## CSEE 3827: Fundamentals of Computer Systems

Finite State Machine Design

## Recall: Sequential circuit



## Example sequential circuit (schematic)



## Reverse engineering a sequential circuit



## State machine

A state machine model of a system's behavior in terms of states and transitions between those states that are triggered by actions.

## State diagrams represent state machines

## one or more states, indicated by nodes


input value that triggers transition on edge

## Finite state machine (FSM)

A state machine that has a finite number of states

* Any finite state machine can be implemented with sequential logic
* All sequential circuits implement finite state machines


## Implementing a finite state machine


2. convert to truth table

| $S$ | in | $S+$ | out |
| :---: | :---: | :---: | :---: |
| 00 | 0 | 10 | 0 |
| 00 | 1 | 01 | 0 |
| 01 | 0 | 10 | 1 |
| 01 | 1 | 01 | 1 |
| 10 | 0 | 10 | 0 |
| 10 | 1 | 01 | 0 |

4. annotate table with flip-flop inputs for next state
5. wire circuit and flipflops together together

| $S$ | in | $S_{+}$ | out | T1 | T2 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | 0 | 10 | 0 | 1 | 0 |
| 00 | 1 | 01 | 0 | 0 | 1 |
| 01 | 0 | 10 | 1 | 1 | 1 |
| 01 | 1 | 01 | 1 | 0 | 0 |
| 10 | 0 | 10 | 0 | 0 | 0 |
| 10 | 1 | 01 | 0 | 1 | 1 |

## From State Machine to Sequential Circuit

- Specify behavior of state machine (including input and output values)
- Encode states
- figure out how many bits needed to store state (one FF per bit)
- assign state values to states
- Select FF type (i.e., D, T, JK, etc.)
- Compute combinational logic functions
- "next state" logic: $S_{+}=F(S$, inputs $)$
- "output" logic:
- Mealy machine: output $=\mathrm{G}(\mathrm{S}$, inputs $)$
- Moore machine: output $=\mathrm{H}(\mathrm{S})$


## Example State Machine

- 1 input, 1 output
- Let $X=\#$ of 1 's input so far, $Y=\#$ of 0 's input so far.
- Output 1 whenever $X-Y=0(\bmod 3)$


State label is current value of $X-Y$ mod 3


Binary Labeling

## Design with D Flip-Flops

- D FF's are easy: we input the value to the FF that we want it set to

| $A_{\text {cur }}$ | $B_{\text {cur }}$ | In | $A_{\text {next }}$ | $B_{\text {next }}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ |



| $A_{\text {cur }}$ | $B_{\text {cur }}$ | In | $A_{\text {next }}$ | $D_{A}$ | $B_{\text {next }}$ | $D_{B}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ | $X$ | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ | $X$ | $X$ |

New columns indicate what values feed into D FF (mimic "next" values for that FF)

## Design with D Flip-Flops Cont'd

- Build K-Maps, get equations for output and FF input vals (in terms of inputs and previous F vals)

| $A_{\text {cur }}$ | $B_{\text {cur }}$ | In | $A_{\text {next }}$ | $D_{A}$ | $B_{\text {next }}$ | $D_{B}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ | $X$ | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ | $X$ | $X$ |

$$
\begin{array}{r}
\mathrm{D}_{\mathrm{A}}=\overline{\mathrm{A}} \bar{B} \overline{\mathrm{I}}+\mathrm{BI} \\
\mathrm{D}_{\mathrm{B}}=\mathrm{A} \overline{\mathrm{~B}}+\overline{\mathrm{A}} \overline{\mathrm{~B}} \mathrm{I} \\
\text { Out }=\mathrm{Al}+\mathrm{B} \overline{\mathrm{I}}
\end{array}
$$

Design with D FF's cont'd


## Design with T Flip Flops

- Value fed into T FF is 0 if FF should maintain value, 1 if it should flop

| $A_{\text {cur }}$ | $B_{\text {cur }}$ | In | $A_{\text {next }}$ | $B_{\text {next }}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ |


| $A_{\text {cur }}$ | $\mathrm{B}_{\text {cur }}$ | $\ln$ | $\mathrm{A}_{\text {next }}$ | $\mathrm{T}_{\mathrm{A}}$ | $\mathrm{B}_{\text {next }}$ | $\mathrm{T}_{\mathrm{B}}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | 0 |  | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 |  | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 |  | 1 |
| 0 | 1 | 1 | 1 | 1 | 0 |  | 0 |
| 1 | 0 | 0 | 0 | 1 | 1 |  | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 |  | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ |  | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ |  | $X$ |

$\mathrm{T}_{\mathrm{A}}$ :


$$
\mathrm{T}_{\mathrm{A}}=\mathrm{A}+\mathrm{BI}+\overline{\mathrm{B}} \overline{\mathrm{I}}
$$

## Design with T Flip Flops

- Value fed into T FF is 0 if FF should maintain value, 1 if it should flop

| $A_{\text {cur }}$ | $B_{\text {cur }}$ | In | $A_{\text {next }}$ | $B_{\text {next }}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ |



| $A_{\text {cur }}$ | $\mathrm{B}_{\text {cur }}$ | In | $\mathrm{A}_{\text {next }}$ | $\mathrm{T}_{\mathrm{A}}$ | $\mathrm{B}_{\text {next }}$ | $\mathrm{T}_{\mathrm{B}}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | X | X | X | X | X |
| 1 | 1 | 1 | X | X | X | X | X |

In
TB:

$$
\text { A cur }^{}
$$

$$
T_{B}=A \bar{B} \bar{I}+\bar{A} \bar{B} I
$$

## Design with JK Flip-Flops

- Note: to change $A_{\text {cur }}$ to the correct $A_{n e x t}$ value, two possible input pairs can be fed into the J,K inputs of a JK Flip-Flop

| $\mathbf{J}$ | $\mathbf{K}$ | $\mathbf{Q ( t + 1 )}$ |
| :---: | :---: | :---: |
| 0 | 0 | $Q(t)$ |
| 0 | 1 | 0 |
| 1 | 0 | 1 |
| 1 | 1 | $\overline{Q(t)}$ |


| Acur | $\mathbf{A n e x t}^{\text {J,K }}$ |  |
| :---: | :---: | :---: |
| 0 | 0 | 0,0 or $0,1(0, X)$ |
| 0 | 1 | 1,0 or $1,1(1, X)$ |
| 1 | 0 | 0,1 or $1,1(X, 1)$ |
| 1 | 1 | 0,0 or $1,0(X, 0)$ |

Design with JK Flip-Flop

| $\mathbf{A}_{\text {cur }}$ | $\mathbf{A}_{\text {next }}$ | $\mathbf{J , K}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0,0 or $0,1(0, X)$ |
| 0 | 1 | 1,0 or $1,1(1, X)$ |
| 1 | 0 | 0,1 or $1,1(X, 1)$ |
| 1 | 1 | 0,0 or $1,0(X, 0)$ |


| $A_{\text {cur }}$ | $B_{\text {cur }}$ | $\ln$ | $A_{\text {next }}$ | $B_{\text {next }}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | X | X | X |
| 1 | 1 | 1 | X | X | X |


| $A_{\text {cur }}$ | $\mathrm{B}_{\text {cur }}$ | In | $A_{\text {next }}$ | $J_{A}$ | $K_{A}$ | $B_{\text {next }}$ | $J_{B}$ | $K_{B}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | $X$ | 0 |  |  | 0 |
| 0 | 0 | 1 | 0 | 0 | $X$ | 1 |  |  | 0 |
| 0 | 1 | 0 | 0 | 0 | $X$ | 0 |  |  | 1 |
| 0 | 1 | 1 | 1 | 1 | $X$ | 0 |  |  | 0 |
| 1 | 0 | 0 | 0 | $X$ | 1 | 1 |  |  | 0 |
| 1 | 0 | 1 | 0 | $X$ | 1 | 0 |  |  | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ | $X$ |  |  | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ | $X$ |  |  | $X$ |

Design with JK Flip-Flop

| Acur | A $_{\text {next }}$ | $\mathbf{J , K}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0,0 or $0,1(0, \mathrm{X})$ |
| 0 | 1 | 1,0 or $1,1(1, \mathrm{X})$ |
| 1 | 0 | 0,1 or $1,1(X, 1)$ |
| 1 | 1 | 0,0 or $1,0(X, 0)$ |


| $A_{\text {cur }}$ | $B_{\text {cur }}$ | In | $A_{\text {next }}$ | $B_{\text {next }}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ |



| $A_{\text {cur }}$ | $\mathrm{B}_{\text {cur }}$ | In | $A_{\text {next }}$ | $J_{A}$ | $K_{A}$ | $\mathrm{~B}_{\text {next }}$ | $J_{B}$ | $K_{B}$ | Out |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 1 | $X$ | 0 | 0 | $X$ | 0 |
| 0 | 0 | 1 | 0 | 0 | $X$ | 1 | 1 | $X$ | 0 |
| 0 | 1 | 0 | 0 | 0 | $X$ | 0 | $X$ | 1 | 1 |
| 0 | 1 | 1 | 1 | 1 | $X$ | 0 | $X$ | 1 | 0 |
| 1 | 0 | 0 | 0 | $X$ | 1 | 1 | 1 | $X$ | 0 |
| 1 | 0 | 1 | 0 | $X$ | 1 | 0 | 0 | $X$ | 1 |
| 1 | 1 | 0 | $X$ | $X$ | $X$ | $X$ | $X$ | $X$ | $X$ |
| 1 | 1 | 1 | $X$ | $X$ | $X$ | $X$ | $X$ | $X$ | $X$ |

Design with JK Flip-Flop


Design with JK FF's cont'd

$$
J_{A}=B I+\bar{B} \bar{T} \quad K_{A}=1 \quad J_{B}=\bar{A} I+A \bar{T} \quad K_{B}=1
$$



In class exercise: design a 3-bit counter

## Moore machine


a circuit in which the output depends only on the current state

## Mealy machine


a circuit in which the outputs depend on the inputs as well as the current state

## A Mealy or Moore circuit?


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## An example Moore circuit

5-16

(a)

| Present <br> state | Inputs |  |  |  |
| :---: | :---: | :---: | :---: | :---: | | Inext |
| :---: | :---: | :---: | :---: |
| state | Output

(b) State table

## FSM timing characteristics

## MEALY



## Advanced FSM design and implementation

Unused states: extra state encodings (e.g., using 3 FFs to represent 6 states leaves 2 unused states) can be treated as "don't care" values and used to simplify the combinational logic

State minimization: two states are equivalent if they transition to the same or equivalent states on the same inputs (while producing the same outputs in the case of a Mealy machine)

