
This step deals with the actual implementation of the database, using commercial DBMS.

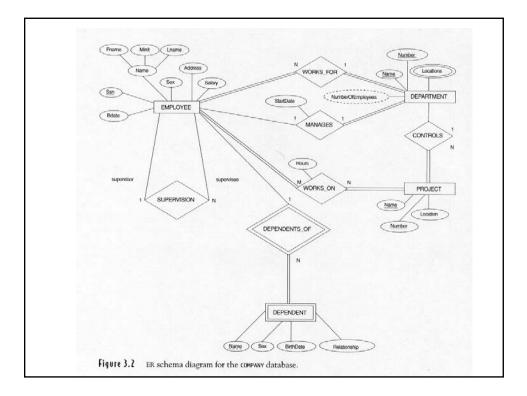
### Step (4) – Physical design

In this phase the internal storage structures, access paths, and file organizations for the database files are specified.



An Example Database Application – cont.
3) Employees information:
Each employee has a name,
a unique SSN,
address,
salary,
sex, and
DOB,
The department to which he is assigned,
The projects that he/she is working on,
Number of hours per week that he/she works on each project,
His/her direct supervisor.
4) Employees dependents:
Each dependent's first name,
sex,
DOB, and
relationship to the employee.
Having this in mind, can you design the schema?
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# Entity Types, Entity Sets, Attributes, and Keys

The ER model describes data as entities, relationships, attributes.

### **Entities and Their Attributes**

An entity is a "thing" in the real world with independent existence. Each entity has attributes which are the particular properties that describe it.

EXPLOYEE : Name, age, SSN, ... (Entity) (Attributes)

A particular entity has values for its attributes.

Several types of attributes occur in the ER model: simple vs. composite single-valued vs. multivalued stored vs. derived.

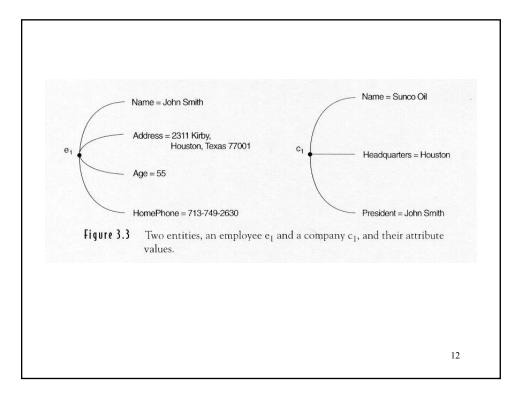
# **Composite vs. Simple (Atomic) Attributes**

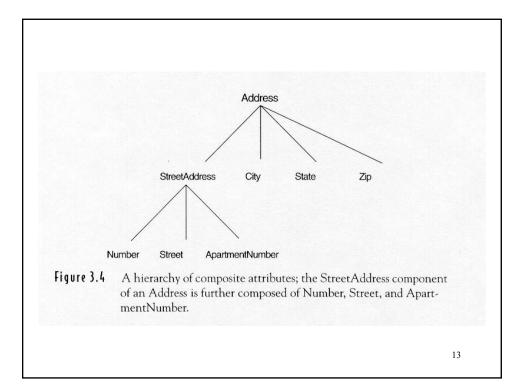
Composite attributes can be divided into smaller subparts.

Example: Address – Street Address, City, State, Zip Code, and Perhaps Country.

The value of a composite attribute is usually the concatenation of subfieelds.

There are times that the user refers to a composite attribute as a whole. In such cases, we do not need to separate the composite attribute in subfields.





### Single-valued vs. Multivalued Attributes

Most attributes have a single value for a particular entity. Example: Age is a single-valued attribute of an individual.

In some cases there are more than one possibilities for the value of an attribute.

Example: Students' class. Color of cars.

### **Stored vs. Derived Attributes**

In some cases the attribute does not exist in the database but its value can be derived by relating other attributes.

Example: Age is something that can be derived from today's date and the birth date. Today's date and the birth date are the two that are stored in the database.

### **Null Values**

In some cases a particular entity may not have an applicable value for an attribute. Example: Apartment Number in the ADDRESS entity.

# **Complex Attributes**

The composite and multivalued attributes can be nested in an arbitrary way.

Example:

Grouping components of a composite attribute between parentheses (), Separating the components with comma, and Displaying multivalued attributes between braces {}.

Example: A person with more than one address and phone number.

{AddressPhone( {Phone(AreaCode, PhoneNumber) }, Address(StreetAddress(Number, Street, ApartmentNumber), City, State, Zip) ) }

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### Entity Types, Entity Sets, Keys, and Value Sets

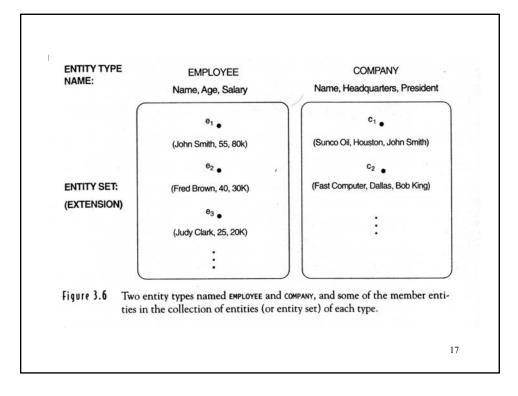
A database may contain a group of entities that are similar. Example: The EMPLOYEE and COMPANY entities both have the same attributes. These have similar names but will take different values.

An **entity type** defines a *collection* (or set) of entities that have the same attributes.

The collection of all entities of a particular entity type in the database at any point in time is called an **entity set**. The entity set usually have the same name as the entity type.

Example: EMPLOYEE refers to both a *type of entity* and the current *set of all employee entities* in the database.

An entity type describes the **schema** or **intention** for a *set of entities* that share the same structure. A collection of entities of a particular entity type are grouped into an entity set, called **extension** of the entity type. 16



# Key Attributes of an Entity Type

The **uniqueness constraint** or **key** is very important in the design of a database.

Example: SSN.

On the ER diagram, the <u>keys</u> are shown <u>underlined</u>. Some entities may have more than one key. In such cases, the combination of the two must be unique.

Example: LastName, FirstName .... (I know what you are going to ask) Example: Cars' registration DB

((ABS 123, TX), TK629, Ford Mustang, convertible, 1998, {red,black})

### Value Sets (Domains) of Attributes

Each simple attribute of an entity type is associated with a **value set** (or **domain** of values). The set consists of the values that may be assigned to that attribute for each individual entity.

Example: Age of employees 16-70.

A:  $E \rightarrow P(V)$ , means: an attribute A of entity E whose value set is V can be defined as a **function** from E to the power set P(V) of V. A power set P(V) of a set V is the set of all subsets of V.

The value of attribute A for entity e is represented as A(e).

For a composite attribute A, the value V is the Cartesian product of  $P(V_1)$ ,  $P(V_2)$ , ...,  $P(V_n)$ , where  $V_1$ ,  $V_2$ ,  $V_3$ , ...,  $V_n$  are the value sets of the simple component attributes that form A:

 $V = P(V_1) \mathbf{x} P(V_2) \mathbf{x} P(V_n) \mathbf{x} \dots \mathbf{x} P(V_n).$ 

Example:

Address: Street Address, P. O. Box, City, ZipCode, Country.

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# **Initial Conceptual Design of the COMPANY Database** We defined the entity types for the COMPANY database. DEPARTMENT DeptName, DeptNumber, Location, Manager, ManagerStartDate Key Key Multiple PROJECT PrjName, PrjNumber, Location, ControllingDepartment\* Key Key **EMPLOYEE** Name, Address, Salary, EmpSSN, Sex, Bdate, DeptNumber, Supervisor Both Composite DEPENDENT SSN,Name, Sex, Bdate, Relationship 20

### DEPARTMENT

*Name, Number, {Location}, Manager, ManagerStartDate* 

**PROJECT** *Name, Number, Location, Controlling Department* 

### **EMPLOYEE**

Name (Fname, Minit, Lname), SSN, Sex, Address, Salary, BirthDate, Department, Supervisor, {WorksOn(Project, Hours) }

### DEPARTMENT

Employee, DependentName, Sex, BirthDate, Relationship

Figure 3.8 – Preliminary design of entity types for the COMPANY database

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**Relationships, Relationship Types, Roles, and Structural Constraints** Whenever an attribute from one entity type refer to an attribute in another entity type, some *implicit relationships* among the various entity types exist.

We need to discuss the:

- Relationship Types, Sets, and Instances
- *Relationship Degree (binary, ternary), Role, and Recursive Relationships*
- Constrains on Relationship Types
- Attributes of Relationship Types

# **Relationship** Types

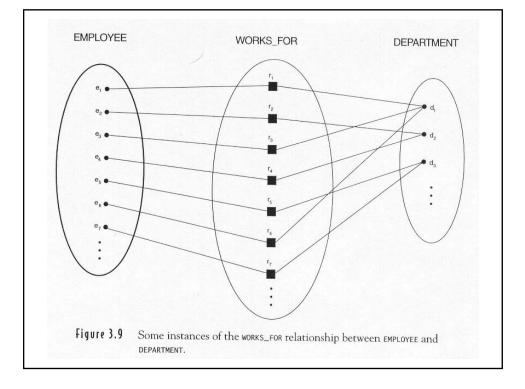
A relationship type R among n entity types  $E_1, E_2, ..., E_n$  defines a set of associations or a **relationship set** among entities from these types.

Mathematically, the relationship set *R* is a set of **relationship instances**  $r_i$ , where each  $r_i$  associates *n* individual entities  $(e_1, e_2, ..., e_n)$ , and each entity  $e_1$  in  $r_1$  is a member of entity type  $E_j$ ,  $1 \le j \le n$ .

Informally, each relationship instance  $r_i$  in R is an association of entities, where the association includes exactly one entity from each participating entity type.

Each  $r_i$  represents the fact that the entities participating in  $r_i$  are related in some way in the corresponding miniworld situation.

Example: What is the relationship between EMPLOYEE and DEPARTMENT? Is order is important ?

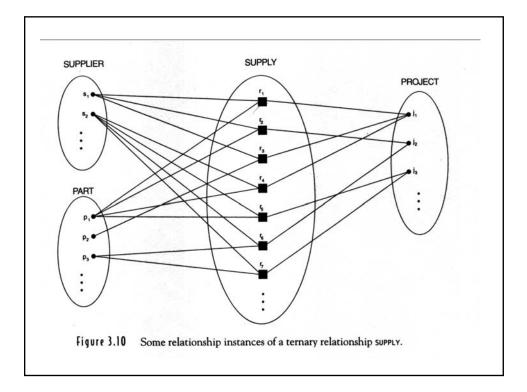


### **Relationship Degree, Role Names, and Recursive Relationships**

Degree of Relationship Type is the number of participating entity types.

Example: The WORKS\_FOR relationship is of degree two, because 2 entities (EMPLOYEE and DEPARTMENT) are participating. A relationship of degree two is called **binary**.

A relationship with degree three is called **ternary**. An example of a ternary relationship is SUPPLY shown on the figure on next page.



# **Relationship as Attributes**

It is sometimes convenient to think of a relationship type in terms of attributes.

An attribute called Department of the EMPLOYEE entity type whose value for each employee entity is the *department entity* that the employee works for.

Hence, the value set for this Department attribute is the *set of all DEPARTMENT entities*.

When we think of a binary relationship as an attribute, we always have two options.

Can you think of an example in the COMPANY DB?

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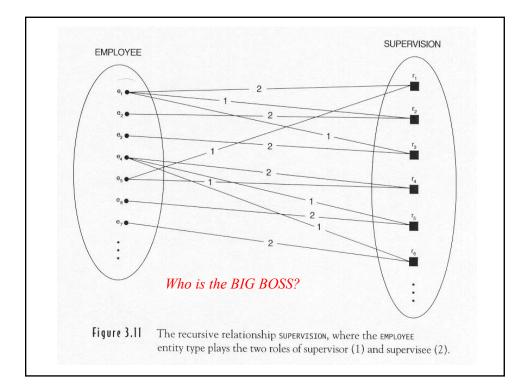
# **Role Name and Recursive Relationships**

Each entity type that participates in a relationship type plays a particular **role** in the relationship.

The **role name** signifies the role that a participating entity from the entity type plays in each relationship instance. It helps to explain what the relationship means.

Example: in the WORKS\_FOR relationship type, EMPLOYEE plays the role of *employee* or *worker* and *DEPARTMENT* plays the role of *department* or *employer*.

In some cases the *same* entity type participate more than once in a relationship type in *different roles*. In such cases, the role name becomes essential for distinguishing the meaning of each participation. Such relationships types are called **recursive relationships**.



# **Constrains on Relationship Types**

Relationship types usually have certain constraints that limit the possible combinations of entities that may participate in the corresponding relationship set.

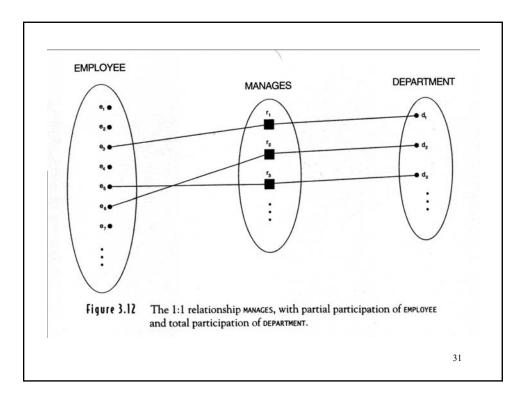
The constraints are determined from the miniworld situation that the relationship represent.

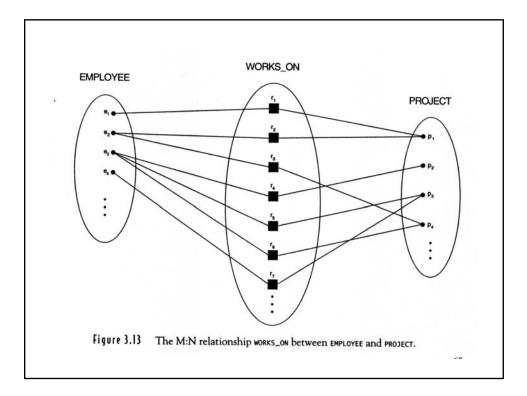
Example: Suppose each employee must work for only one department.

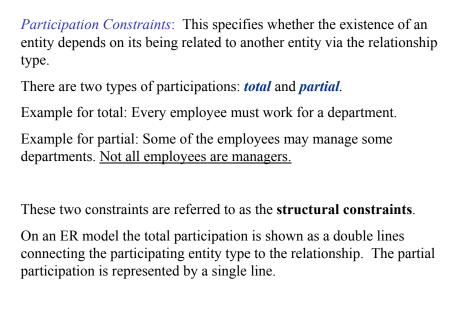
Two main types of relationship constraints: *cardinality ratio* and *participation*.

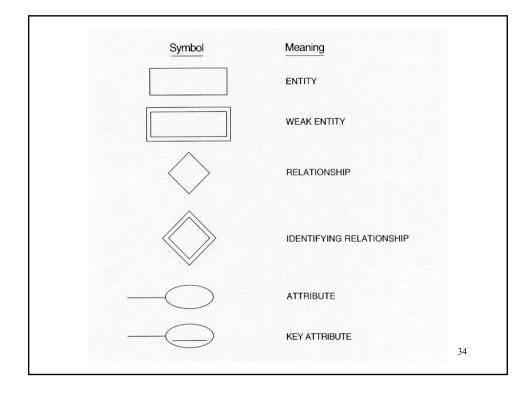
*cardinality ratio:* for a binary relationship specifies the number of relationship instances that an entity can participate in. Example: *DEPARTMENT: EMPLOYEE* is of cardinality ratio 1:N. What does this mean?

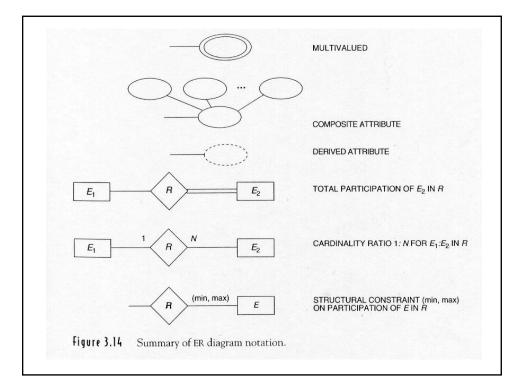
Other possibilities are: 1:1, 1:N, N:1, and M:N (N:N).

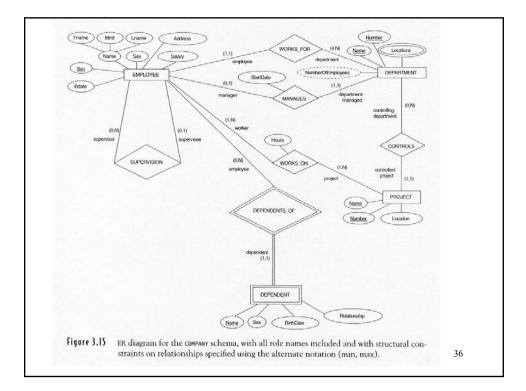












# **Attributes of Relationship Types**

Relationships can also have attributes.

Example: Suppose you add an attribute, Hours, to the WORKS\_FOR relationship for number of hours that the employee works on a particular project.

OR

Add the starting date, StartDate, to the MANAGES relationships.

The 1:1 or 1:N relationship types can be migrated to one of the participating entity types.

Example: The StartDate attribute for the MANGES relationship can be an attribute of either EMPLOYEE or DEPARTMENT.

For a 1:N relationship type, a relationship attribute can be migrated only to the entity type at the N-side of the relationship. Example: The WORKS\_FOR relationship can include the StartDate. This can be added to the EMPLOYEE table as an attribute.

# Weak Entity Types

Entity types that do not have key attributes of their own are called **weak types**.

The **regular entity types** that do have a key attribute are sometimes called **strong entity types**.

Entities belonging to a weak entity type are identified by being related to specific entities from another entity type in combination with some of their attribute values. We call this other entity type the **identifying** or **owner entity type**. It is also called **parent entity type**.

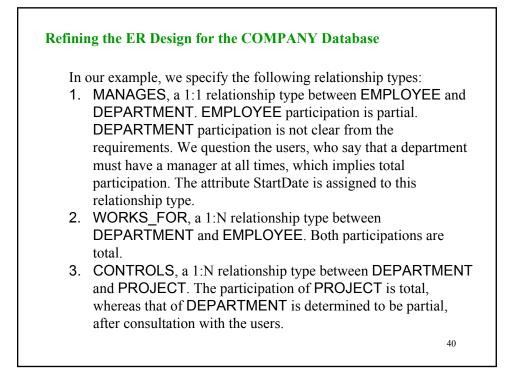
A weak entity type always has a *total participation constraint*. Example: A DRIVER\_LICENSE entity cannot exist unless it is related to a PERSON entity. Note, DRIVER\_LICENSE contains its unique key. Example: DEPENDENT entity ... how ?

# Weak Entity Types

A weak entity type normally has a **partial key**, which is the set of attributes that can uniquely identify weak entries that are <u>related to the same owner entity</u>.

Example: Suppose the children of each employee have unique first names. Thus, the Name in DEPENDENT entity can be used as the partial key.

On an ER model, the weak relationships and entities are represented with double lines.



# Refining the ER Design for the COMPANY Database

- 4. SUPERVISION, a 1:N relationship type between EMPLOYEE (in the supervisor role) and EMPLOYEE (in the supervisee role). Both participations are determined to be partial, after the users indicate that not every employee is a supervisor and not every employee has a supervisor.
- 5. WORKS\_ON, determined to be an M:N relationship type with attribute Hours, after the users indicate that a project can have several employees working on it. Both participations are determined to be total.
- 6. DEPENDENTS\_OF, a 1:N relationship type between EMPLOYEE and DEPENDENT, which is also the identifying relationship for the weak entity type DEPENDENT. The participation of EMPLOYEE is partial, whereas that of DEPENDENT is total.

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# **Design Choices for ER Conceptual Design** In general, the schema design process should be considered an iterative refinement process, where an initial design is created and then iteratively refined until the most suitable design is reached. Some of the refinements that are often used include the following: 1. A concept may be first modeled as an attribute and then refined into a relationship because it is determined that the attribute is a reference to another entity type. It is often the case that a pair of such attributes that are inverses of one another are refined into a binary relationship. 2. Similarly, an attribute that exists in several entity types may be refined into its own independent entity type. For example, suppose that several entity types in a UNIVERSITY database, such as STUDENT, INSTRUCTOR, and COURSE each have an attribute Department in the initial design; the designer may then choose to create an entity type DEPARTMENT with a single attribute DeptName and relate it to the three entity types (STUDENT, INSTRUCTOR, and COURSE) via appropriate relationships. Other attributes/relationships of DEPARTMENT may be discovered later. 42

### **Design Choices for ER Conceptual Design – cont.**

3. An inverse refinement to the previous case may be applied—for example, if an entity type DEPARTMENT exists in the initial design with a single attribute DeptName and related to only one other entity type STUDENT. In this case, DEPARTMENT may be refined into an attribute of STUDENT.

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### **Summary**

We learned about, modeling concepts of a high-level conceptual data model, ER model.

We talked about:

Simple or atomic Composite Multivalued

We also briefly discussed stored versus derived attributes. We then discussed the ER model concepts at the schema or "intension" level:

•Entity types and their corresponding entity sets.

- •Key attributes of entity types.
- •Value sets (domains) of attributes.
- •Relationship types and their corresponding relationship sets.
- •Participation roles of entity types in relationship types.

We presented two methods for specifying the structural constraints on relationship types. The first method distinguished two types of structural constraints:

> Cardinality ratios (1:1, 1:N, M:N for binary relationships) Participation constraints (total, partial)