

# Data Mining Techniques

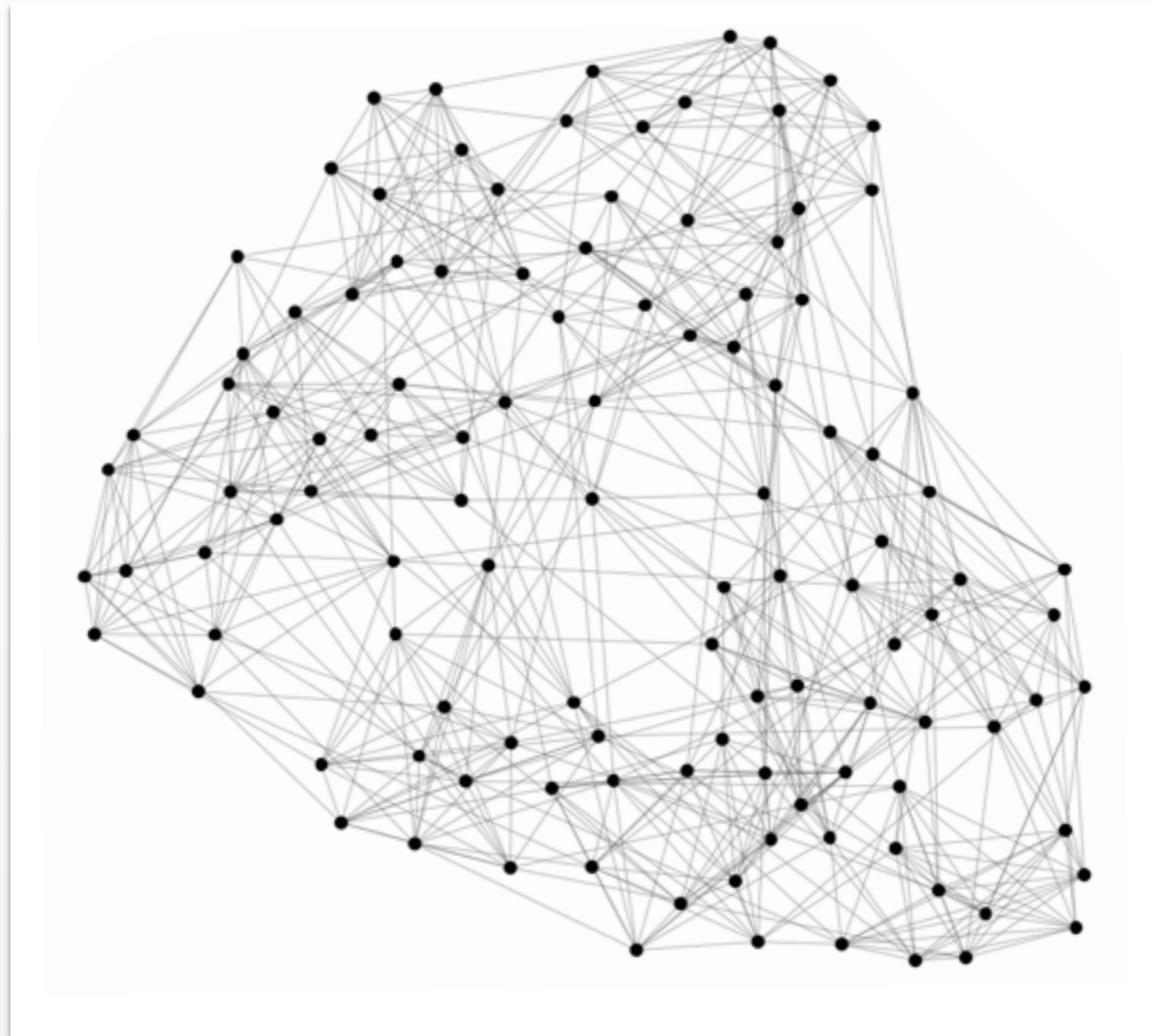
CS 6220 - Section 3 - Fall 2016

## Lecture 19: Social Networks

Jan-Willem van de Meent  
(*credit*: Leskovec et al Chapter 10,  
Aggarwal Chapter 19)



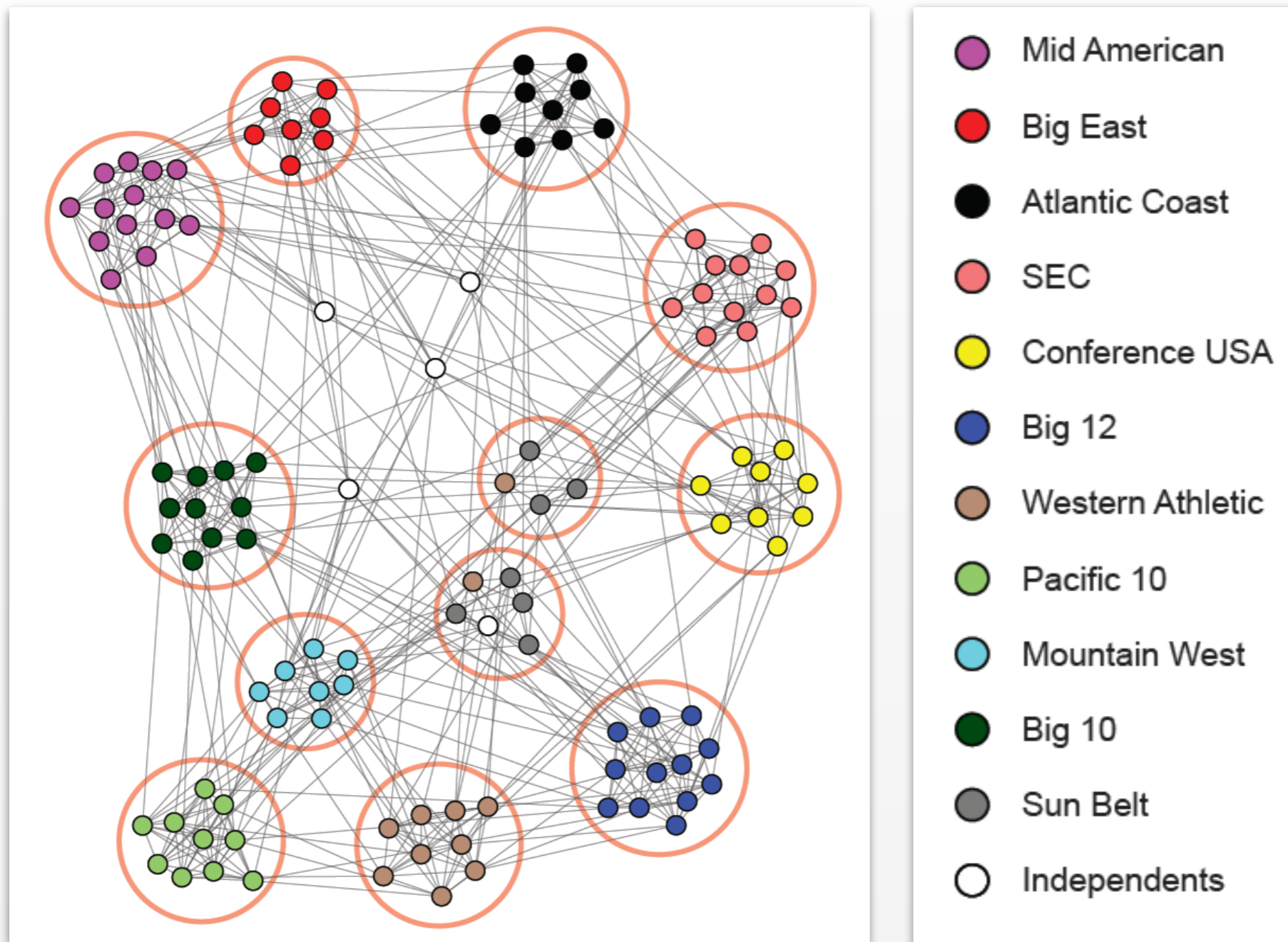
# Community Detection



*Problem:* Can we identify groups of densely connected nodes?

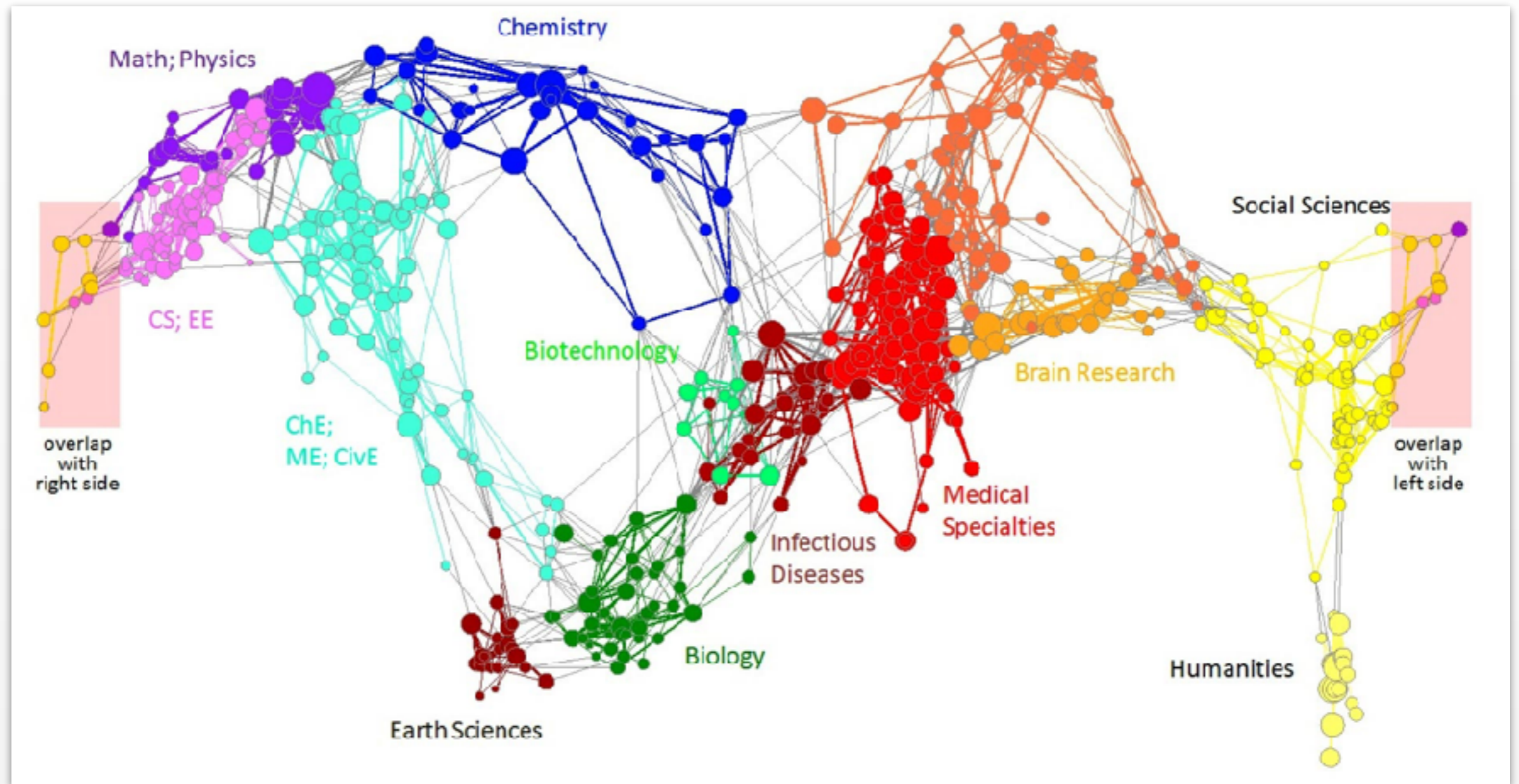
(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

# Communities: Football Conferences



*Nodes:* Football Teams, *Edges:* Matches,  
*Communities:* Conferences

# Communities: Academic Citations

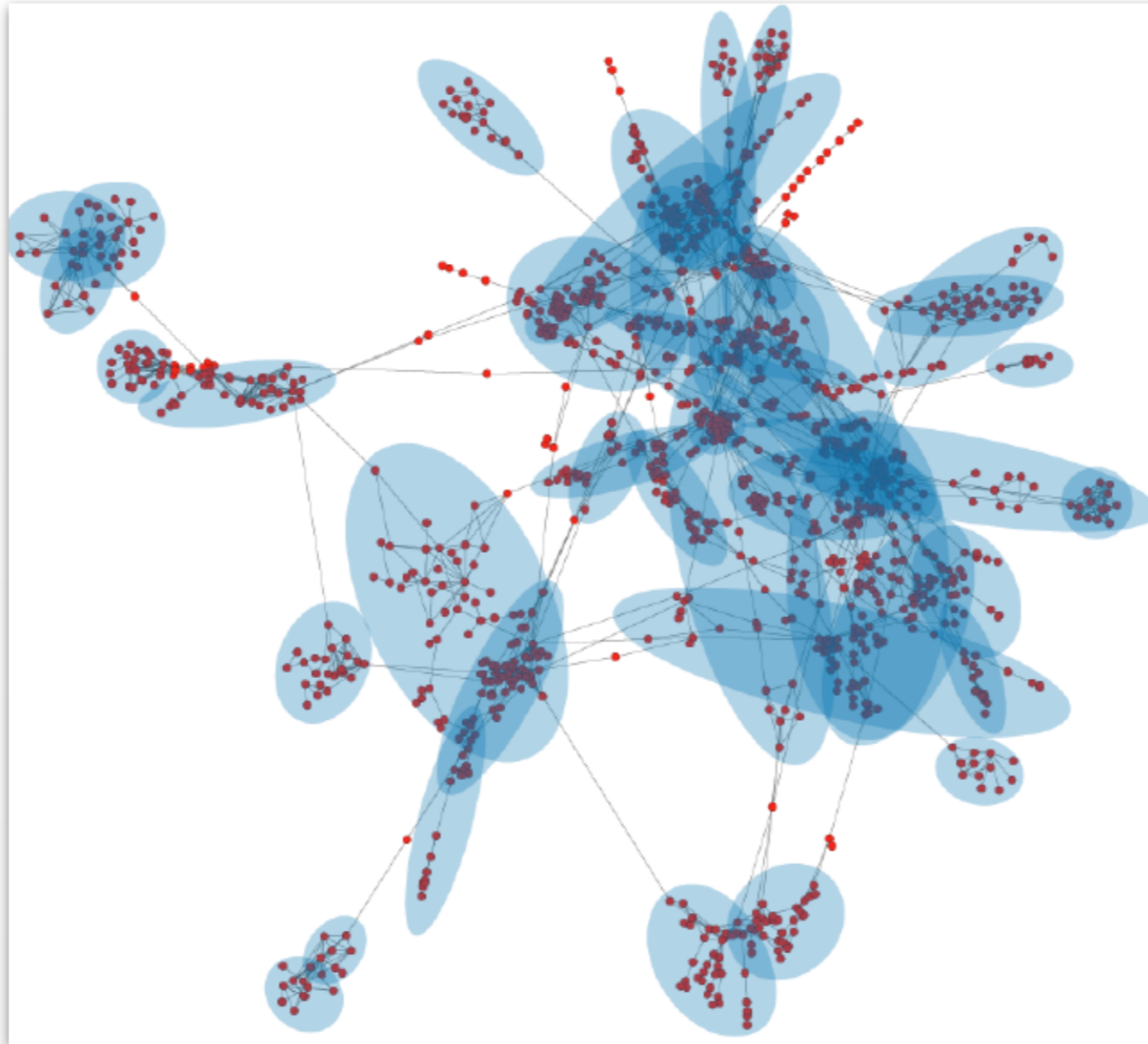


Source: Citation networks and Maps of science [Börner et al., 2012]

*Nodes: Journals, Edges: Citations,  
Communities: Academic Disciplines*

(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

# Communities: Protein-Protein Interactions



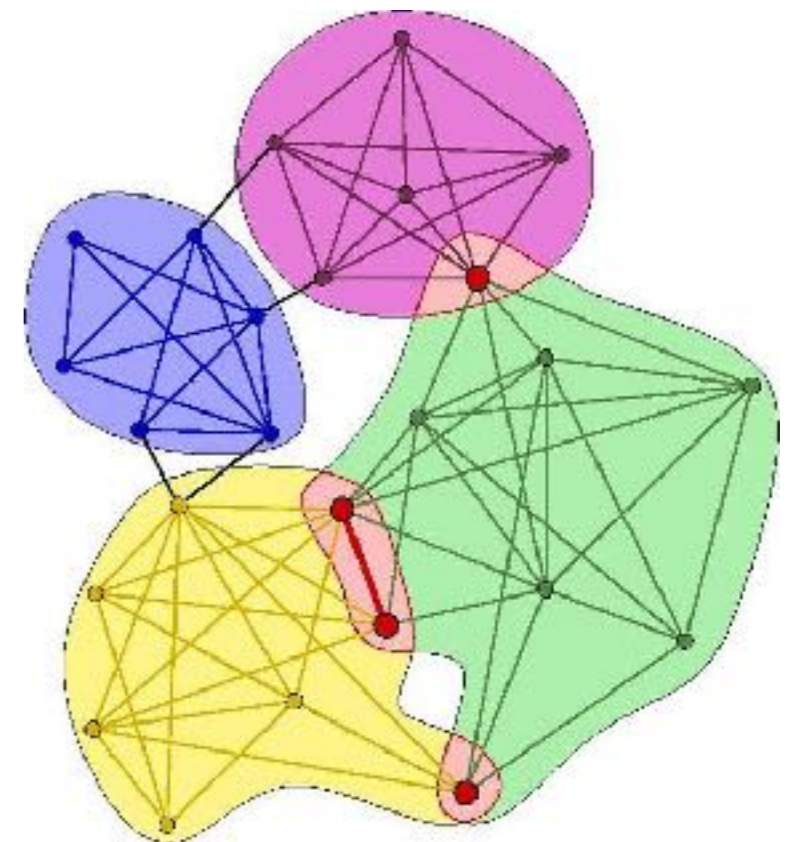
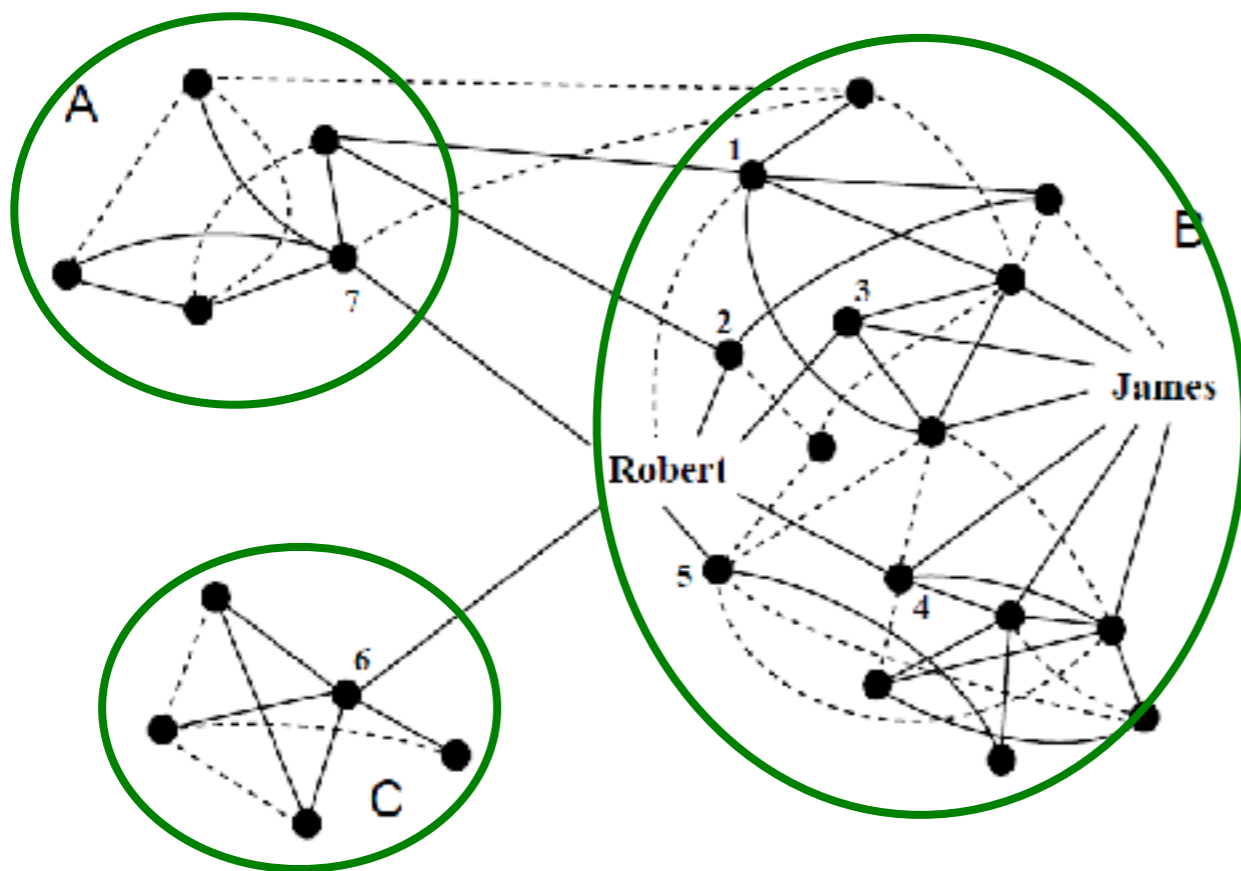
*Nodes:* Proteins, *Edges:* Physical interactions,  
*Communities:* Functional Modules

(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

# Community Detection

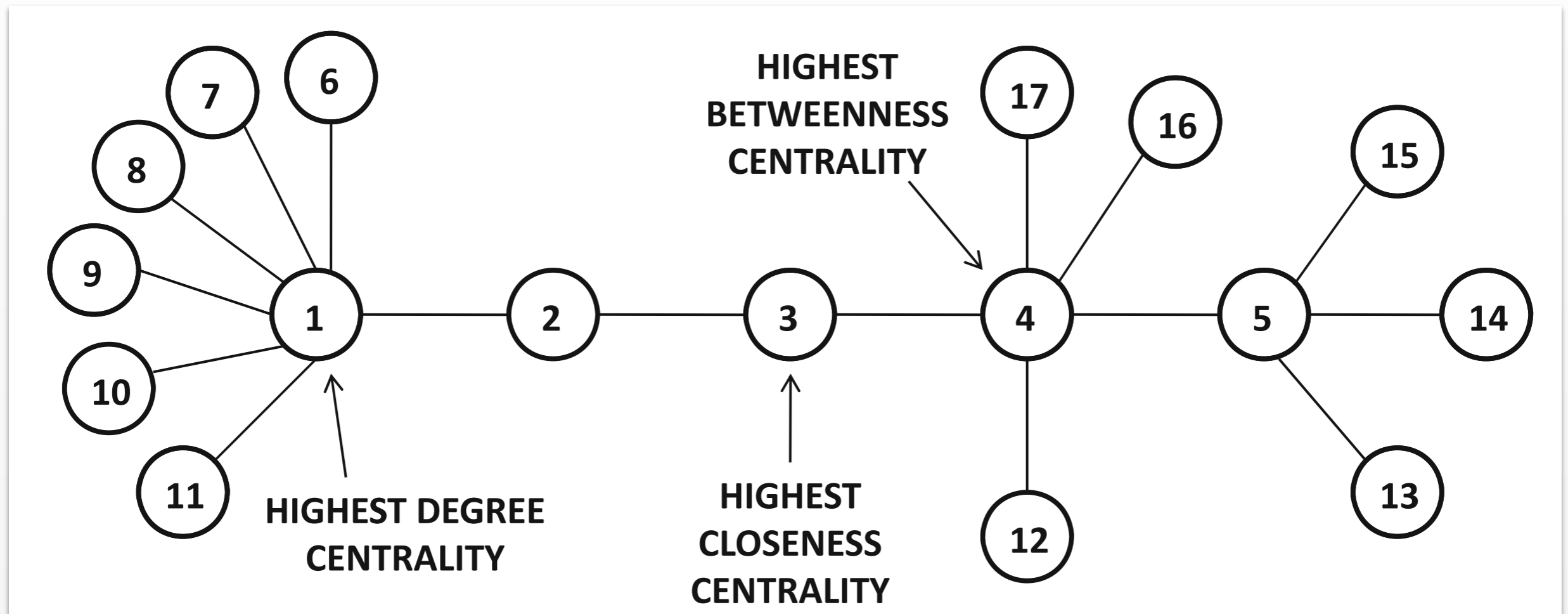
Graph Partitioning

Overlapping Communities



We will work with **undirected** (unweighted) networks

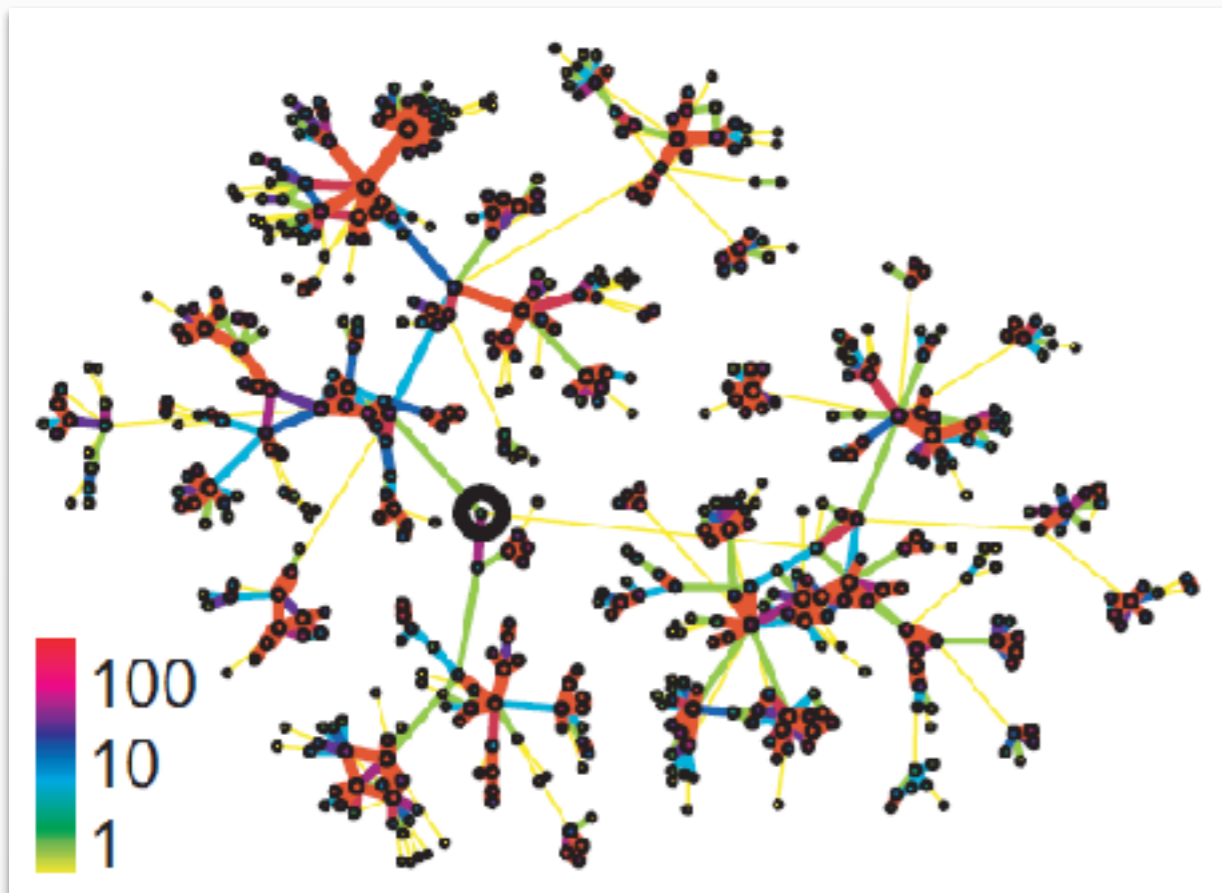
# Centrality Measures



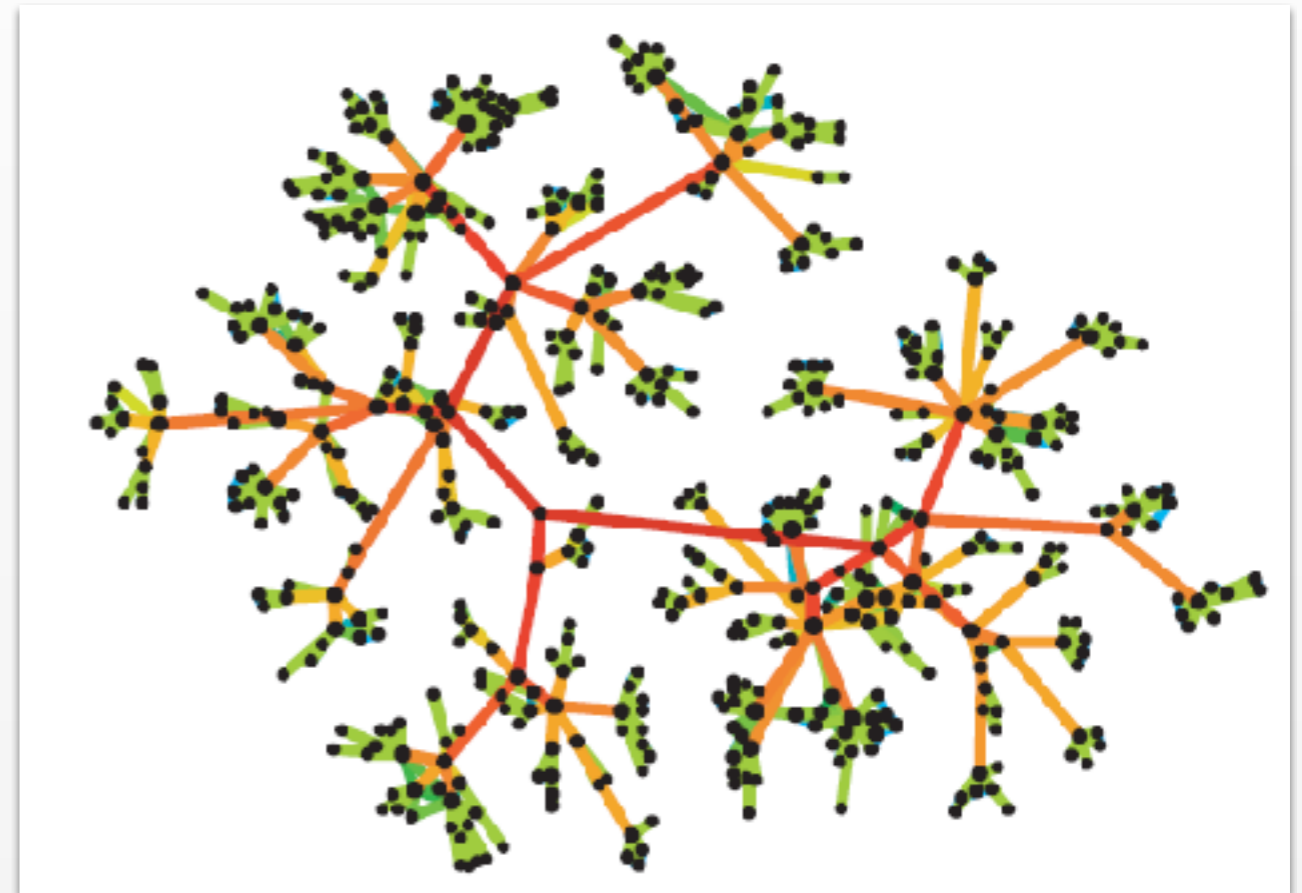
- *Betweenness*: Number of shortest paths
- *Closeness*: Average distance to other nodes
- *Degree*: Number of connections to other nodes

# Betweenness

Edge Strength (call volume)



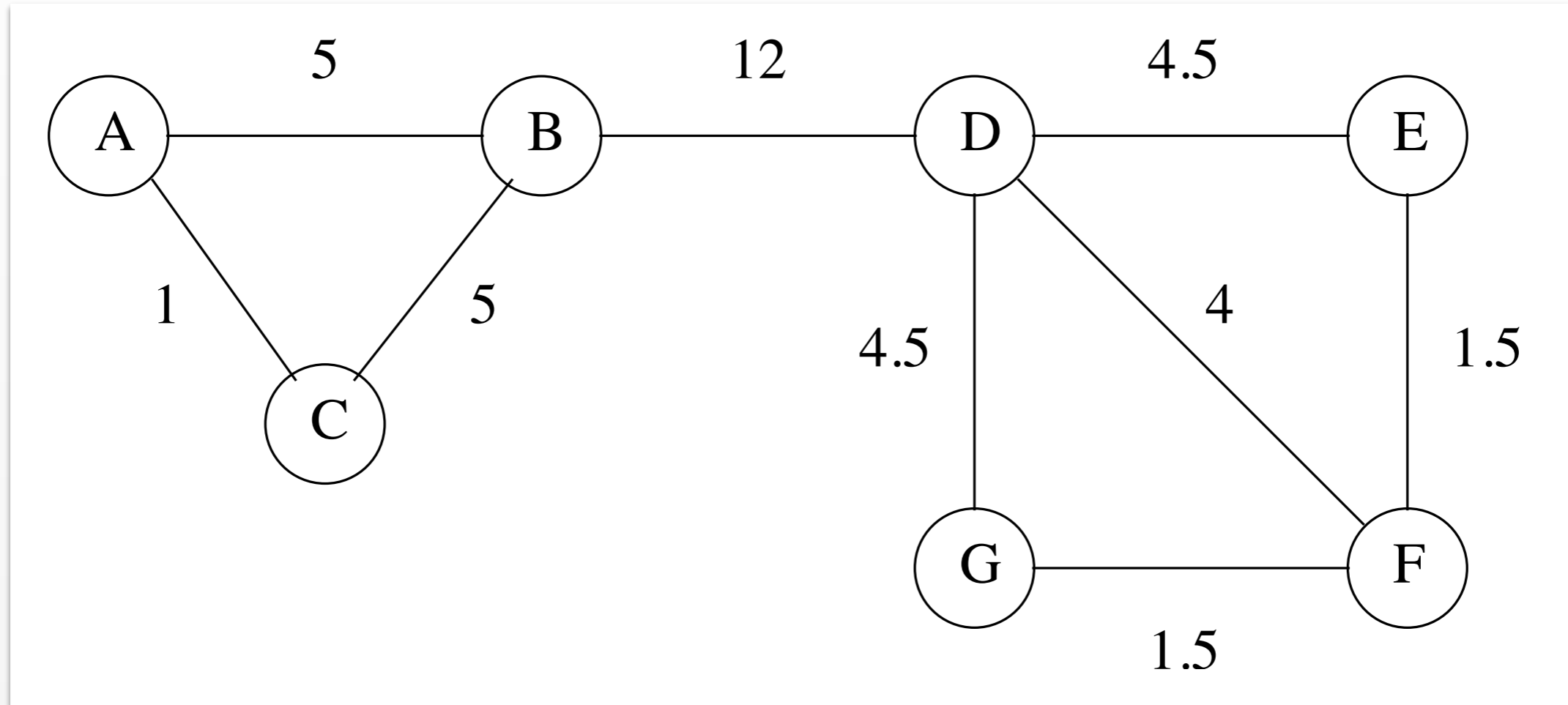
Edge Betweenness



- ***Betweenness***: Number of shortest paths passing through a node or edge



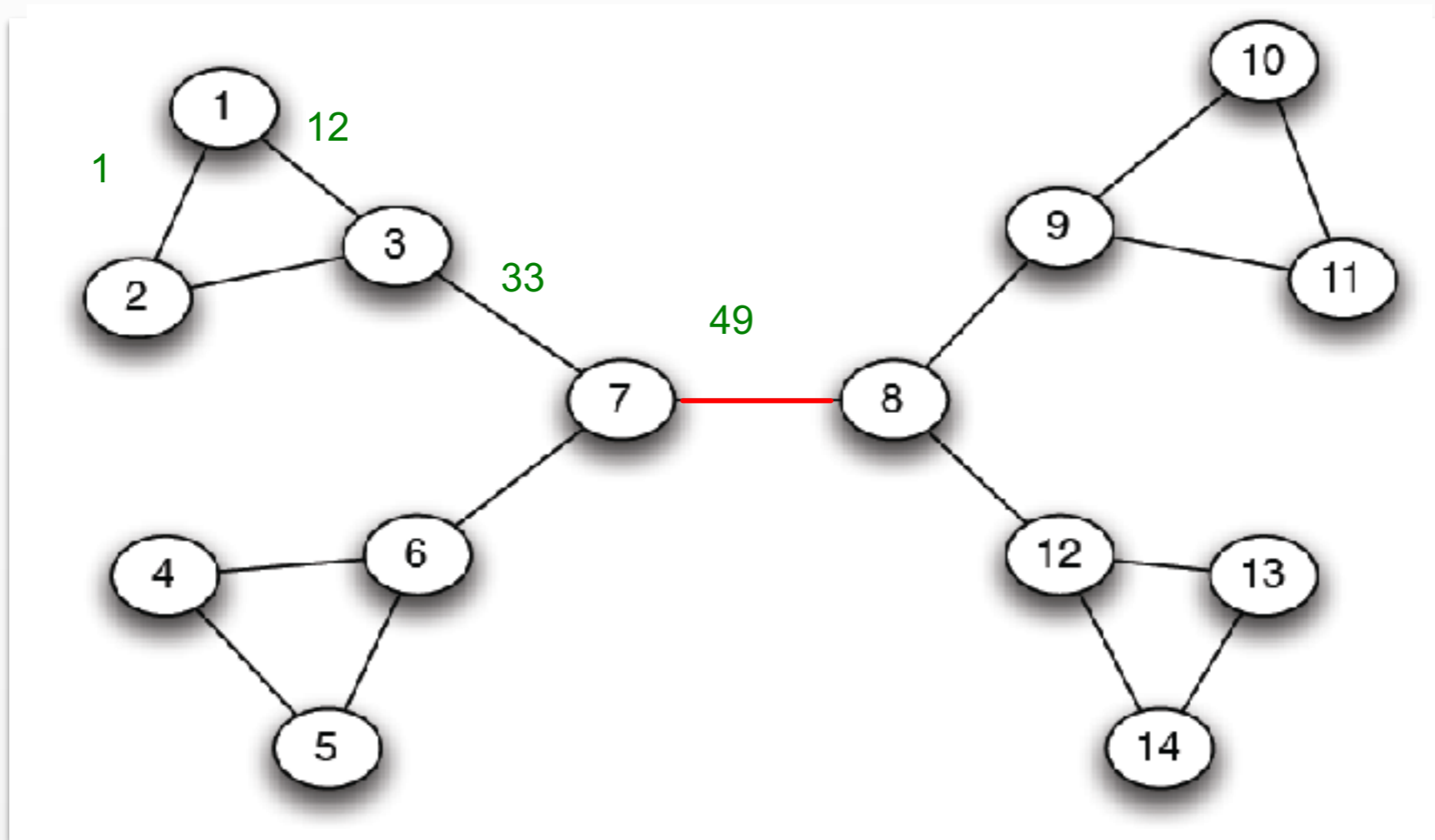
# Edge Betweenness



- Count number of shortest paths passing through each edge  
(*can be done with weighted edges*)
- If there are multiple paths of equal length, then split counts

# Girvan-Newman Algorithm

*(hierarchical divisive clustering according to betweenness)*

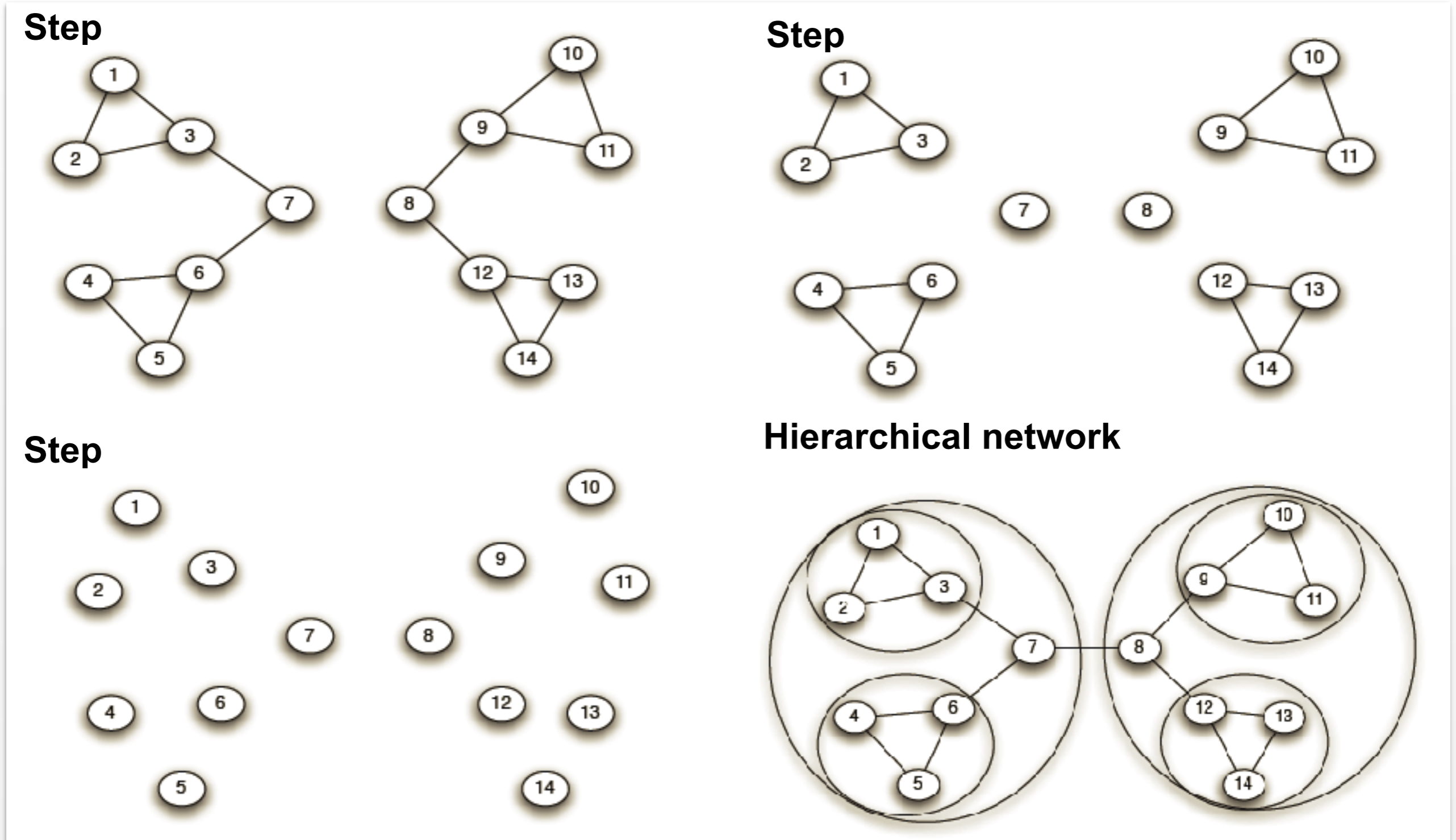


*Repeat until  $k$  clusters found*

1. Calculate betweenness
2. Remove edge(s) with highest betweenness

# Girvan-Newman Algorithm

*(hierarchical divisive clustering according to betweenness)*





# Girvan-Newman

*Two problems*

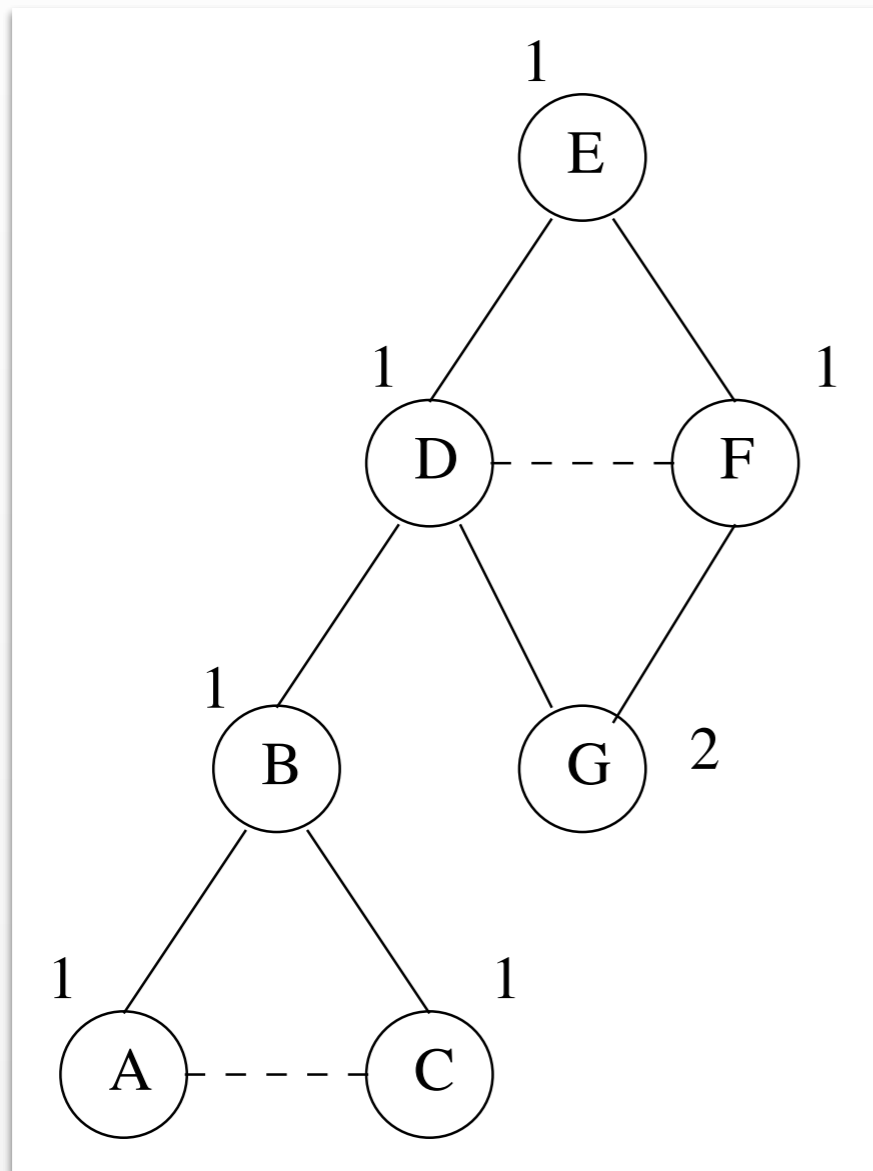
1. How can we compute the betweenness for all edges?
2. How can we choose the number of components  $k$ ?

# Calculating Betweenness

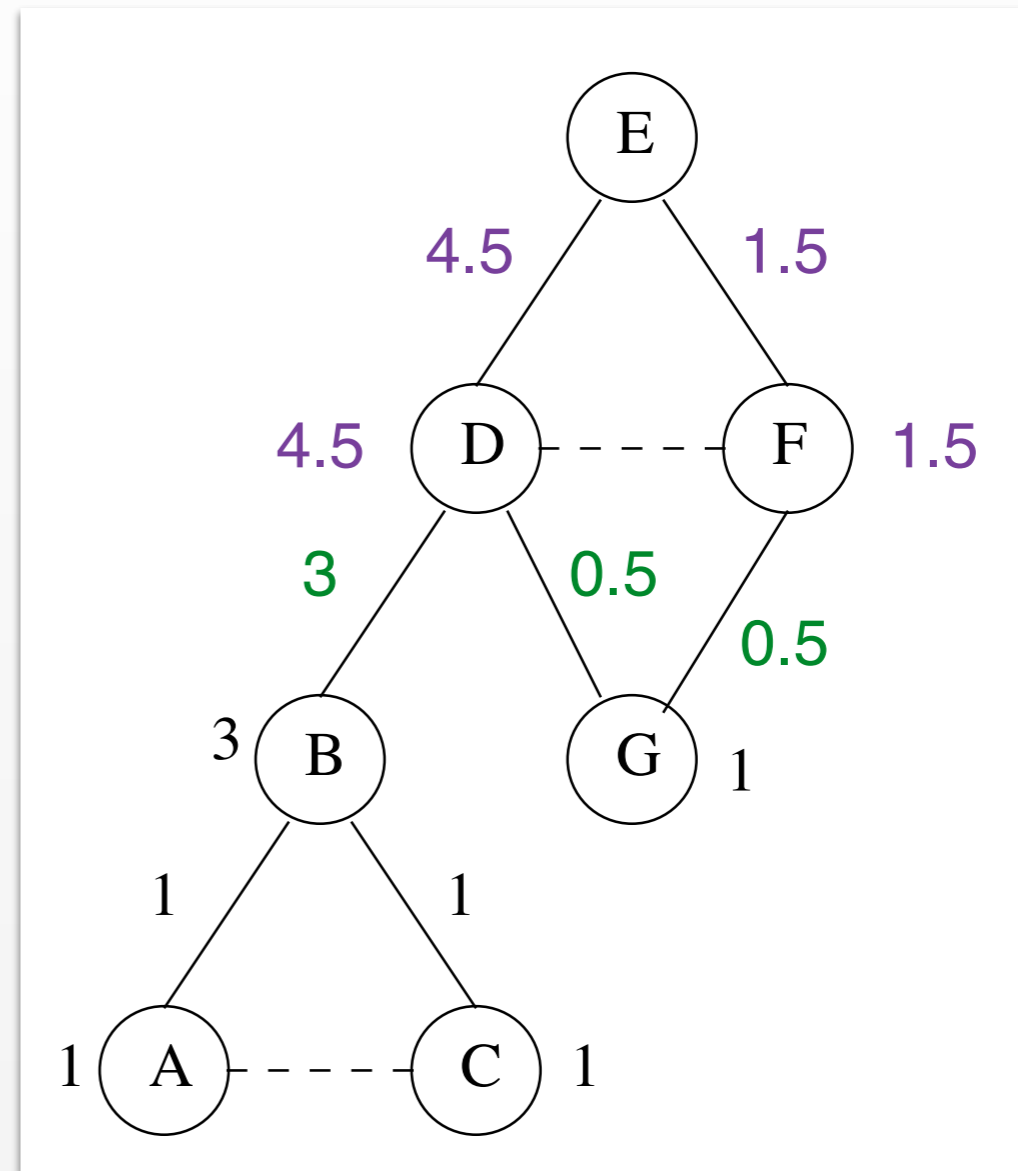
*How can we count all shortest paths?*

- Loop over nodes in graph
  - Perform breadth-first search to find shortest paths to other nodes
  - Increment counts for edges traversed by shortest paths
- Divide final betweenness by 2  
*(since all paths counted twice)*

# Counting Shortest Paths



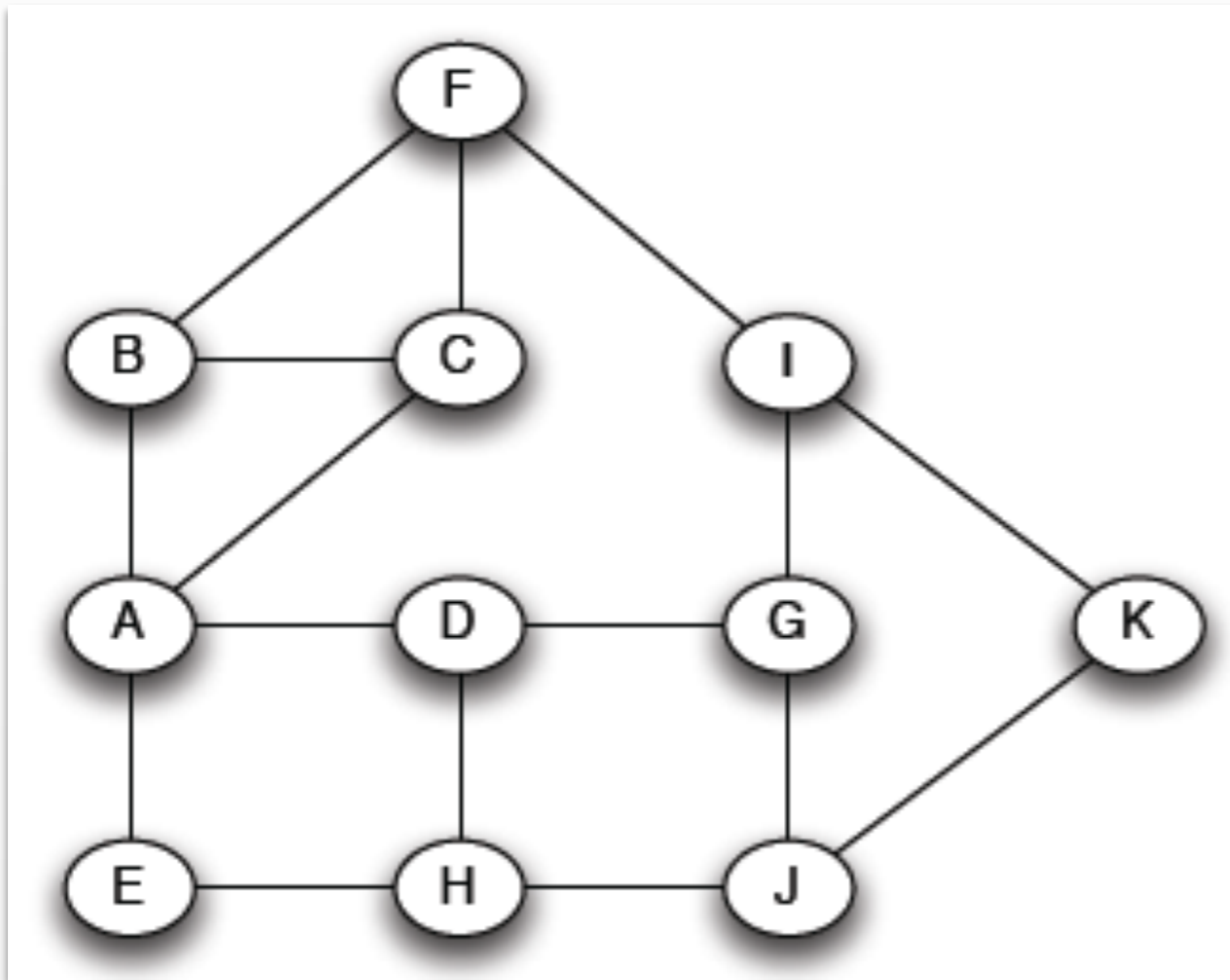
Count number of shortest paths from (E) to each node



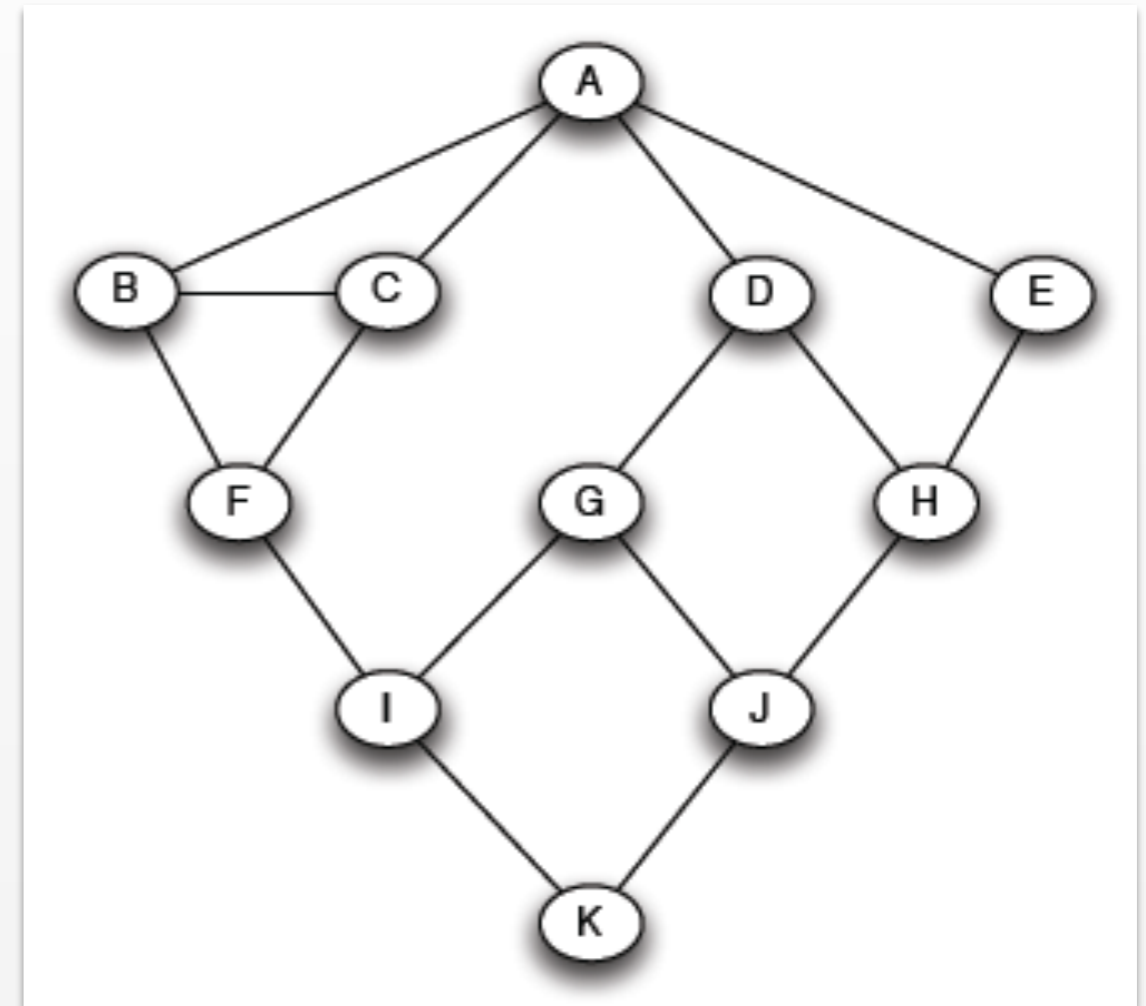
Accumulate credit upwards, dividing across shortest paths

# Counting Paths: Larger Example

Original Graph

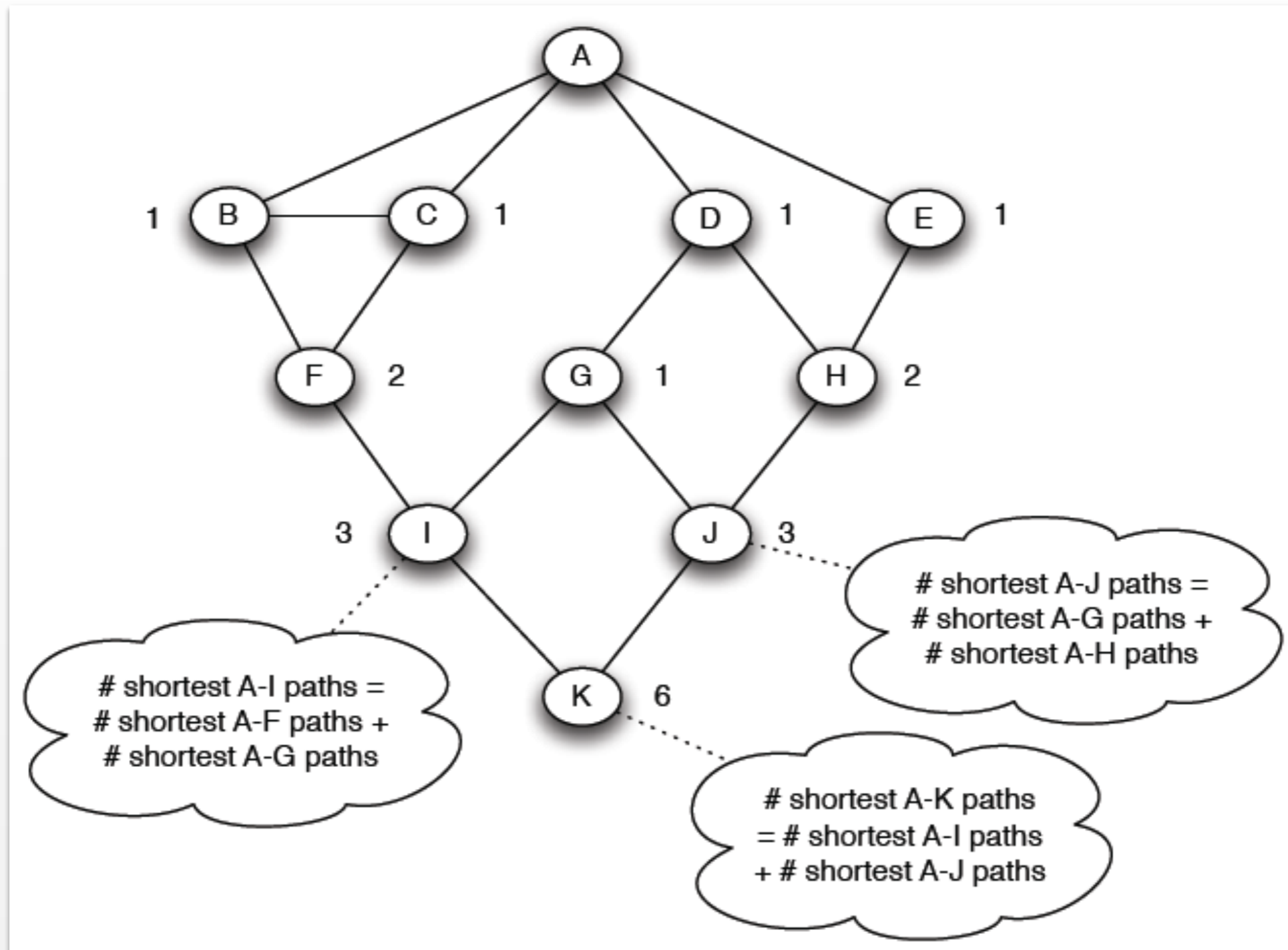


Breadth-first Ordering from A



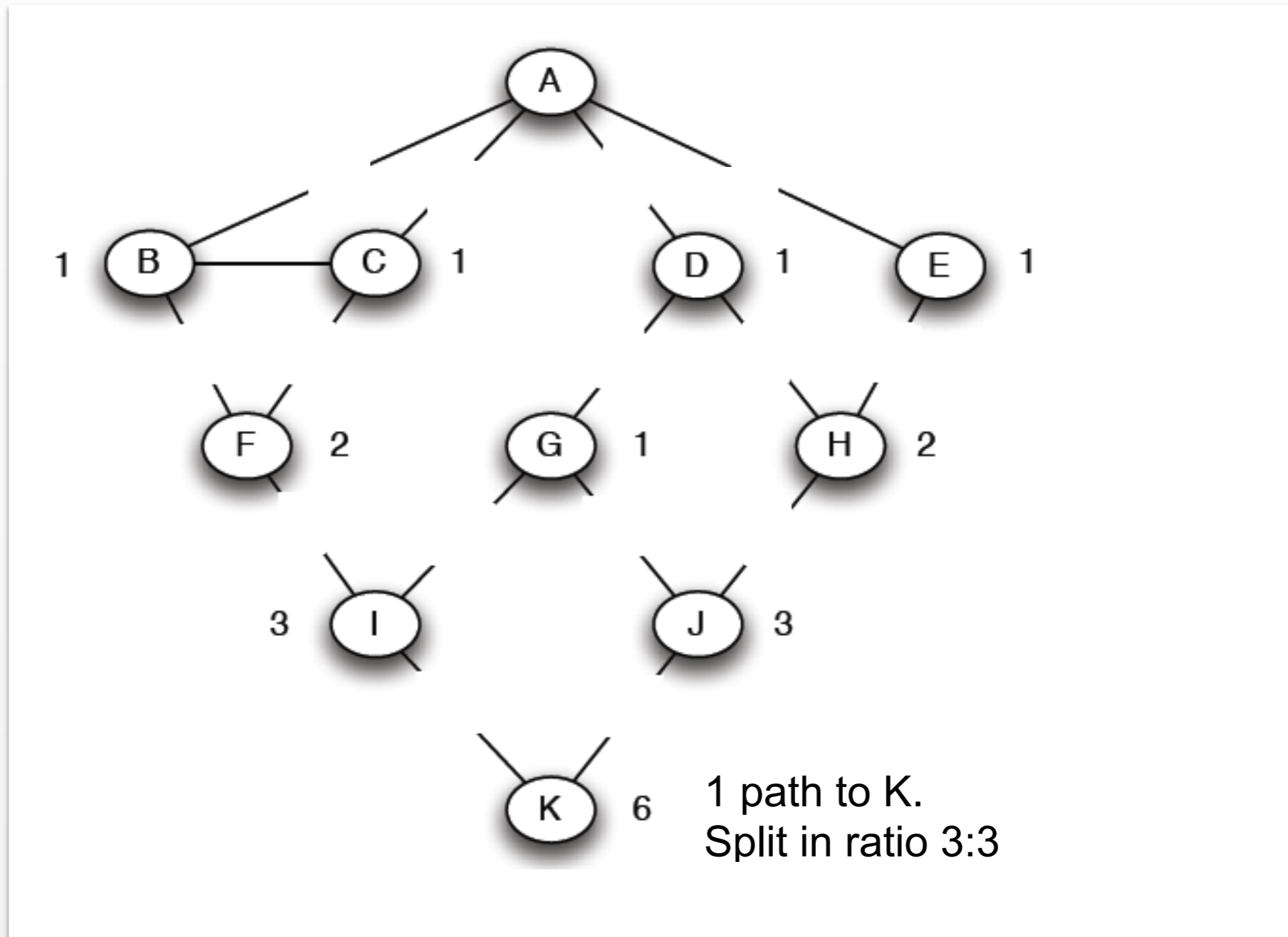


# Counting Paths: Larger Example



Step 1. Count number of shortest paths from to each node

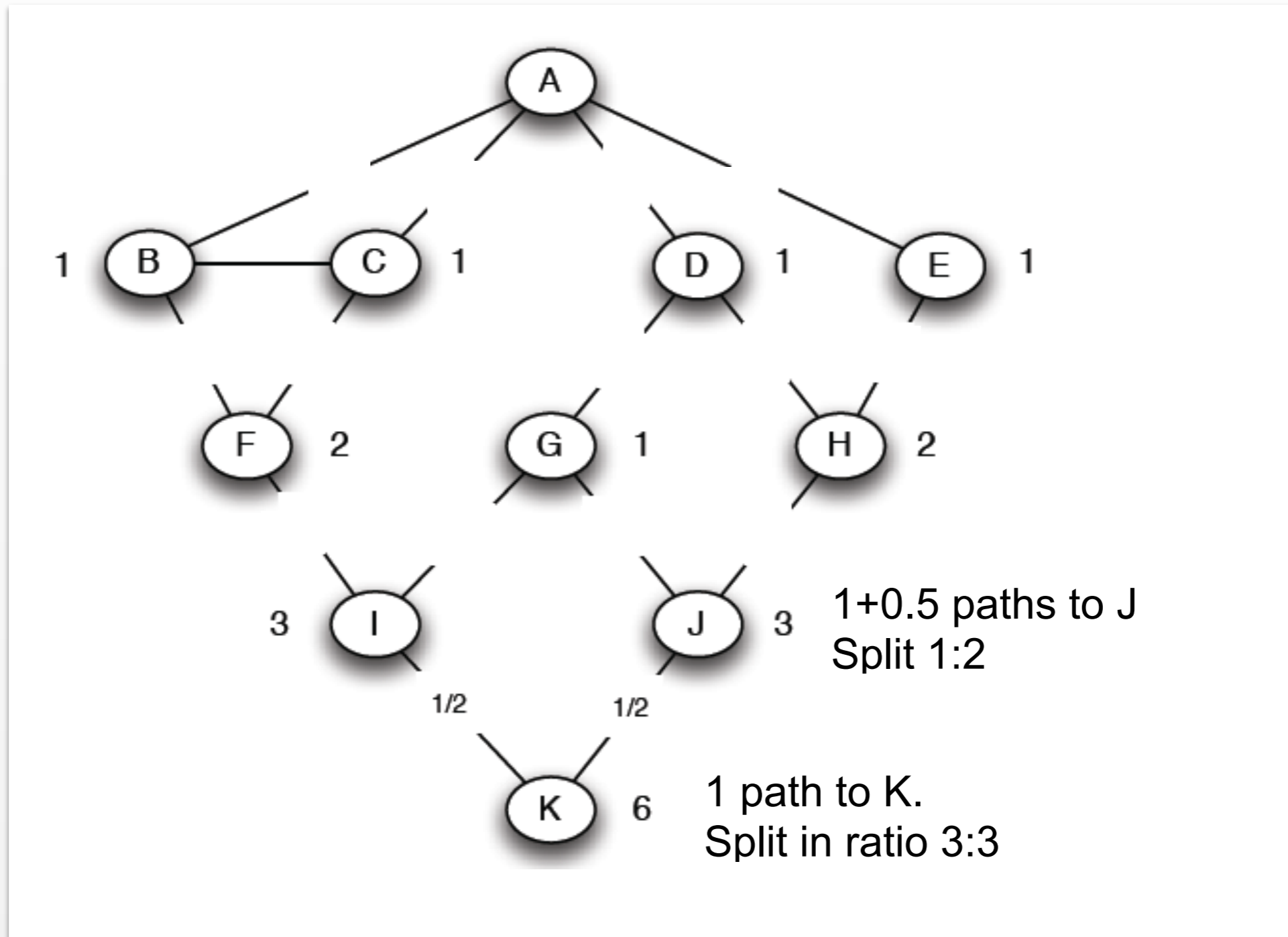
# Counting Paths: Larger Example



Step 2. Propagate credit upwards, splitting according to number of paths to parents

(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

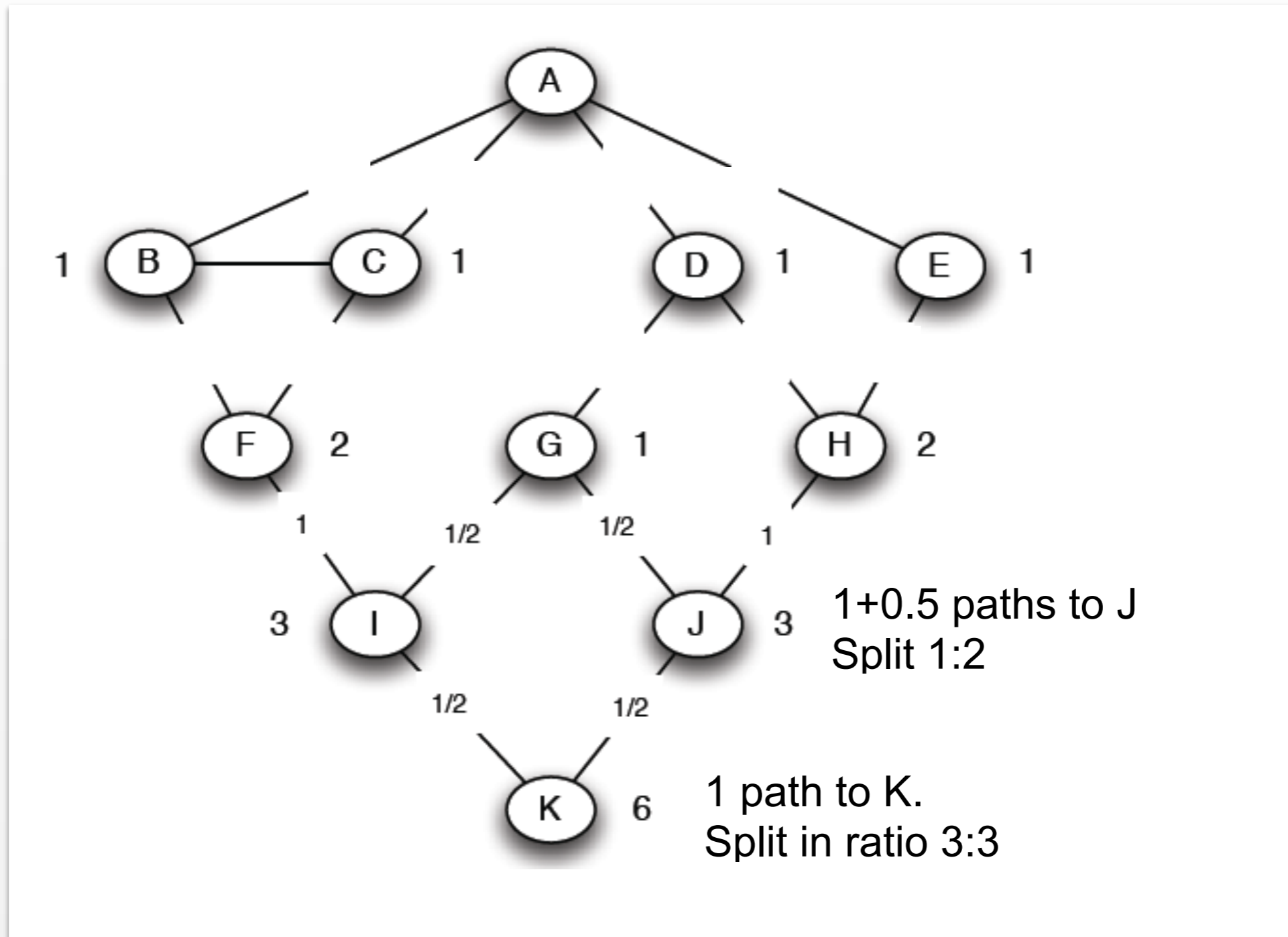
# Counting Paths: Larger Example



Step 2. Propagate credit upwards, splitting according to number of paths to parents

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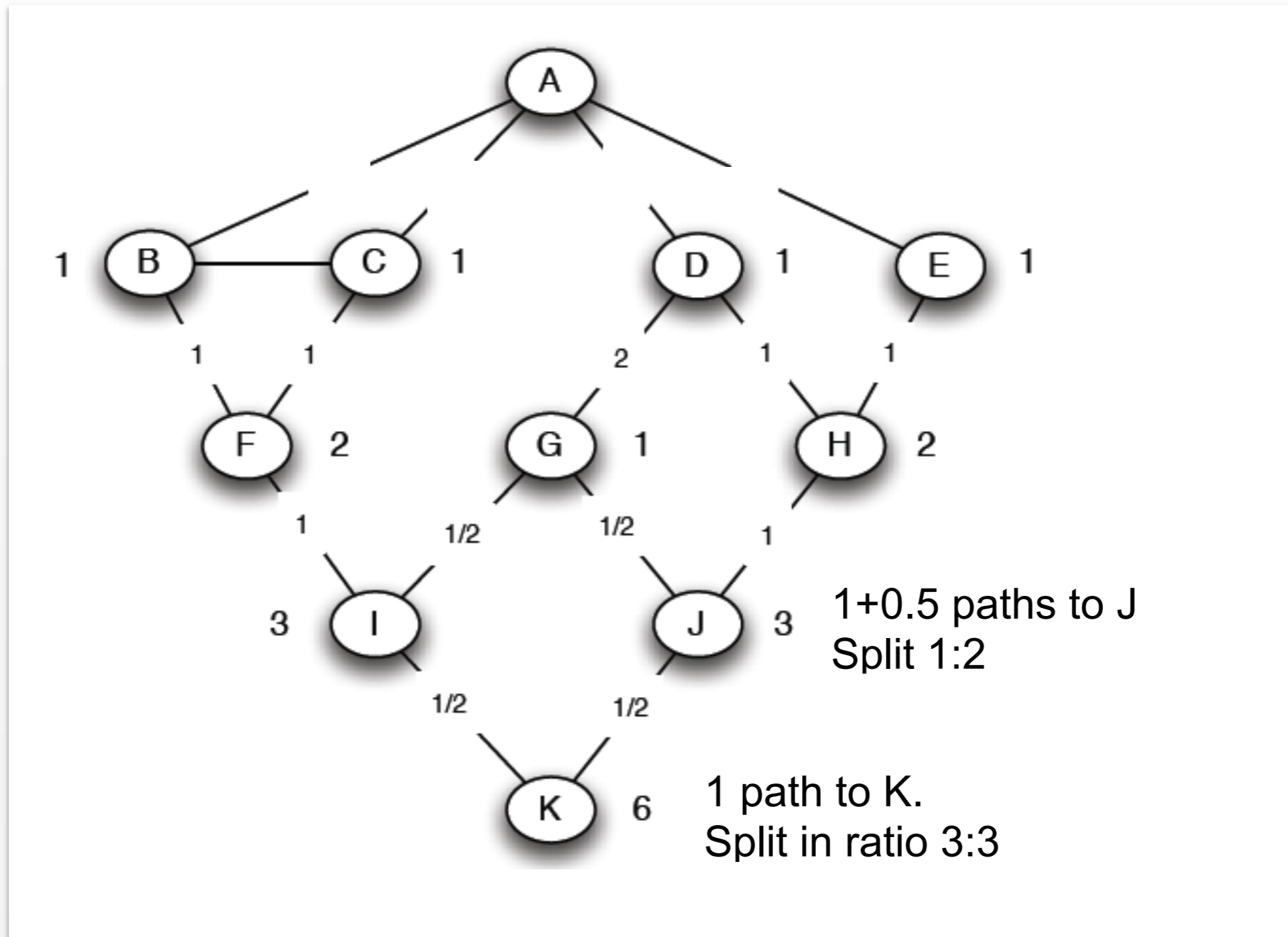
# Counting Paths: Larger Example



Step 2. Propagate credit upwards, splitting according to number of paths to parents

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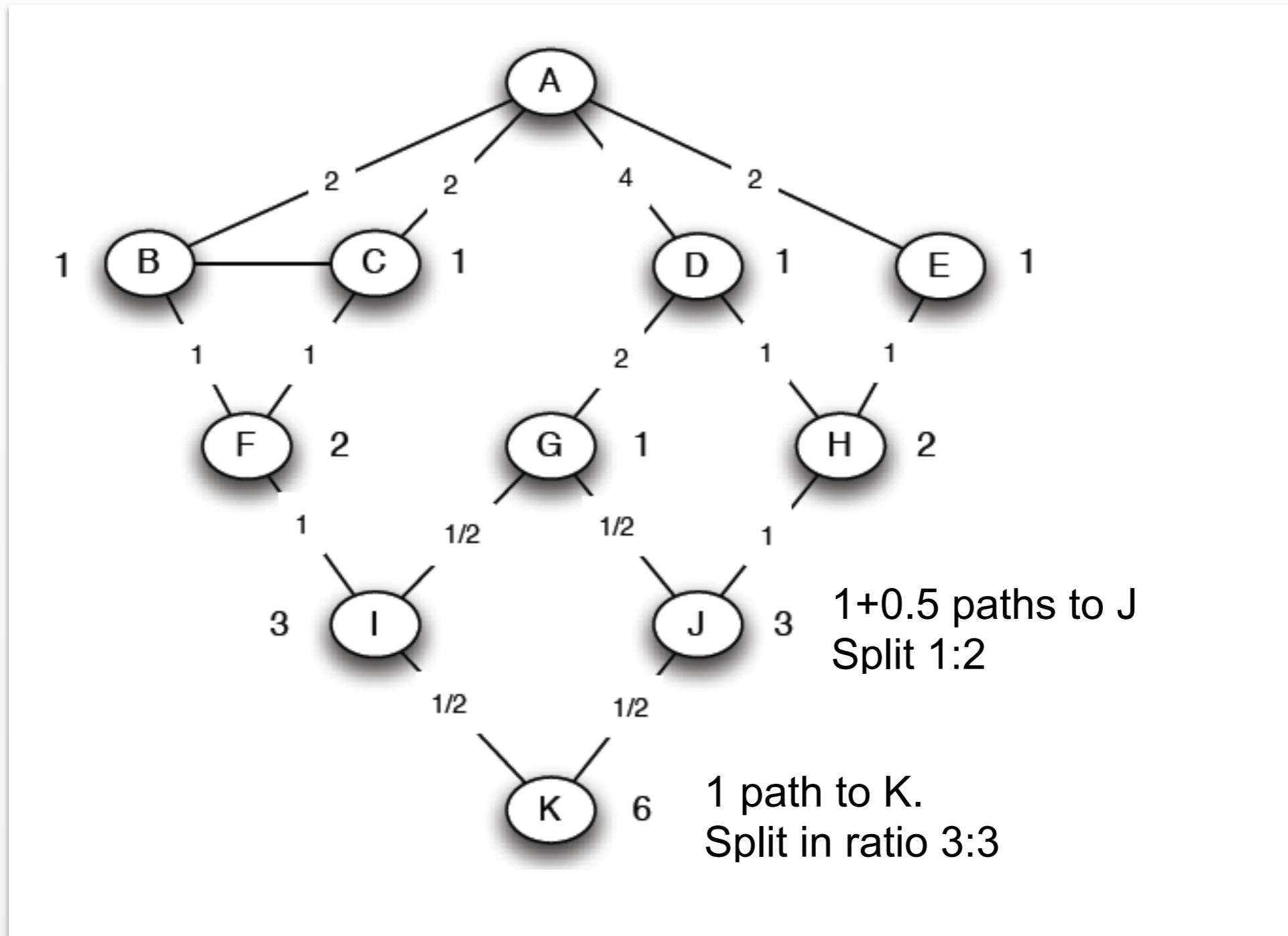
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# Counting Paths: Larger Example

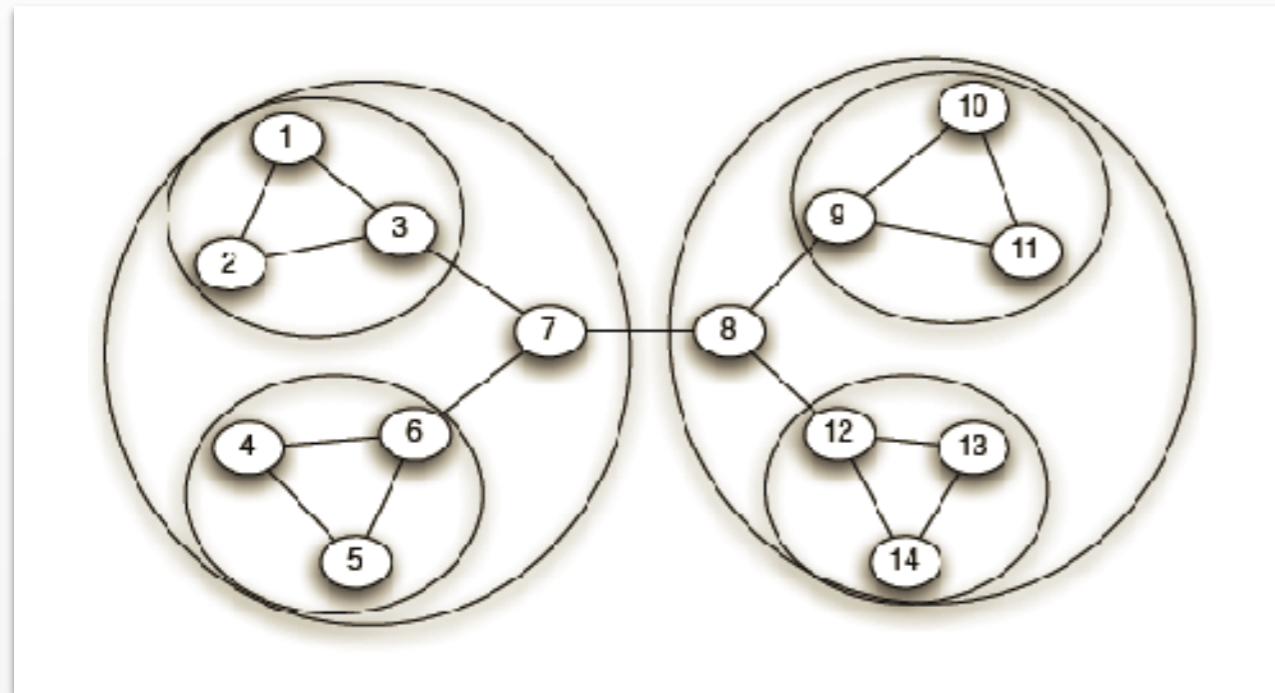


Step 2. Propagate credit upwards, splitting according to number of paths to parents

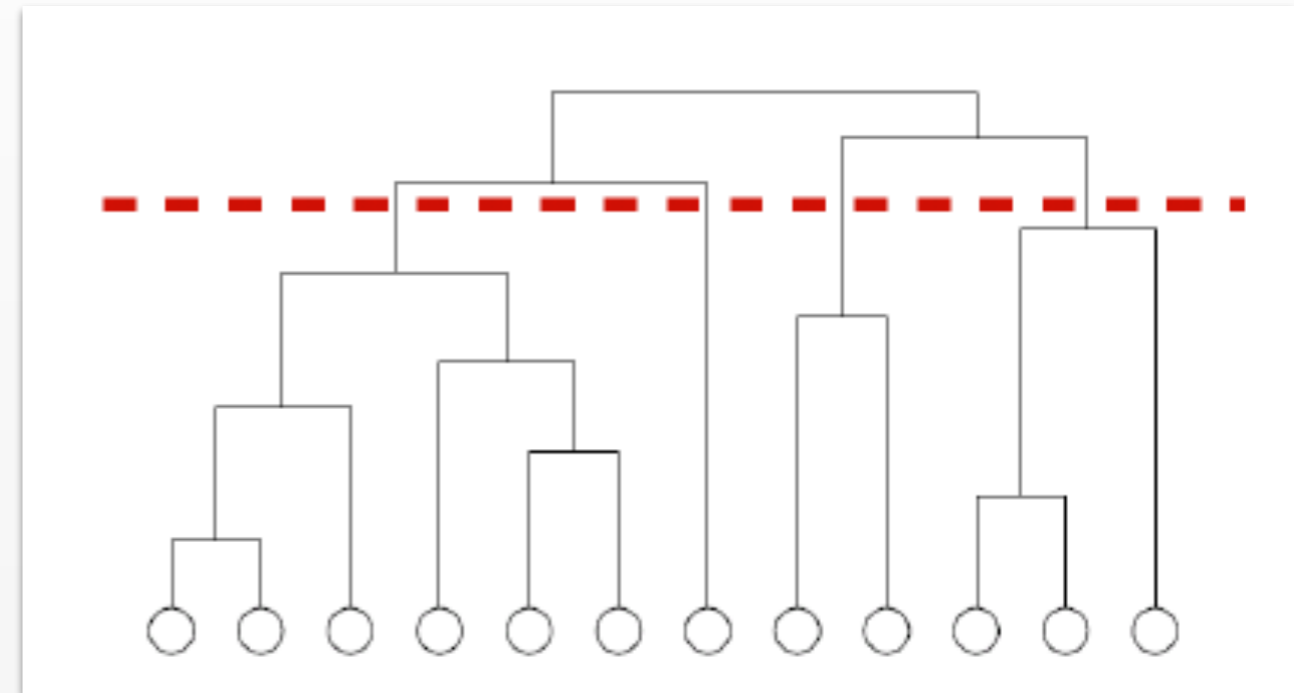
(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

# Determining the Number of Communities

Hierarchical decomposition



Choosing a cut-off



Analogous problem to deciding on number of clusters in hierarchical clustering

# Modularity

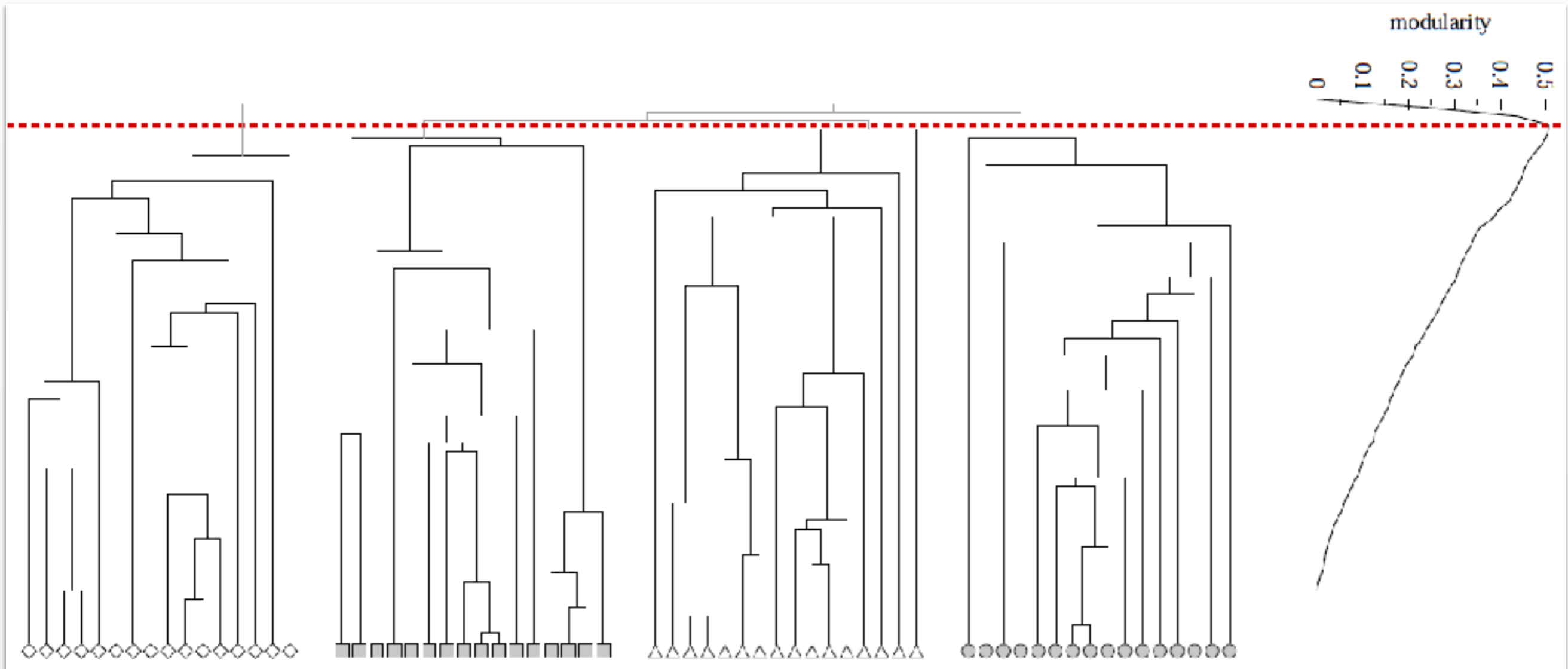
*Idea:* Compare fraction of edges within module to fraction that would be observed for random connections

$$Q = \frac{1}{2m} \sum_{uv} \left[ A_{vw} - \frac{k_v k_w}{2m} \right] \delta(c_u, c_v)$$

- $m$ : Number of edges in graph
- $A_{uv}$ : Adjacency matrix (1 if edge exists 0 otherwise)
- $k_u$ : Degree of node  $u$
- $c_u$ : Cluster assignment for node  $u$



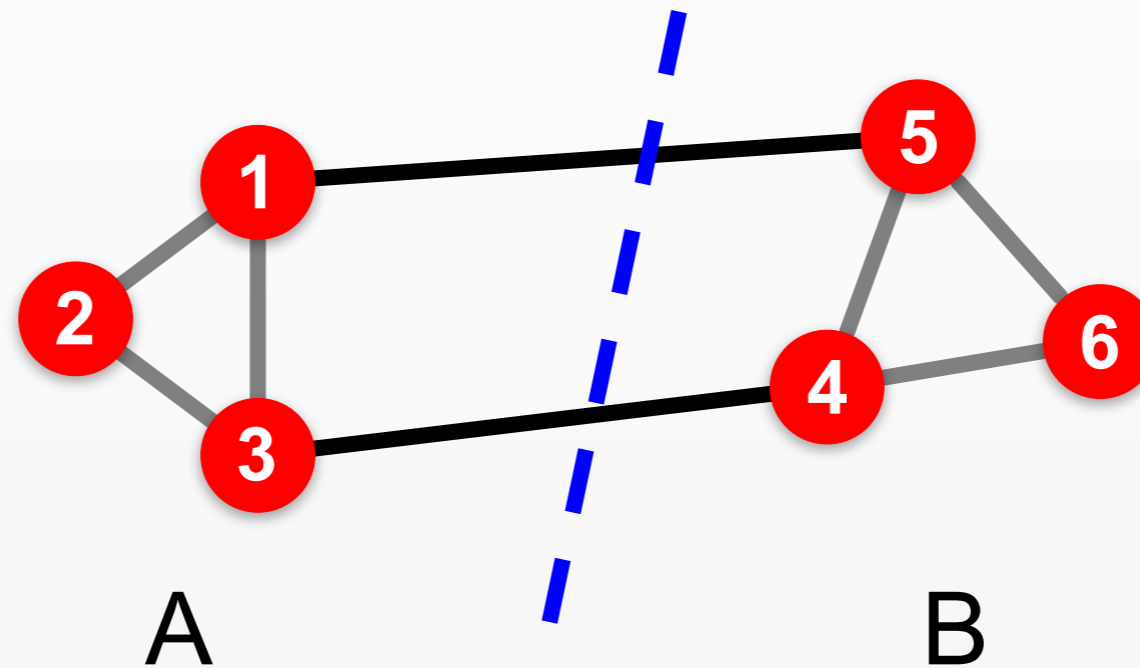
# Modularity



Use modularity to optimize connectivity within modules

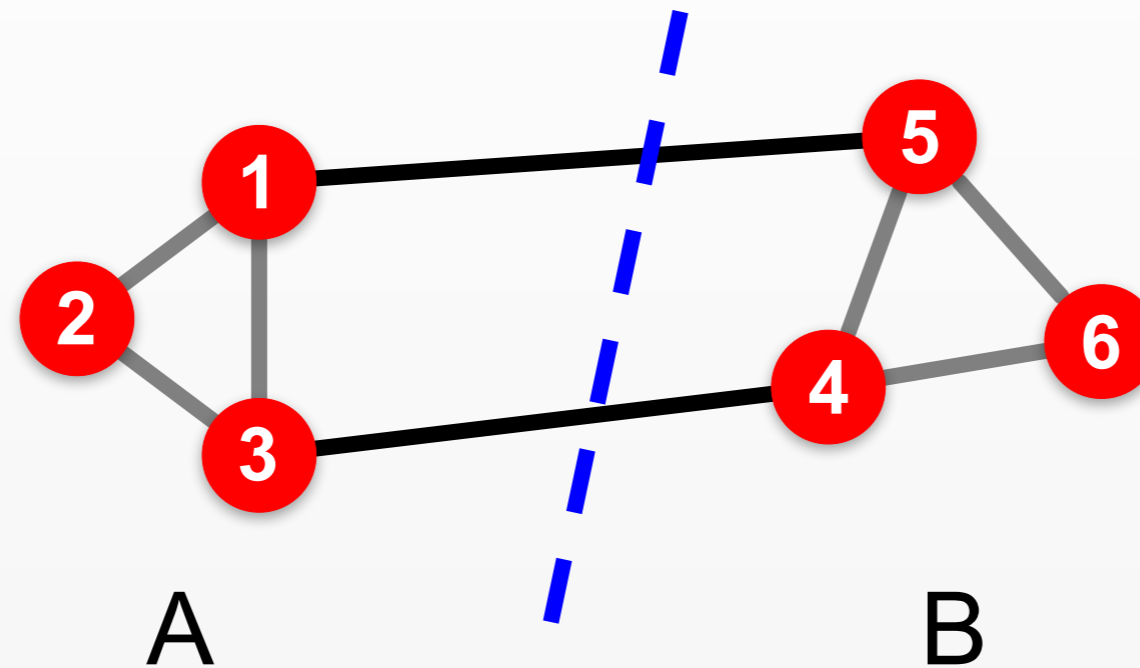
# Spectral Clustering

# Graph Partitioning



- What makes a good partition?
  - Maximize the within-group connections
  - Minimize the between-group connections

# Graph Cuts



Degree

$$d_i = \sum_j A_{ij}$$

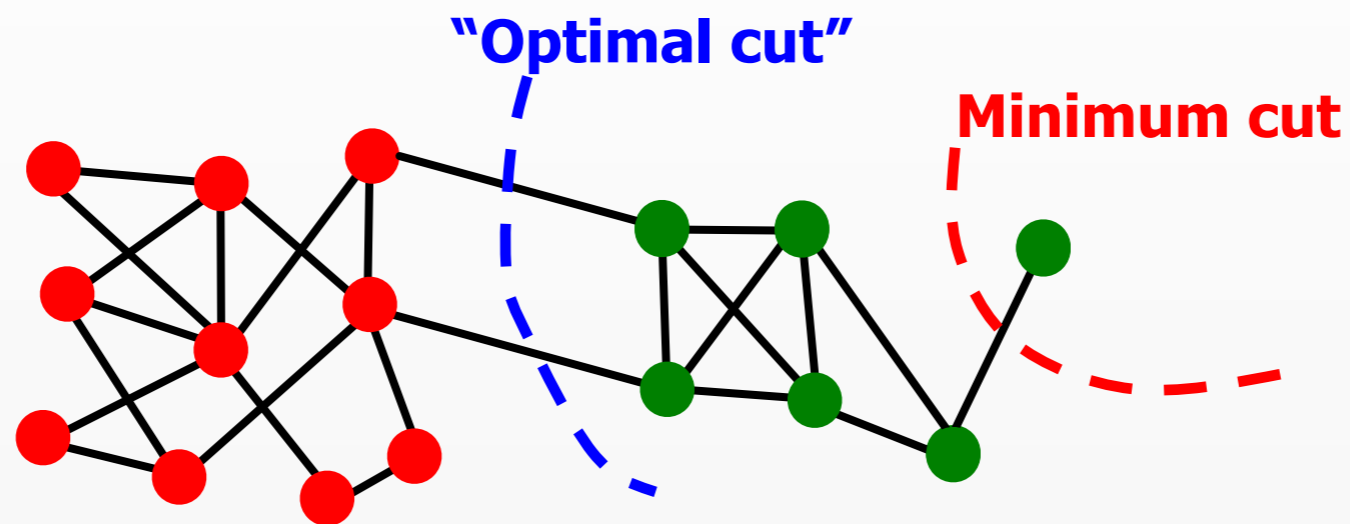
Volume

$$\text{vol}(A) = \sum_j d_i$$

Cut

$$\text{cut}(A, B) = \sum_{i \in A, j \in B} A_{ij}$$

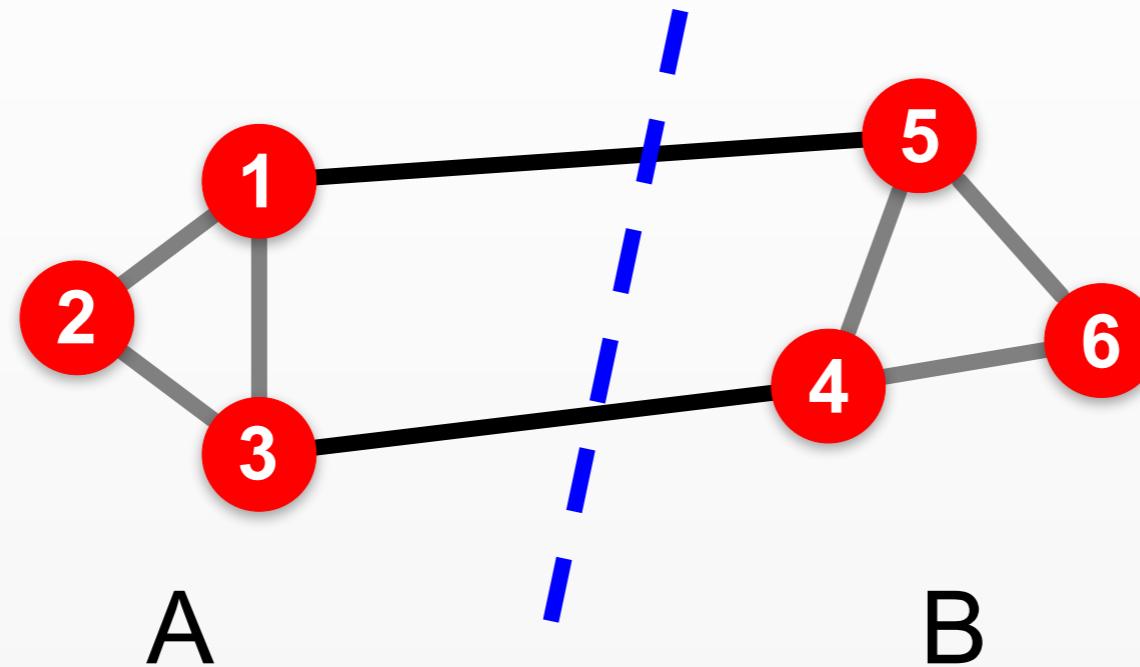
# Minimal Cuts



$$\mathbf{arg\ min}_{A,B} \mathbf{cut}(A,B)$$

*Problem:* minimal cut is not necessarily a good splitting criterion

# Normalized Cuts



$$\text{ncut}(A, B) = \frac{\text{cut}(A, B)}{\text{vol}(A)} + \frac{\text{cut}(A, B)}{\text{vol}(B)}$$

Degree

$$d_i = \sum_j A_{ij}$$

Volume

$$\text{vol}(A) = \sum_j d_i$$

Cut

$$\text{cut}(A, B) = \sum_{i \in A, j \in B} A_{ij}$$

# Find Optimal Cut [Fiedler'73]

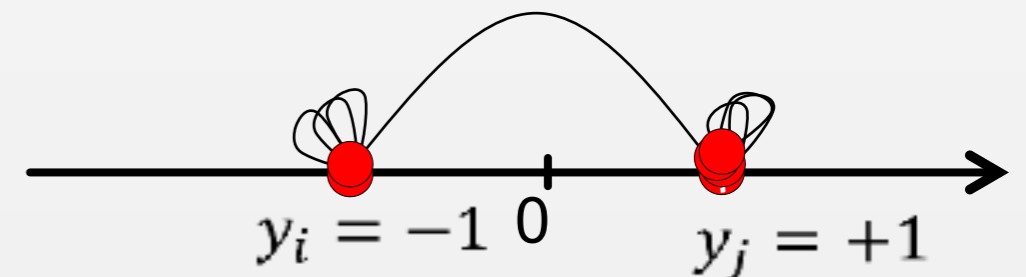
- Back to finding the optimal cut
- Express partition (A,B) as a vector

$$y_i = \begin{cases} +1 & \text{if } i \in A \\ -1 & \text{if } i \in B \end{cases}$$

- We can minimize the cut of the partition by finding a non-trivial vector  $x$  that **minimizes**:

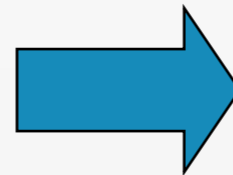
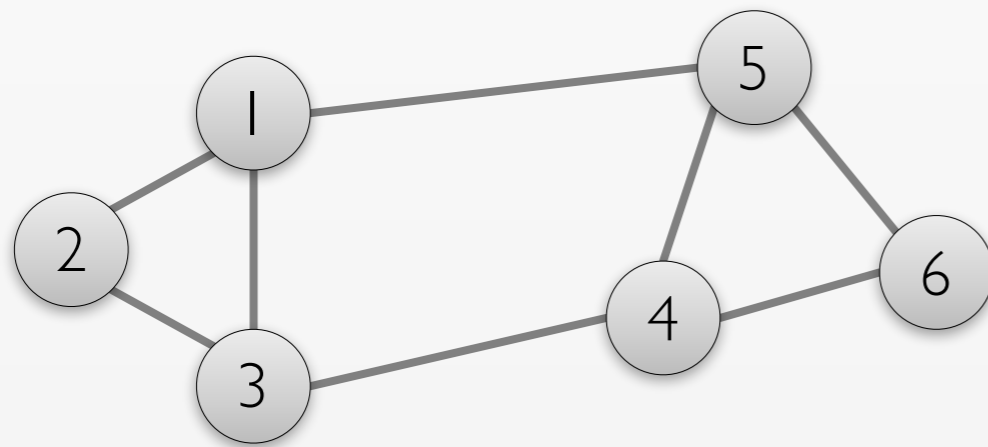
$$y^* = \operatorname{argmin}_{y \in \{-1,1\}^n} \sum_{(i,j) \in E} (y_i - y_j)^2$$

Can't solve exactly. Let's relax  $y$  and allow it to take any real value.



# Matrix Representations

- Adjacency matrix ( $A$ ):
  - $n \times n$  matrix
  - $A=[a_{ij}]$ ,  $a_{ij}=1$  if edge between node  $i$  and  $j$



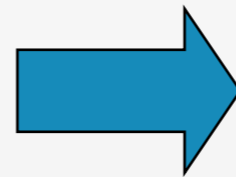
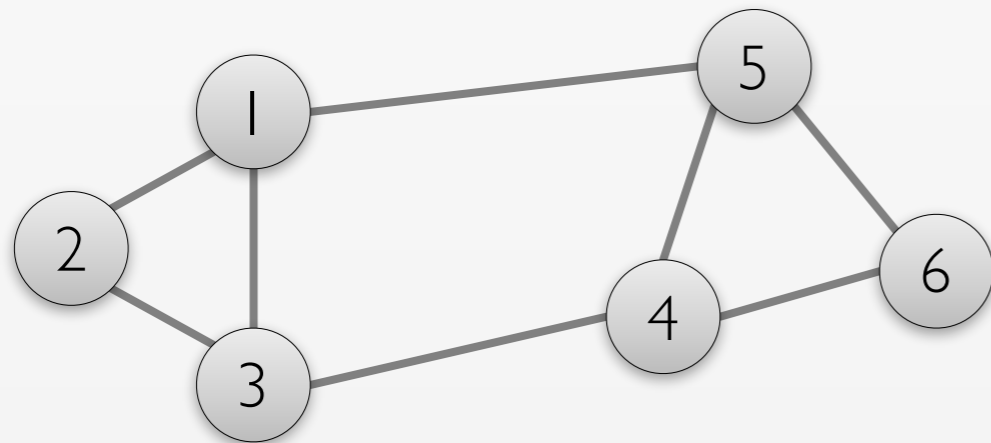
	1	2	3	4	5	6
1	0	1	1	0	1	0
2	1	0	1	0	0	0
3	1	1	0	1	0	0
4	0	0	1	0	1	1
5	1	0	0	1	0	1
6	0	0	0	1	1	0

- Important properties:
  - Symmetric matrix
  - Eigenvectors are real and orthogonal



# Matrix Representations

- Degree matrix ( $D$ ):
  - $n \times n$  diagonal matrix
  - $D=[d_{ii}]$ ,  $d_{ii}$  = degree of node  $i$

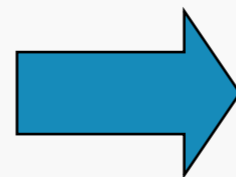
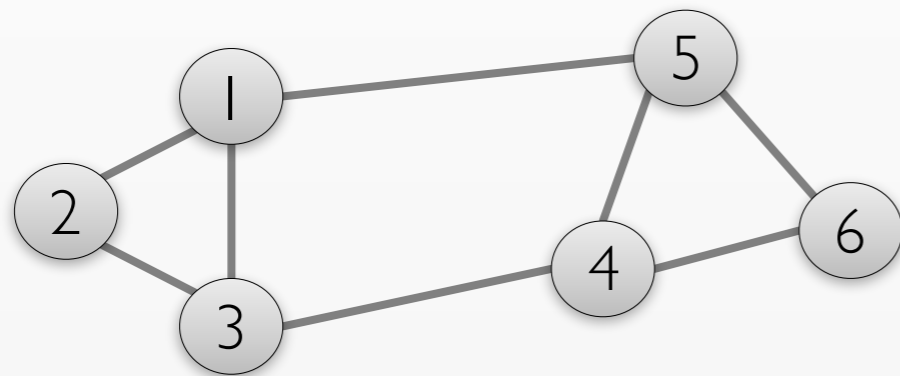


	1	2	3	4	5	6
1	3	0	0	0	0	0
2	0	2	0	0	0	0
3	0	0	3	0	0	0
4	0	0	0	3	0	0
5	0	0	0	0	3	0
6	0	0	0	0	0	2

# Matrix Representations

- **Laplacian matrix (L):**

- $n \times n$  symmetric matrix



	1	2	3	4	5	6
1	3	-1	-1	0	-1	0
2	-1	2	-1	0	0	0
3	-1	-1	3	-1	0	0
4	0	0	-1	3	-1	-1
5	-1	0	0	-1	3	-1
6	0	0	0	-1	-1	2

- **What is trivial eigenpair?**

- $x = (1, \dots, 1)$  then  $L \cdot x = \mathbf{0}$  and so  $\lambda = \lambda_1 = \mathbf{0}$

- **Important properties:**

- **Eigenvalues** are non-negative real numbers
- **Eigenvectors** are real and orthogonal

# Second Eigenvalue

- **Fact: For symmetric matrix  $M$ :**

$$\lambda_2 = \min_x \frac{x^T M x}{x^T x}$$

- **What is the meaning of  $\min x^T L x$  on  $G$ ?**

- $x^T L x = \sum_{i,j=1}^n L_{ij} x_i x_j = \sum_{i,j=1}^n (D_{ij} - A_{ij}) x_i x_j$
- $= \sum_i D_{ii} x_i^2 - \sum_{(i,j) \in E} 2x_i x_j$
- $= \sum_{(i,j) \in E} \underbrace{(x_i^2 + x_j^2)} - 2x_i x_j = \sum_{(i,j) \in E} (x_i - x_j)^2$

Node  $i$  has degree  $d_i$ . So, value  $x_i^2$  needs to be summed up  $d_i$  times.  
But each edge  $(i,j)$  has two endpoints so we need  $x_i^2 + x_j^2$

# Second Eigenvector of Laplacian

- **What else do we know about  $x$ ?**

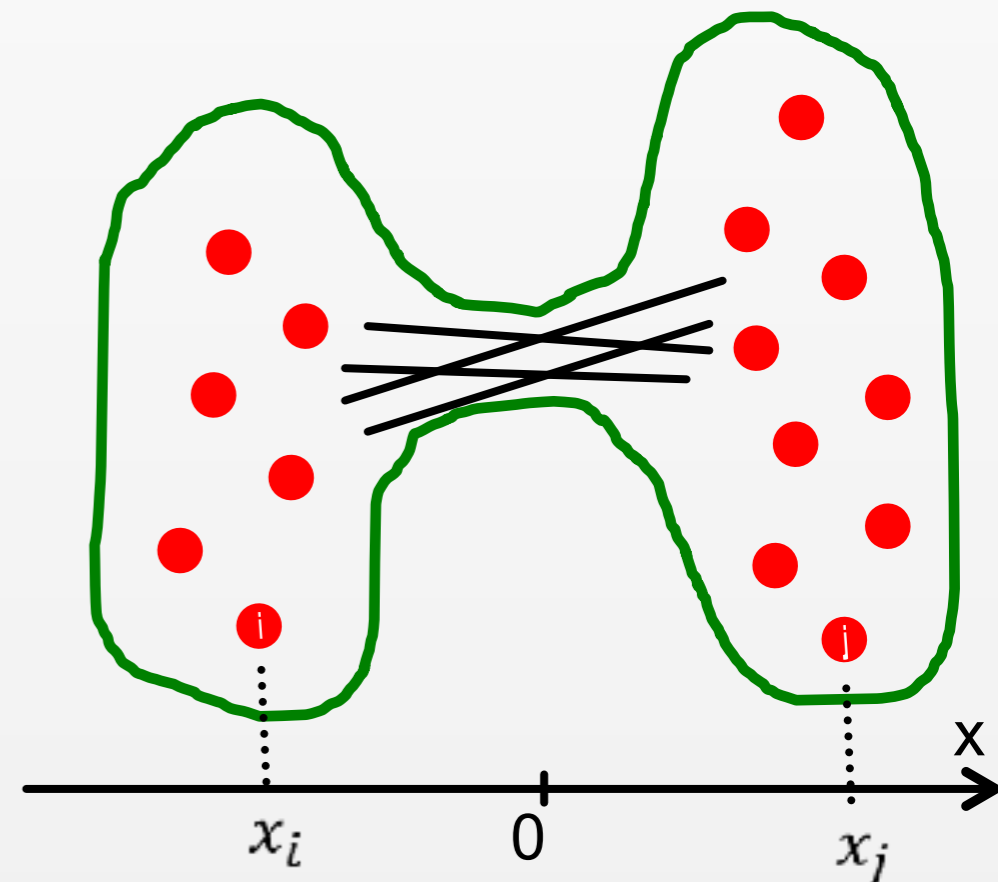
- $x$  is unit vector:  $\sum_i x_i^2 = 1$
- $x$  is orthogonal to 1<sup>st</sup> eigenvector  $(1, \dots, 1)$  thus:  
 $\sum_i x_i \cdot 1 = \sum_i x_i = 0$

- **Remember:**

$$\lambda_2 = \min_{\substack{\text{All labelings} \\ \text{of nodes } i \text{ so} \\ \text{that } \sum x_i = 0}} \frac{\sum_{(i,j) \in E} (x_i - x_j)^2}{\sum_i x_i^2}$$

**We want to assign values  $x_i$  to nodes  $i$  such that few edges cross 0.**

**(we want  $x_i$  and  $x_j$  to subtract each other)**



**Balance to minimize**

# Rayleigh Theorem

$$\min_{y \in \mathbb{R}^n} f(y) = \sum_{(i,j) \in E} (y_i - y_j)^2 = y^T L y$$

- $\lambda_2 = \min_y f(y)$ : The minimum value of  $f(y)$  is given by the 2<sup>nd</sup> smallest eigenvalue  $\lambda_2$  of the Laplacian matrix  $L$
- $x = \arg \min_y f(y)$ : The optimal solution for  $y$  is given by the corresponding eigenvector  $x$ , referred as the **Fiedler vector**

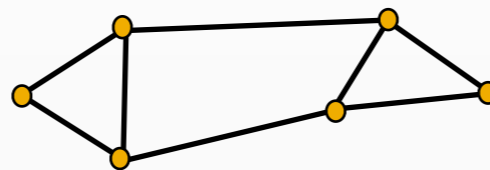
# Spectral Clustering Algorithms

- Three basic stages:
  - 1) Pre-processing
    - Construct a matrix representation of the graph
    - More generally, construct similarity matrix
  - 2) Decomposition
    - Compute eigenvalues and eigenvectors of the matrix
    - Map each point to a lower-dimensional representation based on one or more eigenvectors
  - 3) Grouping
    - Assign points to two or more clusters, based on the new representation

# Spectral Partitioning Algorithm

- 1) **Pre-processing:**

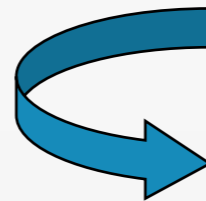
- Build Laplacian matrix  $\mathbf{L}$  of the graph



	1	2	3	4	5	6
1	3	-1	-1	0	-1	0
2	-1	2	-1	0	0	0
3	-1	-1	3	-1	0	0
4	0	0	-1	3	-1	-1
5	-1	0	0	-1	3	-1
6	0	0	0	-1	-1	2

- 2) **Decomposition:**

- Find eigenvalues  $\lambda$  and eigenvectors  $\mathbf{x}$  of the matrix  $\mathbf{L}$



$\lambda =$

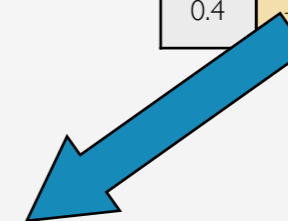
0.0
1.0
3.0
3.0
4.0
5.0

$\mathbf{X} =$

0.4	0.3	-0.5	-0.2	-0.4	-0.5
0.4	0.6	0.4	-0.4	0.4	0.0
0.4	0.3	0.1	0.6	-0.4	0.5
0.4	-0.3	0.1	0.6	0.4	-0.5
0.4	-0.3	-0.5	-0.2	0.4	0.5
0.4	-0.6	0.4	-0.4	-0.4	0.0

- Map vertices to corresponding components of  $\lambda_2$

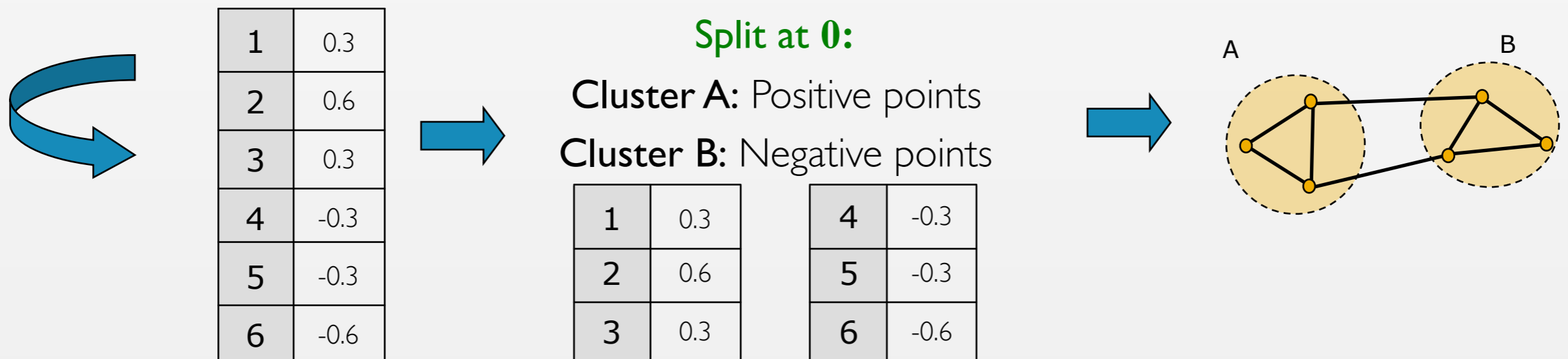
<b>1</b>	0.3
<b>2</b>	0.6
<b>3</b>	0.3
<b>4</b>	-0.3
<b>5</b>	-0.3
<b>6</b>	-0.6



How do we now find the clusters?

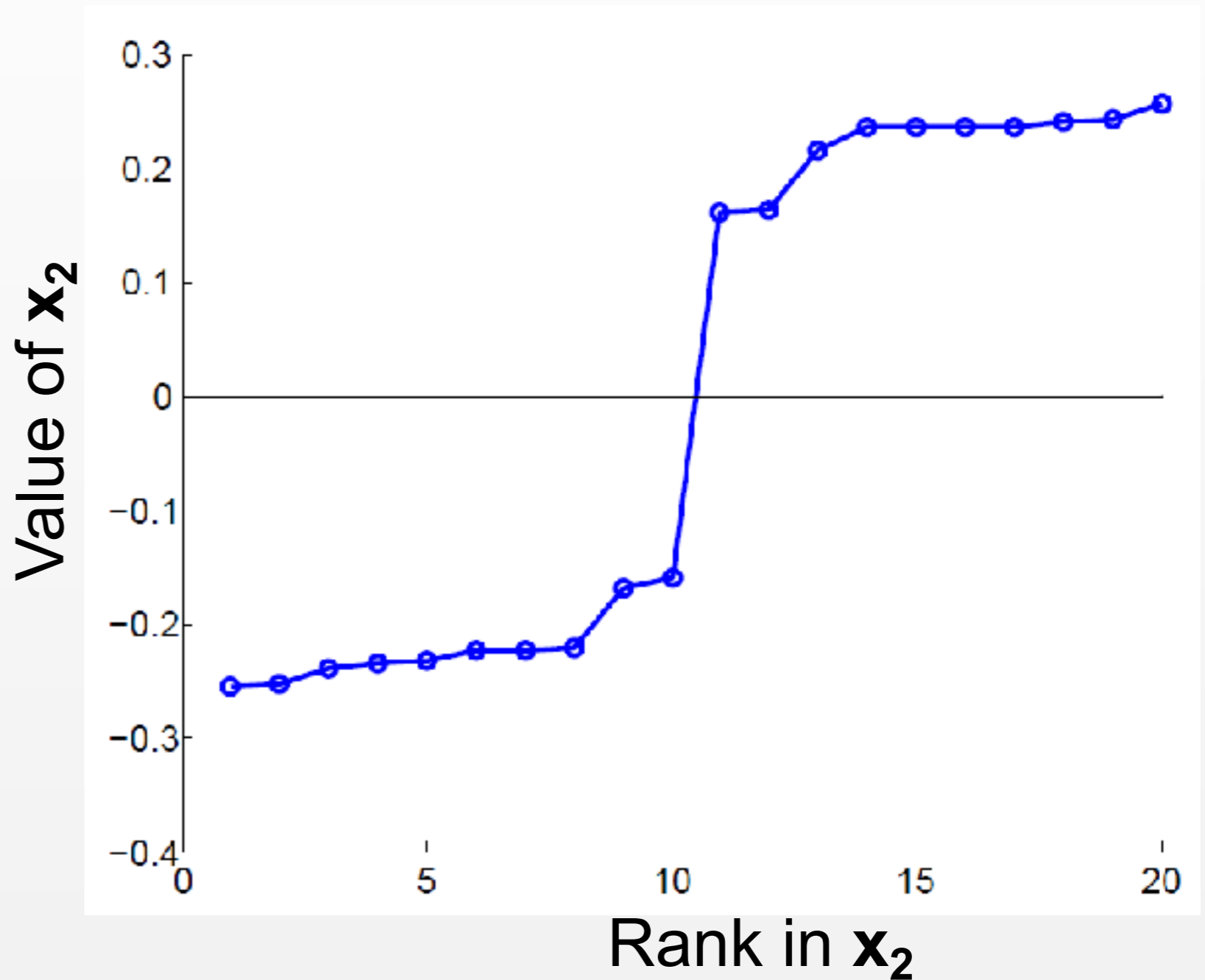
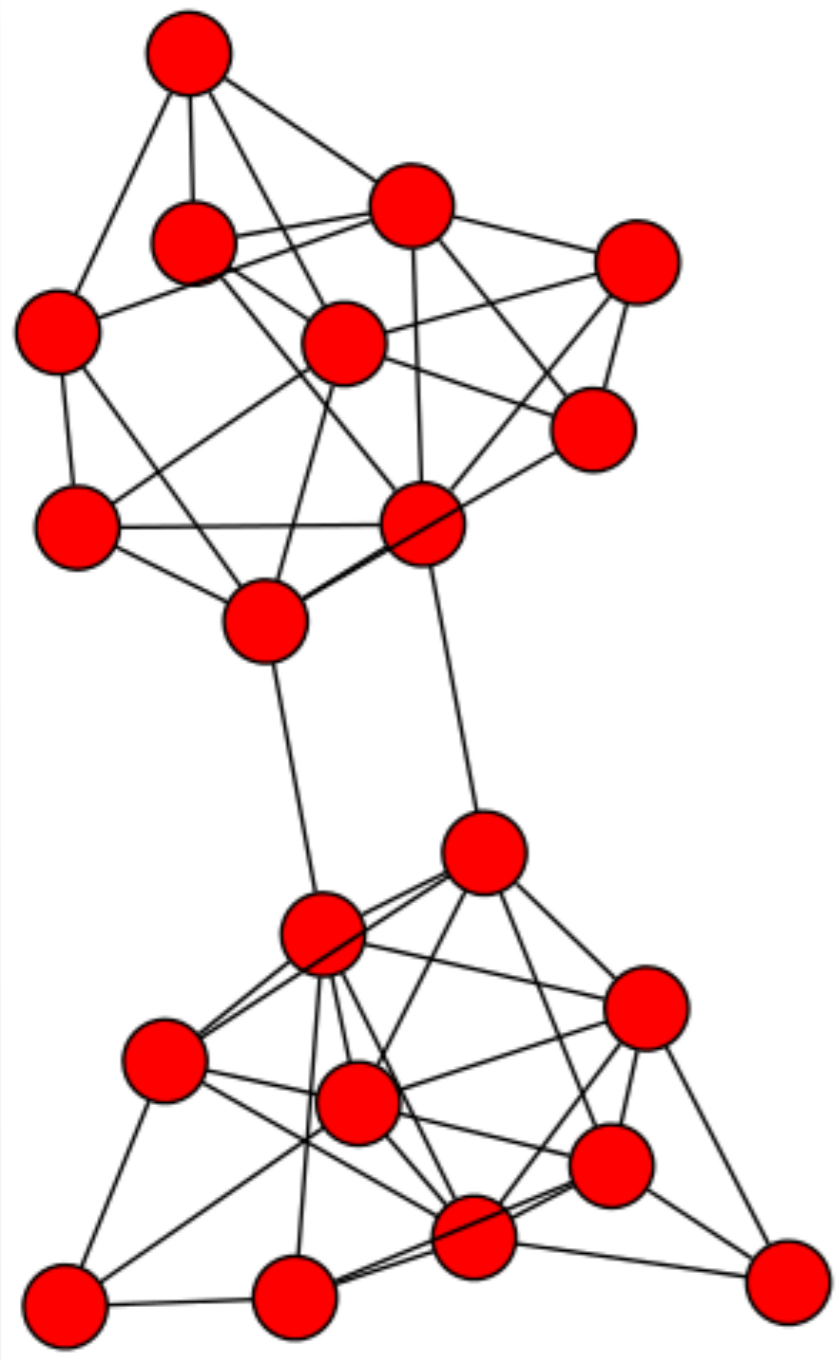
# Spectral Partitioning

- 3) **Grouping:**
  - Sort components of reduced 1-dimensional vector
  - Identify clusters by splitting the sorted vector in two
- **How to choose a splitting point?**
  - Naïve approaches:
    - Split at 0 or median value
  - More expensive approaches:
    - Attempt to minimize normalized cut in 1-dimension (sweep over ordering of nodes induced by the eigenvector)

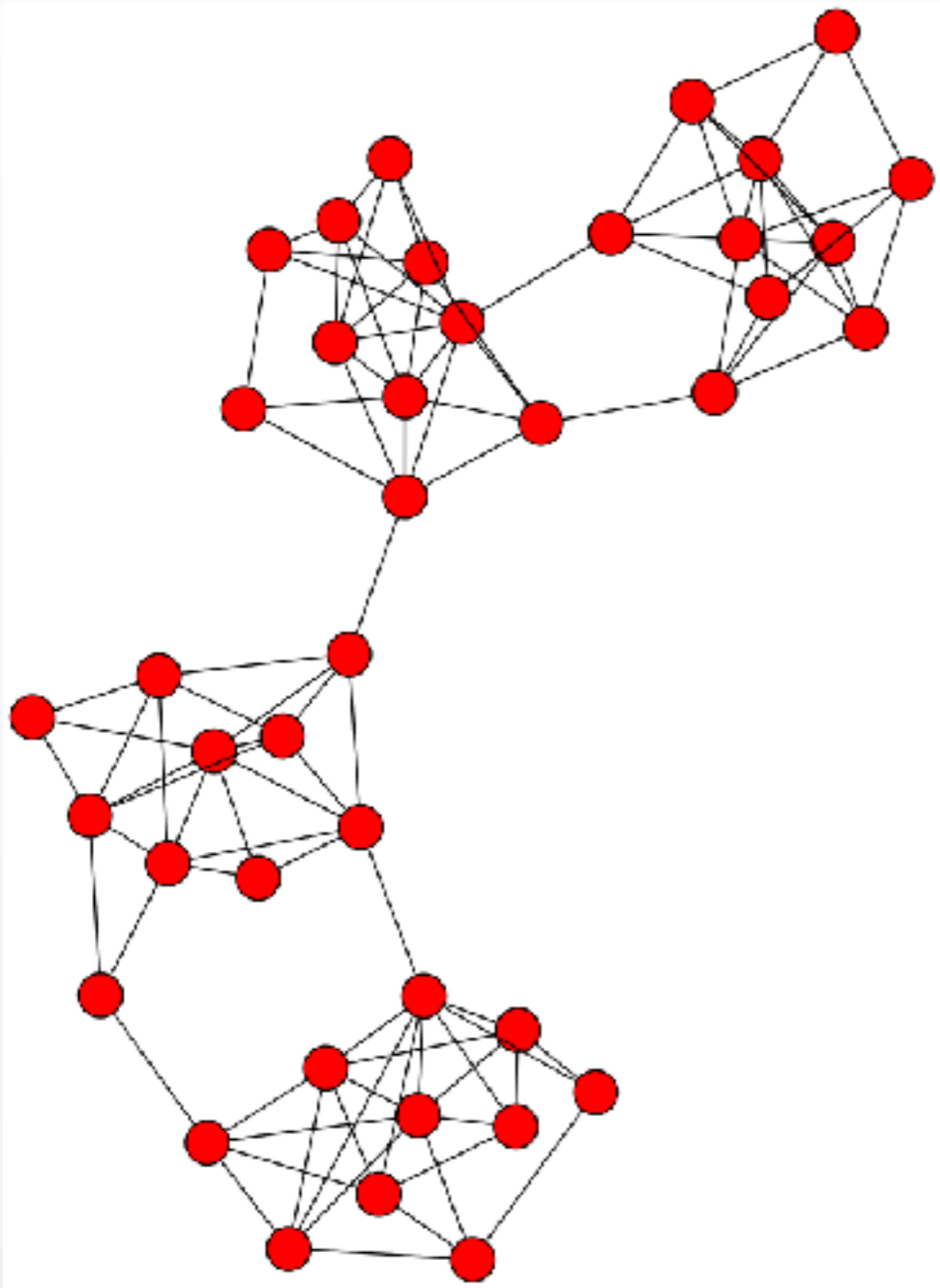




# Example: Spectral Partitioning



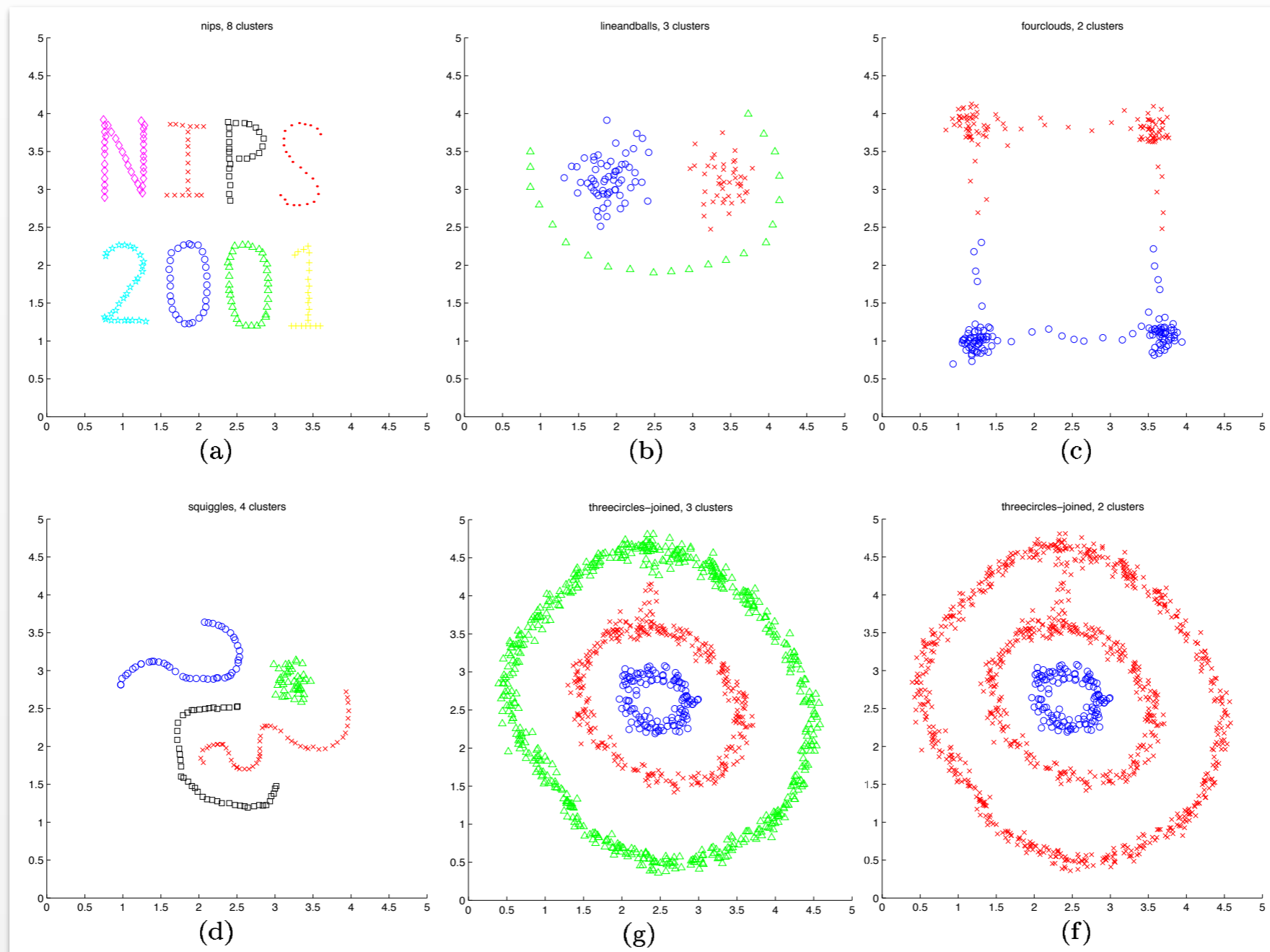
# Example: Spectral Partitioning



# k-Way Spectral Clustering

- How do we partition a graph into  $k$  clusters?
- Two basic approaches:
  - Recursive bi-partitioning [Hagen et al., '92]
    - Recursively apply bi-partitioning algorithm in a hierarchical divisive manner
    - Disadvantages: Inefficient, unstable
  - Cluster multiple eigenvectors [Shi-Malik, '00]
    - Build a reduced space from multiple eigenvectors
    - Commonly used in recent papers

# Spectral Clustering as General-purpose Method



source: Ng, Jordan and Weiss, NIPS 2001

Define “edge weight”  $W$  using some similarity metric  
(e.g. a kernel function)