CS 106B, Lecture 27 Advanced Hashing

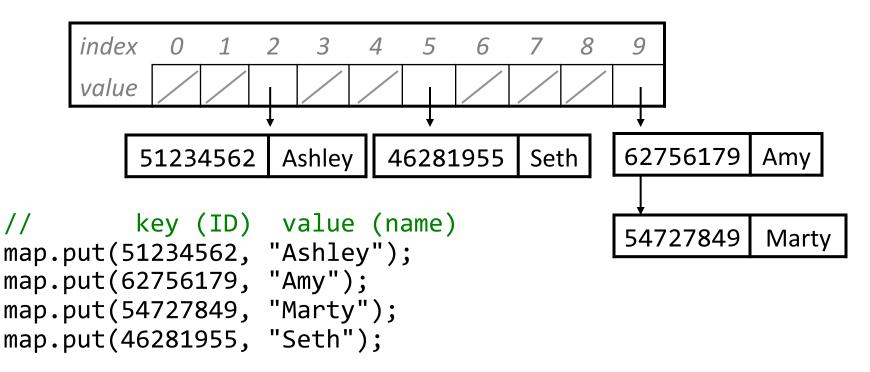
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Plan for Today

- Discuss how HashMaps differ from HashSets
- Another implementation for HashSet/Map: Cuckoo Hashing!
- Discuss qualities of a good hash function.
- Learn about another application for hashing: cryptography.

Hash map (15.4)

• A hash map is like a set where the nodes store key/value pairs:

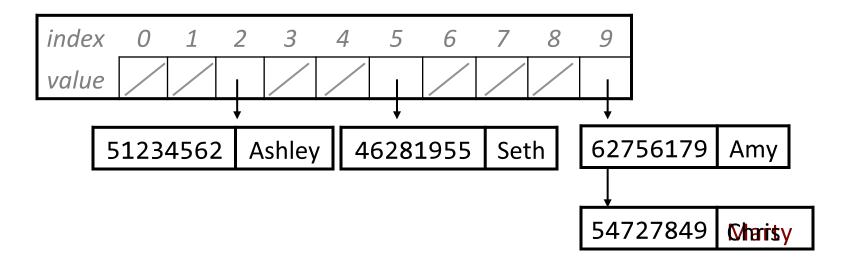


- Must modify the HashNode class to store a key and a value

Hash map vs. hash set

- The hashing is always done on the keys, not the values.
- The contains function is now containsKey; there and in remove, you search for a node whose key matches a given key.
- The add method is now **put**; if the given key is already there, you must replace its old value with the new one.

map.put(54727849, "Chris"); // replace Marty with Chris



Another Way to Hash

• Fun (but soon to be relevant) fact: cuckoo birds lay their eggs in other birds' nests



Source: wikimedia

- What if we made contains **really** fast (look at at most two elements, no matter what)?
- Idea: have two arrays that store elements, where each array has its own hash function
- Try hashing the element into both arrays, and put it in an empty space
- If no space is empty, kick out one of the existing elements and move it to the other array.
- Contains just checks the corresponding spot in both arrays
- Slower add, but faster contains

Insert: 3



Hash Function: 3x % 4

Hash Function: (2x + 1) % 4

Insert: 3



Hash Function: 3x % 4

Hash Function: (2x + 1) % 4

Insert: 6



Hash Function: 3x % 4

Hash Function: (2x + 1) % 4

Insert: 6



Hash Function: (2x + 1) % 4

Insert: 5



Hash Function: (2x + 1) % 4

Insert: 5



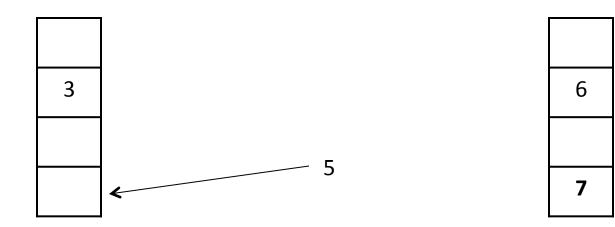
Hash Function: (2x + 1) % 4

Insert: 7



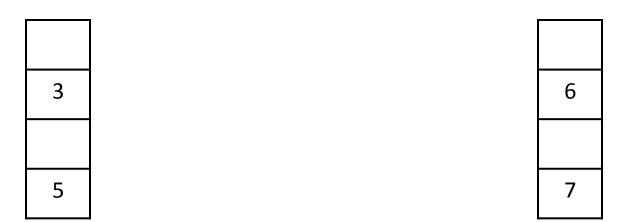
Hash Function: (2x + 1) % 4

Insert: 7



Hash Function: (2x + 1) % 4

Insert: 7



Hash Function: (2x + 1) % 4

Search for 7 (look in both arrays)



Hash Function: (2x + 1) % 4

- What are the advantages or disadvantages of cuckoo hashing versus resolving collisions through chaining?
- What do we need to watch out for? When should we rehash?

Announcements

- Calligraphy announcements
 - Should start the 3rd part today or tomorrow at the latest
 - Starter code and Windows please redownload
 - No late days may be used, no late submissions accepted
- Last class tomorrow go to poll.ly/#/LdVNgWyo/G6z0awRv
- Final is a on Saturday, at 8:30AM, in Cubberley Auditorium
 - Everything from the course through today is fair game, emphasis is on second half materials (starting with pointers)
 - More information:

https://web.stanford.edu/class/cs106b/exams/final.html

- Practice exam is online not guaranteed to match in format, etc.
- Wednesday and Thursday will be final review
- Please give us feedback! cs198.stanford.edu

Hashing strings

- It is easy to hash an integer i (use index *abs(i) % length*).
 - How can we hash other types of values (such as strings)?
- If we could convert strings into integers, we could hash them.
 - What kind of integer is appropriate for a given string?
 - Does it matter what integer we choose? What should it be based on?

index	0	1	2	3	4	5	6	7
character	'H'	'i'		'D'	'0'	'0'	'd'	'!'

hashCode consistency

• A valid hashCode function <u>must</u> be *consistent* (must produce same results on each call)

hashCode(x) == hashCode(x), if x's state doesn't change



hashCode and equality

- A valid hashCode function <u>must</u> be *consistent with equality*.
 - a == b must imply that hashCode(a) == hashCode(b).

Vector<int> v1; Vector<int> v2; v1.add(1); v2.add(3);

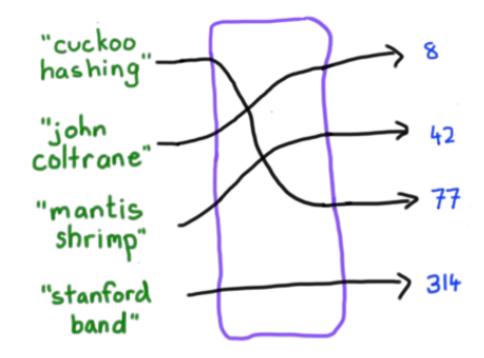
v1.add(3); v2.insert(0, 1);

```
// hashCode(v1) == hashCode(v2)
```

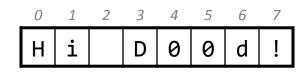
a != b does NOT necessarily imply that
hashCode(a) != hashCode(b) (why not?)

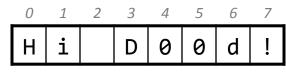
hashCode distribution

- A good hashCode function is *well-distributed*.
 - For a large set of distinct values, they should generally return unique hash codes rather than often colliding into the same hash bucket.
 - This property is desired but not required. Why?

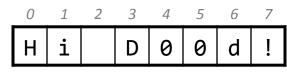


```
int hashCode(string s) { // #1
    return 42;
}
```



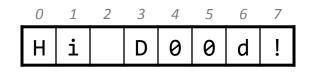


```
int hashCode(string s) { // #2
    return randomInteger(0, 9999999);
}
```



```
int hashCode(string s) { // #3
    return (int) &s; // address of s (a pointer)
}
```

```
int hashCode(string s) { // #4
    return s.length();
}
```



• **Q:** Is this a valid hash function? Is it good?

```
int hashCode(string s) { // #5
    if (s.length() > 0) {
        return (int) s[0]; // ascii of 1st char
    } else {
        return 0;
    }
}
```

2

3

D

4

0

5

0

6

d

7

1

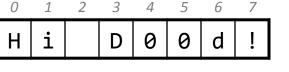
i

Η

- This function sums the characters' ASCII values.
 - Is it valid? Is it good?

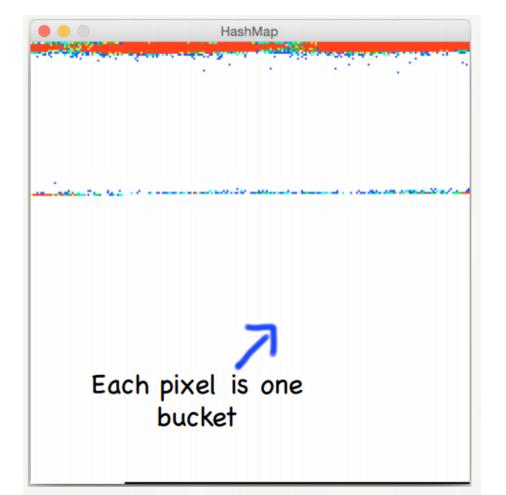
– What will collide?

int hashCode(string s) { // #6
 int hash = 0;
 for (int i = 0; i < s.length(); i++) {
 hash += (int) s[i]; // ASCII of char
 }
 return hash;
}</pre>



Measuring collisions

- Hash function = sum of characters of string.
- Add 50,000,000 article titles to a hash map with 50,000 buckets:





Idea: Weighted sum

hash = s[0] + s[1] + s[2] + ... + s[n]

- Instead of adding, let's give each character a **weight**.
 - Multiply it by increasing powers of some prime number; say, 31.
 - This helps spread the strings' hash codes over the range of int values.

hash = $s[0] + (31 * s[1]) + (31^2 * s[2]) + ... + (31^n * s[n])$

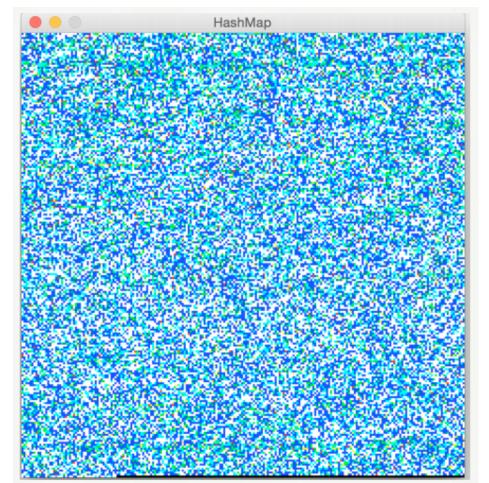
hashCode for strings

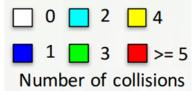
```
int hashCode(string s) {
    int hash = 5381;
    for (int i = 0; i < (int) s.length(); i++) {
        hash = 31 * hash + (int) s[i];
    }
    return hash;
}</pre>
```

- FYI: The above is the actual hash function used for strings in Java.
- As with any general hashing function, collisions are possible.
 - Example: "Ea" and "FB" have the same hash value.

Measuring collisions

- Hash function = sum of characters of string, multiplying by 31.
- Add 50,000,000 article titles to a hash map with 50,000 buckets:





Hashing structs/objects

- By default you cannot add your own structs/objects to hash sets.
 - Our libraries don't know how to hash these objects.

```
struct Point {
    int x;
    int y;
    ...
};
```

```
HashSet<Point> hset;
Point p {17, 35};
hset.add(p);
```

```
ERROR: no matching function for call to
'hashCode(const Point&)'
```

Hashing structs/objects

- To make your own types hashable by our libraries:
 - 1) Overload the == operator.
 - 2) Write a hashCode function that takes your type as its parameter.
 - "Add up" the object's state; scale/multiply parts to distribute the results.

```
struct Point {
    int x;
    int y;
    ...
};
int hashCode(const Point& p) {
    return 1337 * p.y + 31 * p.x;
}
```

bool operator ==(const Point& p1, const Point& p2) {
 return p1.x == p2.x && p1.y == p2.y;

Hashing and Passwords

- We want to store a file of user passwords
 - When a user types a password, see if it matches our file
- Problem: anyone who can see our file can get all the passwords

Password
password123
traceComics
ki88leLuv

Hashing and Passwords

- What if we stored a unique code for each password instead of the string?
 - Hashing!
- Extra requirements for the hash function:
 - Want a large number of possible values (hard to find collisions)
 - Can't find the password from the hash (one-way)
 - Generally use a different hash function (e.g. SHA-256)
- The need for salting

User	Password
Ashley	17851691385
Marty	63158910316
Amy	90713593110

Hashing and Data Integrity

- A common "attack" in cryptography is man-in-the-middle
- How can you ensure that a hacker didn't interfere with the data?
- Get the hash from a **trusted** source since hash functions only rarely have collisions, changes to data will lead to a different hash