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SOCKET PROGRAMMING

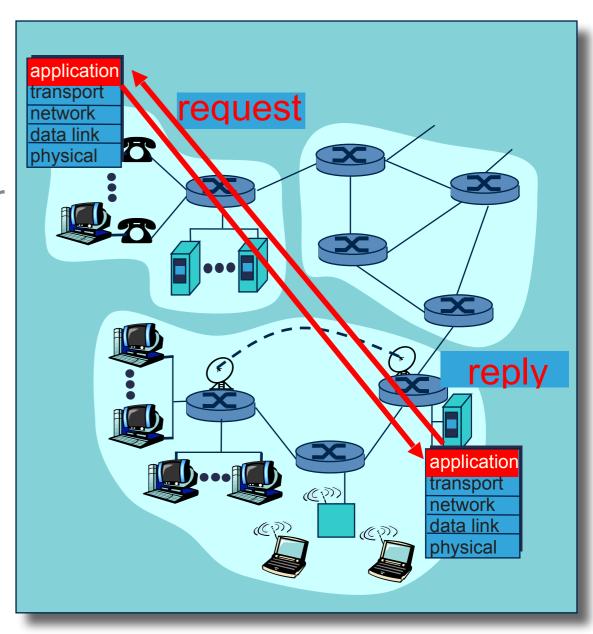
LECTURE OVERVIEW

- Application layer
 - Client-server
 - Application requirements
- Background
 - TCP vs. UDP
 - Byte ordering
- Socket I/O
 - ▶ TCP/UDP server and client
 - I/O multiplexing

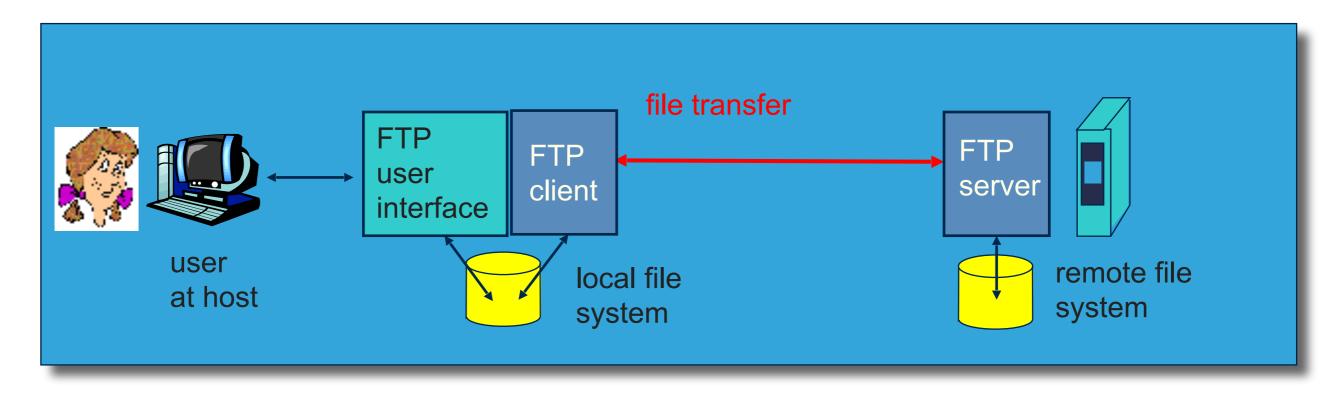
CLIENT-SERVER PARADIGM

Typical network app has two pieces: client and server

- Client
 - Initiates contact with server
 - Typically requests service from server
 - Client implemented in browser for web, mail reader for e-mail
- Server
 - Provides requested service to client
 - e.g. Sends web page, delivers e-mail



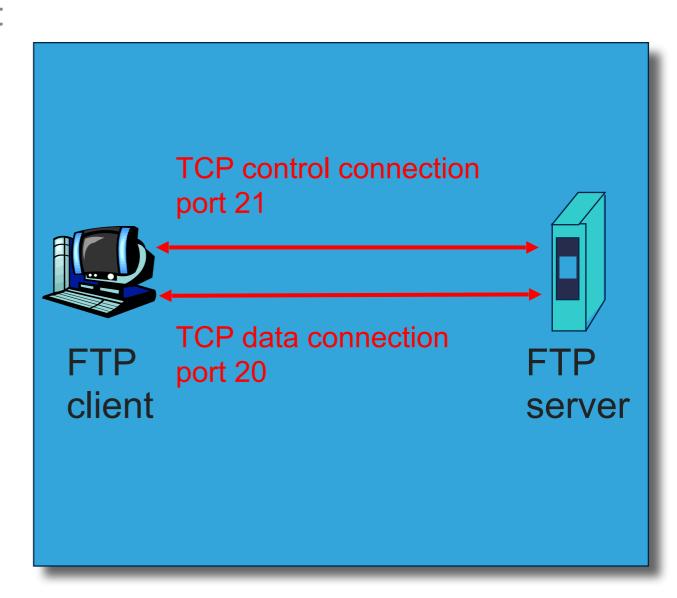
FTP: THE FILE TRANSFER PROTOCOL



- Transfer file to/from remote host
- Client/server model
 - Client: side that initiates transfer (either to/from remote)
 - Server: remote host
- ftp: RFC 959
- ftp server: port 21

SEPARATE CONTROL, DATA CONNECTIONS

- Ftp client contacts ftp server at port 21, specifying TCP as transport protocol
- Two parallel TCP connections opened:
 - Control: exchange commands, responses between client and server
 - Data: file data to/from server
 - Out-of-band protocol
- Ftp server maintains "state": current directory, earlier authentication



FTP COMMANDS, RESPONSES

Sample Commands:

sent as ASCII text over control channel

USER username

PASS password

LIST return list of files in current directory

RETR filename retrieves (gets) file

STOR filename stores (puts) file onto remote host

Sample Return Codes:

status code and phrase

331 Username OK, password required

125 data connection already open; transfer starting

425 Can't open data connection

452 Error writing file

TRANSPORT SERVICE REQUIREMENTS

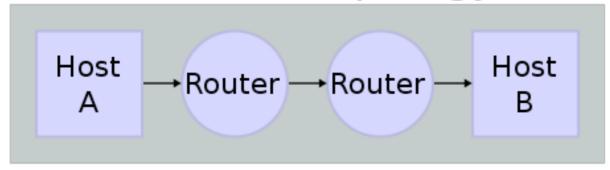
- Data loss
 - Some apps (e.g. audio) can tolerate loss
 - ▶ Other apps (e.g. file transfer, telnet) require 100% reliable transfer
- Timing
 - Some apps (e.g. games) require low delay to be effective
- Bandwidth
 - Some apps (e.g. multimedia) require minimum bandwidth to be effective
 - Some apps (e.g. "elastic apps") use whatever bandwidth they can

TRANSPORT SERVICE REQUIREMENTS

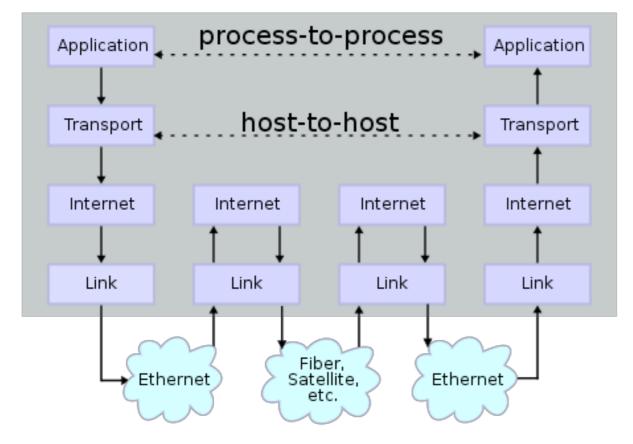
Application	Data loss	Bandwidth	Time Sensitive
file transfer	no loss	elastic	no
e-mail	no loss	elastic	no
web documents	no loss	elastic	no
real-time audio/	loss-tolerant	audio: 5Kb-1Mb	yes, 100 msec
video		video:10Kb-5Mb	
stored audio/video	loss-tolerant	same as above	yes, few secs
interactive games	loss-tolerant	few Kbps	yes, 100 msec
financial apps	no loss	elastic	yes and no

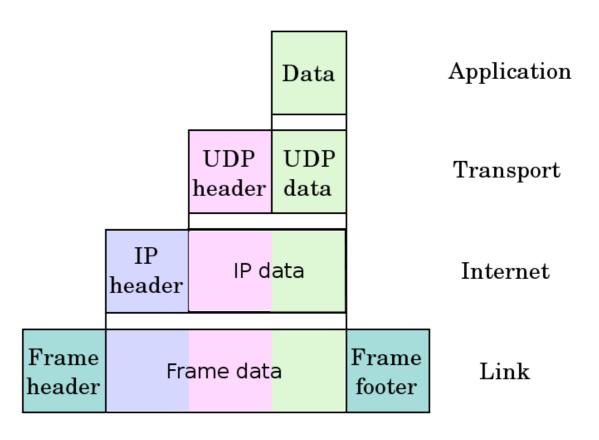
PACKET FORMAT

Network Topology

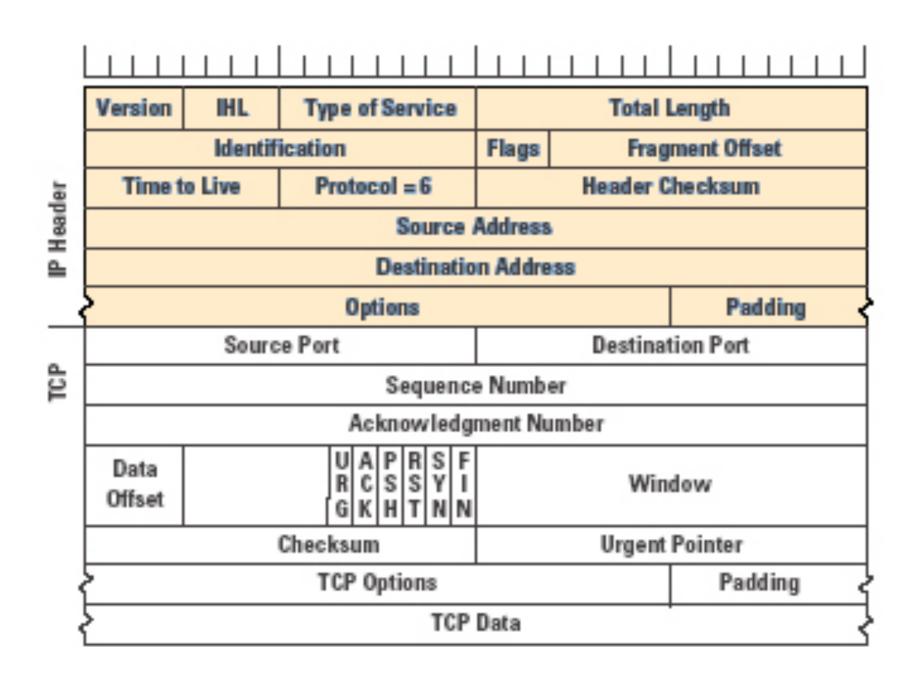


Data Flow





PACKET FORMAT

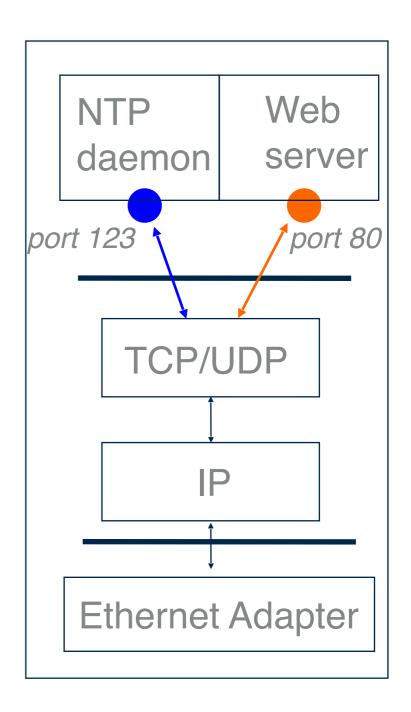


NAMES AND ADDRESSES

- Each attachment point on Internet is given a unique address
 - Based on location within network (like phone numbers)
- Humans prefer to deal with names not addresses
 - Domain Name Service (DNS) provides mapping of name to address
 - Name based on administrative ownership of host

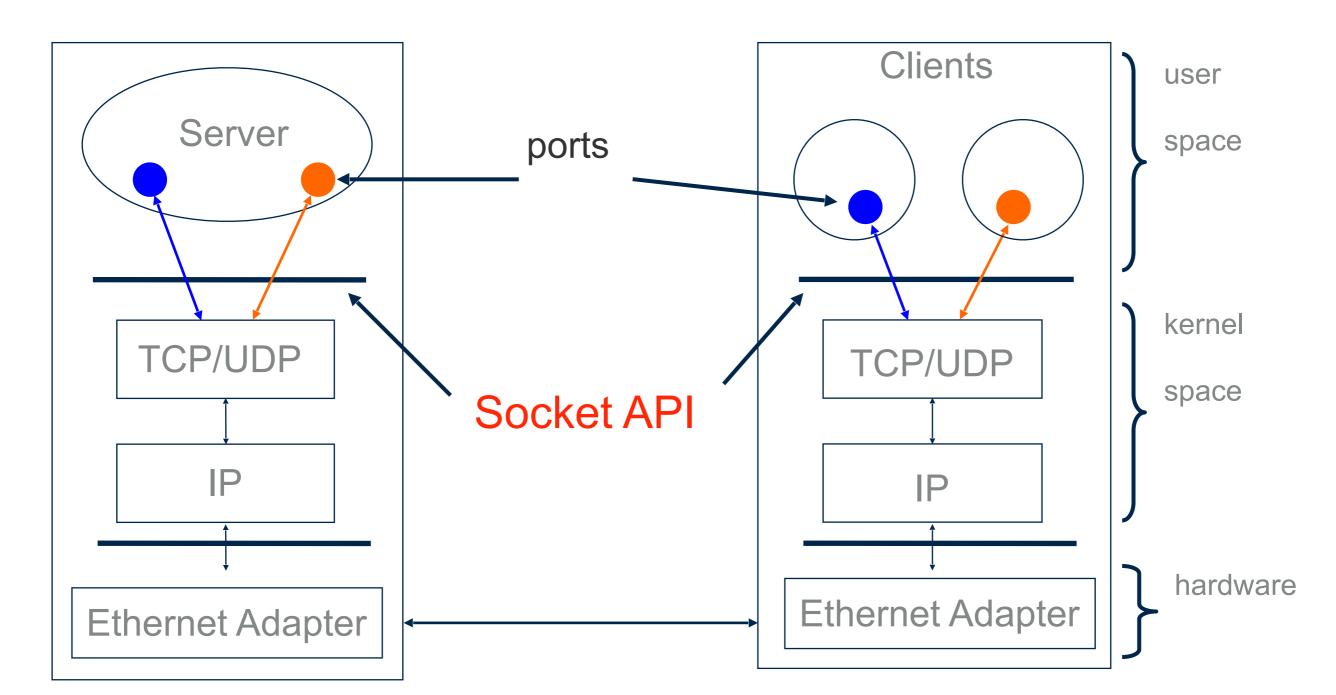
CONCEPT OF PORT NUMBERS

- Port numbers are used to identify "entities" on a host
- Port numbers can be:
 - Well-known (port 0-1023)
 - Assigned (port 1024-49151)
 - Dynamic or private (port 49152-65535)
- Servers/daemons usually use well-known ports
 - Any client can identify the server/service
 - HTTP = 80, FTP = 21, Telnet = 23, ...
- Other common services use assigned ports
- Clients should use dynamic ports
 - Assigned by kernel at runtime



SERVER AND CLIENT

Server and Client exchange messages over the network through a common Socket API

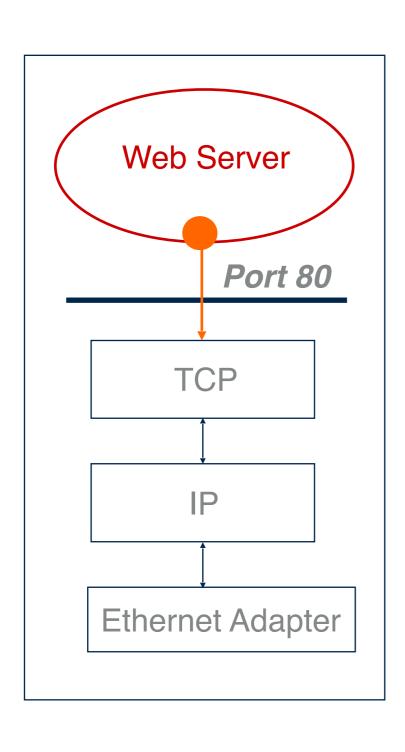


WHAT IS A SOCKET?

A socket is a file descriptor that lets an application read/write data from/to the network

- socket returns an integer (socket descriptor)
 - ▶ fd < 0 indicates that an error occurred
- ▶ AF_INET: associates a socket with the Internet protocol family
- SOCK_STREAM: selects the TCP protocol, SOCK_DGRAM: selects the UDP protocol

TCP SERVER



What does a web server need to do so that a web client can connect to it?

SOCKET I/O: SOCKET()

Since web traffic uses TCP, the web server must create a socket of type SOCK_STREAM

SOCKET I/O: BIND()

A socket can be bound to a port

```
int fd;
                                   /* socket descriptor */
struct sockaddr in srv;    /* used by bind() */
/* create the socket */
srv.sin family = AF_INET; /* use the Internet addr family */
srv.sin port = htons(80); /* bind socket 'fd' to port 80*/
/* bind: a client may connect to any of my addresses */
srv.sin addr.s addr = htonl(INADDR ANY);
if (bind (fd, (struct sockaddr*) &srv, sizeof (srv)) < 0) {
      perror("bind"); exit(1);
```

Still not quite ready to communicate with a client...

SOCKET I/O: LISTEN()

listen indicates that the server will accept a connection

Still not quite ready to communicate with a client...

SOCKET I/O: ACCEPT()

accept blocks waiting for a connection

- accept returns a new socket (newfd) with the same properties as the original socket (fd)
 - newfd < 0 indicates that an error occurred</p>

SOCKET I/O: ACCEPT() CONTINUED...

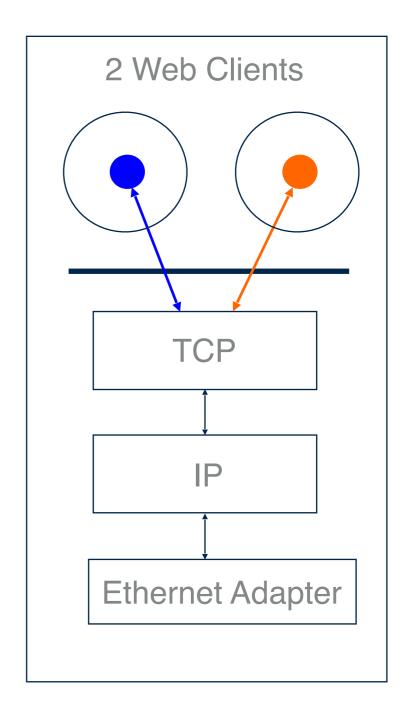
- How does the server know which client it is?
 - cli.sin_addr.s_addr contains the client's IP address
 - cli.sin_port contains the client's port number
- Now the server can exchange data with the client using read and write on the descriptor newfd
- Why does accept need to return a new descriptor?

SOCKET I/O: READ()

read *blocks* on data from the client but does not guarantee that sizeof(buf) is read

TCP CLIENT

How does a web client connect to a web server?



DEALING WITH IP ADDRESSES

▶ IP Addresses are commonly written as strings ("128.2.35.50"), but programs deal with IP addresses as integers.

Converting strings to numerical address:

```
struct sockaddr_in srv;
srv.sin_addr.s_addr = inet_addr("128.2.35.50");
if(srv.sin_addr.s_addr == (in_addr_t) -1) {
    fprintf(stderr, "inet_addr failed!\n"); exit(1);
}
```

Converting a numerical address to a string:

```
struct sockaddr_in srv;
char *t = inet_ntoa(srv.sin_addr);
if(t == 0) {
    fprintf(stderr, "inet_ntoa failed!\n"); exit(1);
}
```

TRANSLATING NAMES TO ADDRESSES

- getaddrinfo provides interface to DNS
- Returns addrinfo structs given a host and service
- getnameinfo provides host and service given addrinfo
- Functions are not IPv4 or IPv6 dependent

```
#include <netdb.h>
int st;
struct addrinfo *results; /*ptr to linked list of address info*/
struct addrinfo hints;
char *name = "www.cs.cmu.edu";
if (st = getaddrinfo(name, "80", &hints, &results) != 0) {
    fprintf(stderr, "getaddrinfo failed!\n"); exit(1);
}
```

SOCKET I/O: CONNECT()

connect allows a client to connect to a server

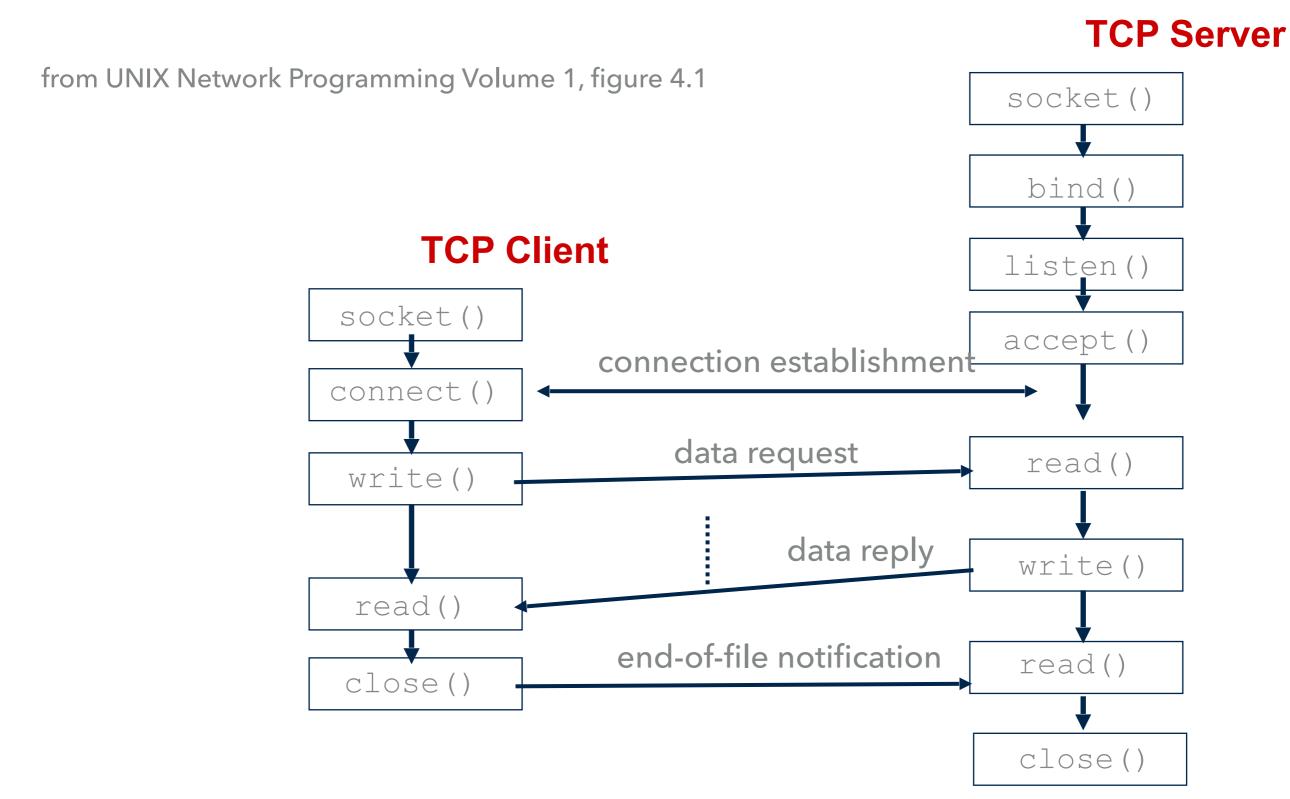
```
int fd;
                                   /* socket descriptor */
                                 /* used by connect() */
struct sockaddr in srv;
/* create the socket */
/* connect: use the Internet address family */
srv.sin family = AF INET;
/* connect: socket 'fd' to port 80 */
srv.sin port = htons(80);
/* connect: connect to IP Address "128.2.35.50" */
srv.sin addr.s addr = inet addr("128.2.35.50");
if(connect(fd, (struct sockaddr*) &srv, sizeof(srv)) < 0) {</pre>
       perror("connect"); exit(1);
```

SOCKET I/O: WRITE()

write can be used with a socket

```
int fd;
                                 /* socket descriptor */
                                 /* used by connect() */
struct sockaddr in srv;
                               /* used by write() */
char buf[512];
                                 /* used by write() */
int nbytes;
/* 1) create the socket */
/* 2) connect() to the server */
/* Example: A client could "write" a request to a server */
if((nbytes = write(fd, buf, sizeof(buf))) < 0) {</pre>
      perror("write");
      exit(1);
```

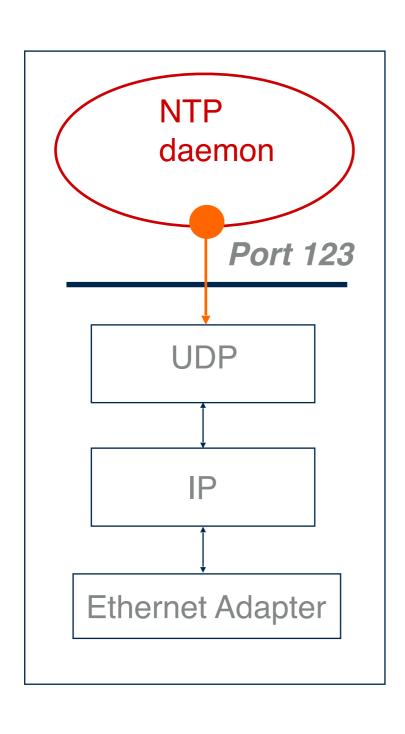
TCP CLIENT-SERVER INTERACTION



UDP PROPERTIES

- Does not assume any handshake or prior communication
- Stateless protocol with no information/session retention
- Uses datagrams or self-contained packets of information
 - No need for prior information exchange

UDP SERVER EXAMPLE



What does a UDP server need to do so that a UDP client can connect to it?

SOCKET I/O: SOCKET()

The UDP server must create a datagram socket...

- socket returns an integer (socket descriptor)
 - fd < 0 indicates that an error occurred</p>
- ► AF_INET associates the socket with the Internet protocol family
- SOCK_DGRAM selects the UDP protocol

SOCKET I/O: BIND()

A socket can be bound to a port

```
int fd;
                                   /* socket descriptor */
struct sockaddr in srv; /* used by bind() */
/* create the socket */
/* bind: use the Internet address family */
srv.sin family = AF INET;
/* bind: socket 'fd' to port 80*/
srv.sin port = htons(80);
/* bind: a client may connect to any of my addresses */
srv.sin addr.s addr = htonl(INADDR ANY);
if (bind (fd, (struct sockaddr*) &srv, sizeof (srv)) < 0) {
      perror("bind"); exit(1);
```

Now the UDP server is ready to accept packets...

SOCKET I/O: RECVFROM()

- read does not provide the client's address to the UDP server
- recvfrom receives messages from a socket

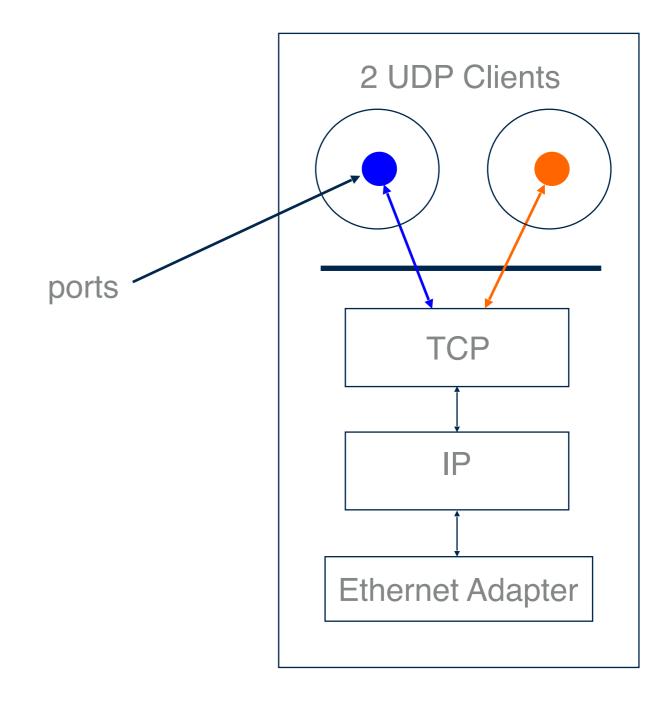
```
int fd;
                           /* socket descriptor */
struct sockaddr in srv; /* used by bind() */
                    /* used by recvfrom() */
struct sockaddr in cli;
                         /* used by recvfrom() */
char buf[512];
int nbytes;
                           /* used by recvfrom() */
/* 1) create the socket */
/* 2) bind to the socket */
nbytes = recvfrom(fd, buf, sizeof(buf), 0 /* flags */,
             (struct sockaddr*) &cli, &cli len);
if(nbytes < 0) {
     perror("recvfrom"); exit(1);
```

SOCKET I/O: RECVFROM() CONTINUED...

- The actions performed by recvfrom
 - Returns the number of bytes to read (nbytes)
 - Copies nbytes of data into buf
 - Returns the address of the client (cli)
 - Returns the length of cli (cli_len)

UDP CLIENT EXAMPLE

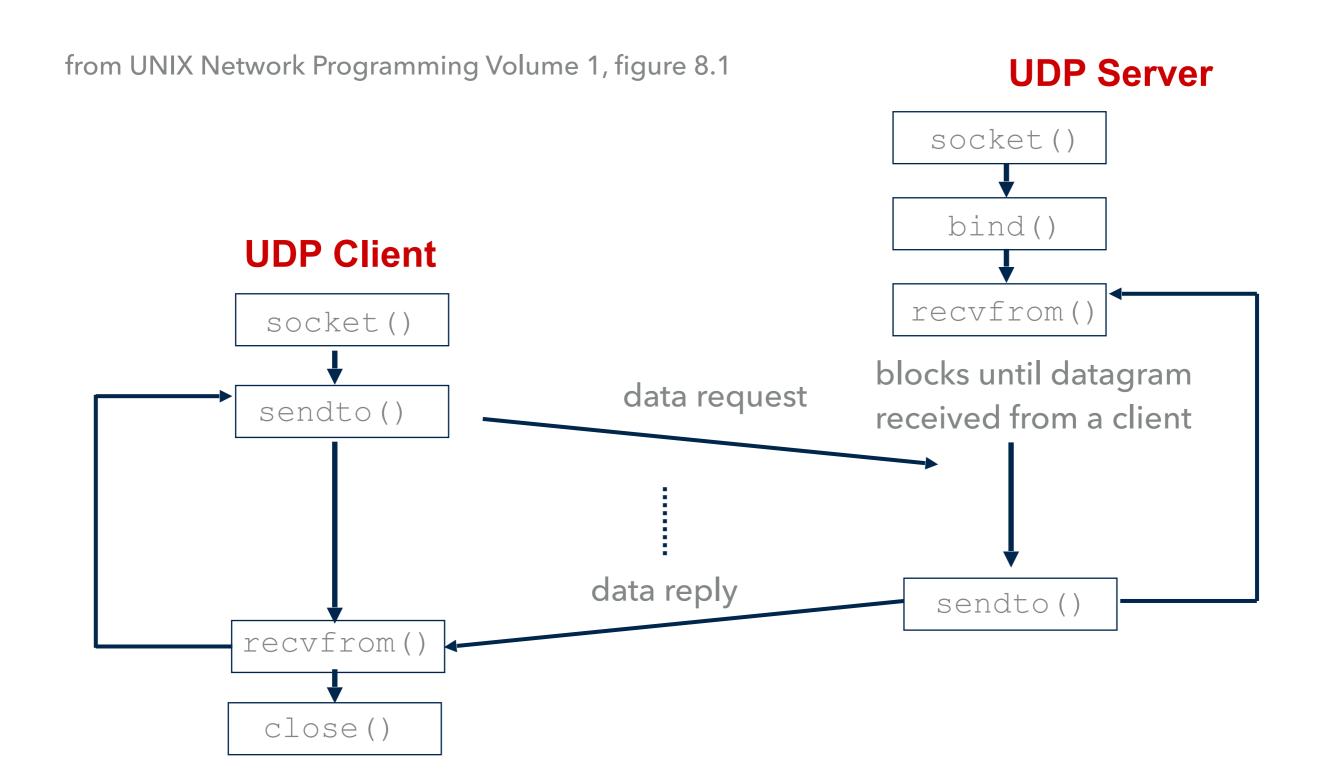
How does a UDP client communicate with a UDP server?



SOCKET I/O: SENDTO()

- write is not allowed
- UDP client does not bind a port number
 - Port number is dynamically assigned when the first sendto is called

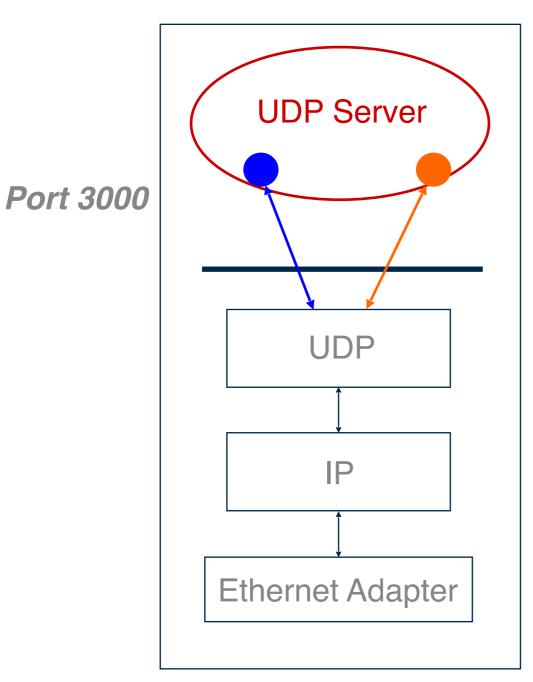
UDP CLIENT-SERVER INTERACTION



SIDE NOTE: UDP BROADCAST AND MULTICAST

- These examples have been point-to-point (one source, one destination) sending of data but UDP supports point-to-multipoint (one source, multiple destinations)
- May not work in all circumstances and primarily for LANs
 - Broadcast only supported in IPV4
 - Multicast not supported by all switches and hubs
 - Only way to do it across the Internet is with additional work-arounds
- ▶ IP Multicast added to IPV4 and fully integrated in IPV6
 - Primarily for multimedia content

THE UDP SERVER



Port 2000

How can the UDP server service multiple ports simultaneously?

UDP SERVER: SERVICING TWO PORTS

What problems does this code have?

```
int s1;
                                 /* socket descriptor 1 */
int s2;
                                 /* socket descriptor 2 */
/* 1) create socket s1 */
/* 2) create socket s2 */
/* 3) bind s1 to port 2000 */
/* 4) bind s2 to port 3000 */
while(1) {
      recvfrom(s1, buf, sizeof(buf), ...);
      /* process buf */
      recvfrom(s2, buf, sizeof(buf), ...);
      /* process buf */
```

SOCKET I/O: SELECT()

- maxfds: number of descriptors to be tested
 - descriptors (0, 1, ... maxfds-1) will be tested
- readfds: a set of fds we want to check if data is available
 - returns a set of fds ready to read
 - if input argument is *NULL*, not interested in that condition
- writefds: returns a set of fds ready to write
- **exceptfds**: returns a set of *fds* with exception conditions

SOCKET I/O: SELECT()

timeout

- If NULL, wait forever and return only when one of the descriptors is ready for I/O
- otherwise, wait up to a fixed amount of time specified by timeout
 - if we don't want to wait at all, create a timeout structure with timer value equal to 0

SOCKET I/O: SELECT() IN UDP

select allows synchronous I/O multiplexing

```
/* socket descriptors */
int s1, s2;
                  /* used by select() */
fd set readfds;
/* create and bind s1 and s2 */
while(1) {
        FD_ZERO(&readfds); /* initialize the fd set */
        FD_SET(s1, &readfds);    /* add s1 to the fd set */
        FD SET(s2, &readfds); /* add s2 to the fd set */
        if(select(s2+1, &readfds, 0, 0, 0) < 0) {
                perror("select");
                exit(1);
        if(FD ISSET(s1, &readfds)) {
                recvfrom(s1, buf, sizeof(buf), ...);
                /* process buf */
        /* do the same for s2 */
```

SOCKET I/O: SELECT() IN TCP

```
int fd, next=0;
                                           /* original socket */
                                   /* new socket descriptors */
int newfd[10];
while(1) {
      fd set readfds;
       FD ZERO (&readfds);
      FD SET (fd, &readfds);
       /* Now use FD SET to initialize other newfd's
          that have already been returned by accept() */
       select(maxfd+1, &readfds, 0, 0, 0);
       if(FD ISSET(fd, &readfds)) {
              newfd[next++] = accept(fd, ...);
       /* do the following for each descriptor newfd[n] */
       if(FD ISSET(newfd[n], &readfds)) {
              read(newfd[n], buf, sizeof(buf));
              /* process data */
```

EVENT-DRIVEN APPROACHES

- Use of asynchronous event notifications
 - Potentially faster and more flexible than select
- Provide notifications when events occur on file descriptors
- Designed to handle event loop in a fast, non-blocking way
- Libraries like libevent, libev, libuv, etc

BASIC PACKET BUILDING FOR A BUFFER

```
struct packet {
     u int32 t type;
     u int16 t length;
     u int16 t checksum;
     u int32 t address;
              char buf[1024];
struct packet *pkt;
pkt = (struct packet*) buf;
pkt->type = htonl(1);
pkt->length = htons(2);
pkt->checksum = htons(3);
pkt->address = htonl(4);
```

EXTENDING FUNCTIONALITY THROUGH PACKETS

- Possible to use TCP and UDP to get functionality of both protocols
- Also possible to add packet information and packet handling to UDP communication for greater reliability
 - e.g. Index checks on packets to verify order and delivery
- System needs and constraints determine how to approach problem
 - Don't reinvent TCP
 - But maybe a little more reliability is worth latency tradeoffs...

PAYLOAD CONSIDERATIONS

- What information needs to be in the packet?
- How large is the payload?
- What is the latency of serializing/deserializing the payload?
- How often do the server and clients need to know about this information?
- Is my payload secure and safe?

PACKET INFORMATION

- What information is in what packet should be architected with care
 - Cannot afford to send out the entire world state every frame
- Provide initial information about world schema to client upon connection
- Provide ongoing updates relative to this schema as the world state changes

DISCUSS

- Consider these client-server network scenarios. What should be in the packet? What needs to happen when the packet is received?
 - A player in an MMO trades with another player
 - A player in a battle royale equips a new weapon
 - A player in a go game places a stone
 - A player in an arena shooter uses a hit scan gun
 - A player in an arena shooter uses a ballistic gun

PACKET FORMAT

- XML and JSON are too verbose for the frequency data is being sent
- Text information is not tightly packed
- Ideally use a binary format
 - Low latency games may use a custom binary format rather than an existing library

PROBLEMS WITH MEMCPY

- Directly copying the struct data into the packet is very cheap
 - Works well on simple projects like what we're creating where only 4 or
 5 people will play it
- Major issues at a commercial level
 - Must ensure cross-platform/cross-compiler support for memory layout
 - Must handle endian-ness
 - Must handle pointers
 - Major security risk if struct data is simply trusted

READING AND WRITING PER-FIELD

- Create serialization library that reads and writes from the struct to the packet
 - Need to be able to read/write from every struct type
 - Need to be able to read/write into every packet type
- Can additionally perform better bitpacking here to ensure good packet properties

ADDITIONAL RESOURCES

- Gaffer on Games < https://www.gafferongames.com/>
 - Tons of in-depth articles on physics, networking, and networked physics