## Combinational Circuit Design

### EE 200

### **Digital Logic Circuit Design**

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### Presentation Outline

- How to Design a Combinational Circuit
- Designing a BCD to Excess-3 Code Converter
- Designing a BCD to 7-Segment Decoder
- Hierarchical Design
- Iterative Design

### **Combinational Circuit**

✤ A combinational circuit is a block of logic gates having:

*n* inputs:  $x_1, x_2, \dots, x_n$ 

*m* outputs:  $f_1, f_2, \dots, f_m$ 

- Each output is a function of the input variables
- Each output is determined from present combination of inputs
- Combination circuit performs operation specified by logic gates



### How to Design a Combinational Circuit

#### 1. Specification

 $\diamond$  Specify the inputs, outputs, and what the circuit should do

#### 2. Formulation

♦ Convert the specification into truth tables or logic expressions for outputs

### 3. Logic Minimization

♦ Minimize the output functions using K-map or Boolean algebra

### 4. Technology Mapping

- ♦ Draw a logic diagram using ANDs, ORs, and inverters
- ♦ Map the logic diagram into the selected technology
- ♦ Considerations: cost, delays, fan-in, fan-out

#### 5. Verification

♦ Verify the correctness of the design, either manually or using simulation

### Designing a BCD to Excess-3 Code Converter

#### 1. Specification

- ♦ Convert BCD code to Excess-3 code
- ♦ Input: BCD code for decimal digits 0 to 9
- $\diamond$  Output: Excess-3 code for digits 0 to 9

#### 2. Formulation

- $\diamond$  Done easily with a truth table
- $\diamond$  BCD input: *a*, *b*, *c*, *d*
- $\diamond$  Excess-3 output: *w*, *x*, *y*, *z*
- ♦ Output is don't care for 1010 to 1111

BCD			Excess-3				
а	b	С	d	w	X	у	z
0	0	0	0	0	0	1	1
0	0	0	1	0	1	0	0
0	0	1	0	0	1	0	1
0	0	1	1	0	1	1	0
0	1	0	0	0	1	1	1
0	1	0	1	1	0	0	0
0	1	1	0	1	0	0	1
0	1	1	1	1	0	1	0
1	0	0	0	1	0	1	1
1	0	0	1	1	1	0	0
1010	t	0	1111	Х	Х	Х	Х

### Designing a BCD to Excess-3 Code Converter

#### **3. Logic Minimization using K-maps**



Minimal Sum-of-Product expressions:

w = a + bc + bd, x = b'c + b'd + bc'd', y = cd + c'd', z = d'

Additional 3-Level Optimizations: extract common term (c + d)

$$w = a + b(c + d)$$
,  $x = b'(c + d) + b(c + d)'$ ,  $y = cd + (c + d)'$ 

### Designing a BCD to Excess-3 Code Converter

#### 4. Technology Mapping

Draw a logic diagram using ANDs, ORs, and inverters

Other gates can be used, such as NAND, NOR, and XOR







#### Combinational Circuit Design

### Designing a BCD to Excess-3 Code Converter

#### 5. Verification

Can be done manually

Extract output functions from circuit diagram Find the truth table of the circuit diagram Match it against the specification truth table Verification process can be automated Using a simulator for complex designs

![](_page_7_Figure_5.jpeg)

#### Truth Table of the Circuit Diagram

BCD			Ex	CE	ess	s-3
abcd	c+d	b(c+d)	W	X	У	Z
0000	0	0	0	0	1	1
0001	1	0	0	1	0	0
0010	1	0	0	1	0	1
0011	1	0	0	1	1	0
0100	0	0	0	1	1	1
0101	1	1	1	0	0	0
0110	1	1	1	0	0	1
0111	1	1	1	0	1	0
1000	0	0	1	0	1	1
1001	1	0	1	1	0	0

### BCD to 7-Segment Decoder

- Seven-Segment Display:
  - ♦ Made of Seven segments: light-emitting diodes (LED)
  - ♦ Found in electronic devices: such as clocks, calculators, etc.

![](_page_8_Figure_4.jpeg)

- ✤ BCD to 7-Segment Decoder
  - ♦ Accepts as input a BCD decimal digit (0 to 9)
  - $\diamond\,$  Generates output to the seven LED segments to display the BCD digit
  - $\diamond$  Each segment can be turned on or off separately

![](_page_8_Picture_9.jpeg)

BCD to

![](_page_8_Picture_13.jpeg)

### Designing a BCD to 7-Segment Decoder

#### 1. Specification:

- $\diamond$  Input: 4-bit BCD (A, B, C, D)
- $\diamond$  Output: 7-bit (*a*, *b*, *c*, *d*, *e*, *f*, *g*)
- Display should be OFF for
  Non-BCD input codes

#### 2. Formulation

- $\diamond$  Done with a truth table
- $\diamond$  Output is zero for 1010 to 1111

# $\begin{array}{c} a \\ f \\ e \\ g \\ c \end{array}$

#### **Truth Table**

BCD input	7-Segment decoder
ABCD	abcdefg
0000	1 1 1 1 1 1 0
0001	0110000
0010	1 1 0 1 1 0 1
0011	1 1 1 1 0 0 1
0100	0110011
0101	1011011
0110	1011111
0111	1110000
1000	1 1 1 1 1 1 1
1001	1 1 1 1 0 1 1
1010 to 1111	0000000

### Designing a BCD to 7-Segment Decoder

#### 3. Logic Minimization Using K-Maps

![](_page_10_Figure_2.jpeg)

![](_page_10_Figure_3.jpeg)

 $\begin{array}{c|c} CD & \text{K-map for } c \\ AB & 00 & 01 & 11 & 10 \\ 00 & 1 & 1 & 1 & 1 \\ 01 & 1 & 1 & 1 & 1 \\ 11 & 1 & 1 & 1 & 1 \\ 10 & 1 & 1 & 1 & 1 \end{array}$ 

a = A'C + A'BD + AB'C' + B'C'D' b = A'B' + B'C' + A'C'D' + A'CD c = A'B + B'C' + A'DExtracting common terms

Let  $T_1 = A'B$ ,  $T_2 = B'C'$ ,  $T_3 = A'D$ 

Optimized Logic Expressions  $a = A'C + T_1 D + T_2 A + T_2 D'$   $b = A'B' + T_2 + A'C'D' + T_3C$   $c = T_1 + T_2 + T_3$  $T_1, T_2, T_3$  are shared gates

#### 3. Logic Minimization Using K-Maps

![](_page_11_Figure_2.jpeg)

### Designing a BCD to 7-Segment Decoder

#### 4. Technology Mapping

Many Common AND terms:  $T_0$  thru  $T_9$  $T_0 = A'C, T_1 = A'B, T_2 = B'C'$  $T_3 = A'D, T_4 = AB'C', T_5 = B'C'D'$  $T_6 = A'B'C$ ,  $T_7 = A'CD'$  $T_8 = A'BC'$ ,  $T_9 = A'BD'$ **Optimized Logic Expressions**  $a = T_0 + T_1 D + T_4 + T_5$  $b = A'B' + T_2 + A'C'D' + T_3C$  $c = T_1 + T_2 + T_3$  $d = T_4 + T_5 + T_6 + T_7 + T_8 D$  $e = T_{5} + T_{7}$  $f = T_4 + T_5 + T_8 + T_9$  $g = T_4 + T_6 + T_8 + T_9$ 

![](_page_12_Figure_3.jpeg)

### Verification Methods

- Manual Logic Analysis
  - ♦ Find the logic expressions and truth table of the final circuit
  - ♦ Compare the final circuit truth table against the specified truth table
  - ♦ Compare the circuit output expressions against the specified expressions
  - ♦ Tedious for large designs + Human Errors

#### Simulation

- ♦ Simulate the final circuit, possibly written in HDL (such as Verilog)
- $\diamond$  Write a test bench that automates the verification process
- ♦ Generate test cases for ALL possible inputs (exhaustive testing)
- ♦ Verify the output correctness for ALL input test cases
- ♦ Exhaustive testing can be very time consuming for many inputs

### Hierarchical Design

Why Hierarchical Design?

To simplify the implementation of a complex circuit

What is Hierarchical Design?

Decompose a complex circuit into smaller pieces called blocks

Decompose each block into even smaller blocks

- Repeat as necessary until the blocks are small enough
- Any block not decomposed is called a primitive block

The hierarchy is a tree of blocks at different levels

- The blocks are verified and well-document
- They are placed in a library for future use

### Example of Hierarchical Design

Top Level: 16-input odd function: 16 inputs, one output

♦ Implemented using Five 4-input odd functions

Second Level: 4-input odd function that uses three XOR gates

![](_page_15_Figure_4.jpeg)

### Top-Down versus Bottom-Up Design

- A top-down design proceeds from a high-level specification to a more and more detailed design by decomposition and successive refinement
- A bottom-up design starts with detailed primitive blocks and combines them into larger and more complex functional blocks
- Design usually proceeds top-down to a known set of building blocks, ranging from complete processors to primitive logic gates

### **Iterative Design**

- Using identical copies of a smaller circuit to build a large circuit
- Example: Building a 4-bit adder using 4 copies of a full-adder
- The cell (iterative block) is a full adder

Adds 3 bits:  $a_i$ ,  $b_i$ ,  $c_i$ , Computes: Sum  $s_i$  and Carry-out  $c_{i+1}$ 

Carry-out of cell *i* becomes carry-in to cell (*i*+1)

![](_page_17_Figure_6.jpeg)

### Full Adder

Full adder adds 3 bits: a, b, and c	Truth
Two output bits:	abc d
1. Carry bit: cout	000
2. Sum bit: <b>sum</b>	001
Sum bit is 1 if the number of 1's in	010
the input is odd (odd function)	011
$sum = (a \oplus b) \oplus c$	100
Carry bit is 1 if the number of 1's in	101
the input is 2 or 3	110
cout = a·b + (a⊕b)·c	111

#### **Table**

а	b	С	cout	sum
	-	•		
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1