

# Digital Design

CS/EEE/ECE/INSTR F215

Lecture 4: Boolean algebra



**Birla Institute of Technology & Science, Pilani**  
Hyderabad Campus

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lead



# Boolean Algebra

## Postulates of Boolean Algebra

$$1. x + 0 = x$$

$$2. x + x' = 1$$

$$3. x + y = y + x$$

$$4. x(y + z) = xy + xz$$

## Duality Principle

$$+ (OR) \rightarrow \cdot (AND)$$

$$\cdot (AND) \rightarrow + (OR)$$

$$1 \rightarrow 0$$

$$0 \rightarrow 1$$



# Boolean Algebra

## Applying Duality

$$1. x + 0 = x$$

$$x \cdot 1 = x$$

$$2. x + x' = 1$$

$$x \cdot x' = 0$$

$$3. x + y = y + x$$

$$x \cdot y = y \cdot x$$

$$4. x(y + z) = xy + xz$$

$$x + (y \cdot z) = (x + y) \cdot (x + z)$$



# Boolean Algebra

## Some important Theorems

$$1. x + x = x$$

$$x \cdot x = x$$

$$2. (x')' = x$$

$$3. x + (y + z) = (x + y) + z$$

$$x \cdot (y \cdot z) = (x \cdot y) \cdot z$$

$$4. (x + y)' = x'y' \quad \text{DeMorgan's Th.} \quad (x \cdot y)' = x' + y'$$

$$4. x + xy = x$$

$$x \cdot (x + y) = x$$



# Boolean Algebra

Theorems can be proved in two ways

-using Postulates

$$x + xy = x$$

$$\begin{aligned} LHS &= x \cdot 1 + xy \\ &= x \cdot (1 + y) \\ &= x \cdot 1 \\ &= x \end{aligned}$$



# Boolean Algebra

Theorems can be proved in two ways

-using Truth tables

$$x + xy = x$$

$x$	$y$	$xy$	$x+xy$
0	0	0	0
0	1	0	0
1	0	0	1
1	1	1	1

=

$x$
0
0
1
1

*LHS* = *RHS*



# Boolean Functions

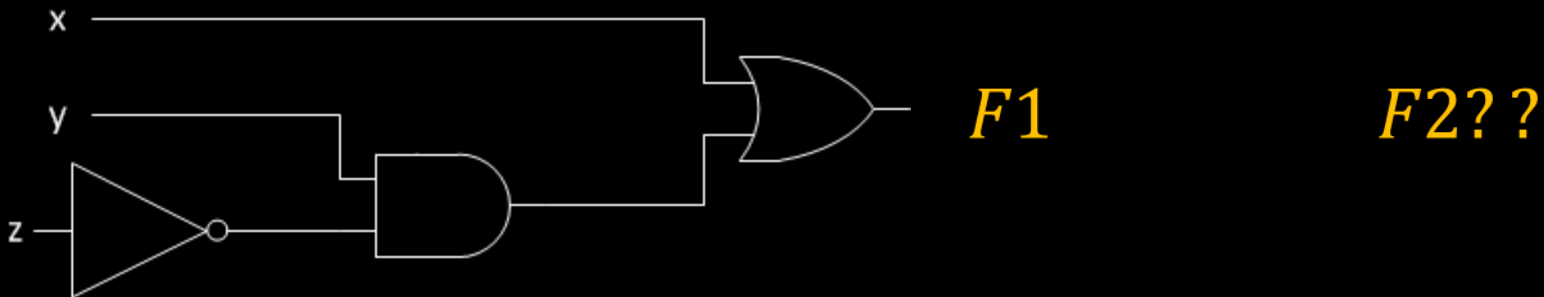
Consists of **binary variables**, constants and Logic symbols

$$F1 = x + yz'$$

$$F2 = (x + yz')'$$

Single variable in normal or complemented form - **literal**

Can be transformed into **circuit**





# Boolean Functions

Boolean function can be represented by truth table in only one way

<i>x</i>	<i>y</i>	<i>z</i>	<i>F1</i>	<i>F2</i>
0	0	0	0	1
0	0	1	0	1
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

$$F1 = x + yz'$$

$$F2 = (x + yz')'$$





# Boolean Functions

Boolean function can be represented in **many ways in algebraic form**

$$F2 = (x + yz')'$$

**Both represent same function ?**

$$F2 = x'yz' + x'yz + xz'$$

**Implementation of both same ? Try it out**

**Need for simplification**



# Boolean Functions

## Ways of simplification

-Algebraic manipulation

-K Maps

-QM(Quine-McCluskey) Method



# Boolean Functions

Algebraic manipulation done using postulates and theorems

For example:  $F1 = x(x' + y)$       1 Not, 1 OR, 1 AND

$$F1 = xx' + xy$$

$$F1 = xy$$
      1 AND



# Boolean Functions

## Canonical Forms-Type1

x	y	F1
0	0	0
0	1	0
1	0	0
1	1	1

$$F1 = xy + F2 = x'y'$$

x	y	F2
0	0	1
0	1	0
1	0	0
1	1	0

x	y	F3
0	0	1
0	1	0
1	0	0
1	1	1

$$F3 = ???$$

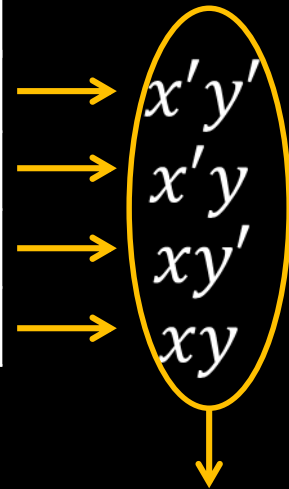
$$F3 = xy + x'y'$$



# Boolean Functions

Each input entry will be represented by a term

$x$	$y$
0	0
0	1
1	0
1	1



Minterms

Standard Product

Designation

$m_0$

$m_1$

$m_2$

$m_3$

$F3$
1
0
1
1

$$F3 = x'y' + xy' + xy$$

$$F3 = m_0 + m_2 + m_3$$

$$F3 = \sum(0,2,3)$$

Sum of Products

(Canonical form-Type1)



# Boolean Functions

## Canonical Forms-Type2

x	y	F1
0	0	0
0	1	1
1	0	1
1	1	1

$$F1 = x + y \quad F2 = x' + y'$$

x	y	F2
0	0	1
0	1	1
1	0	1
1	1	0

x	y	F3
0	0	0
0	1	1
1	0	1
1	1	0

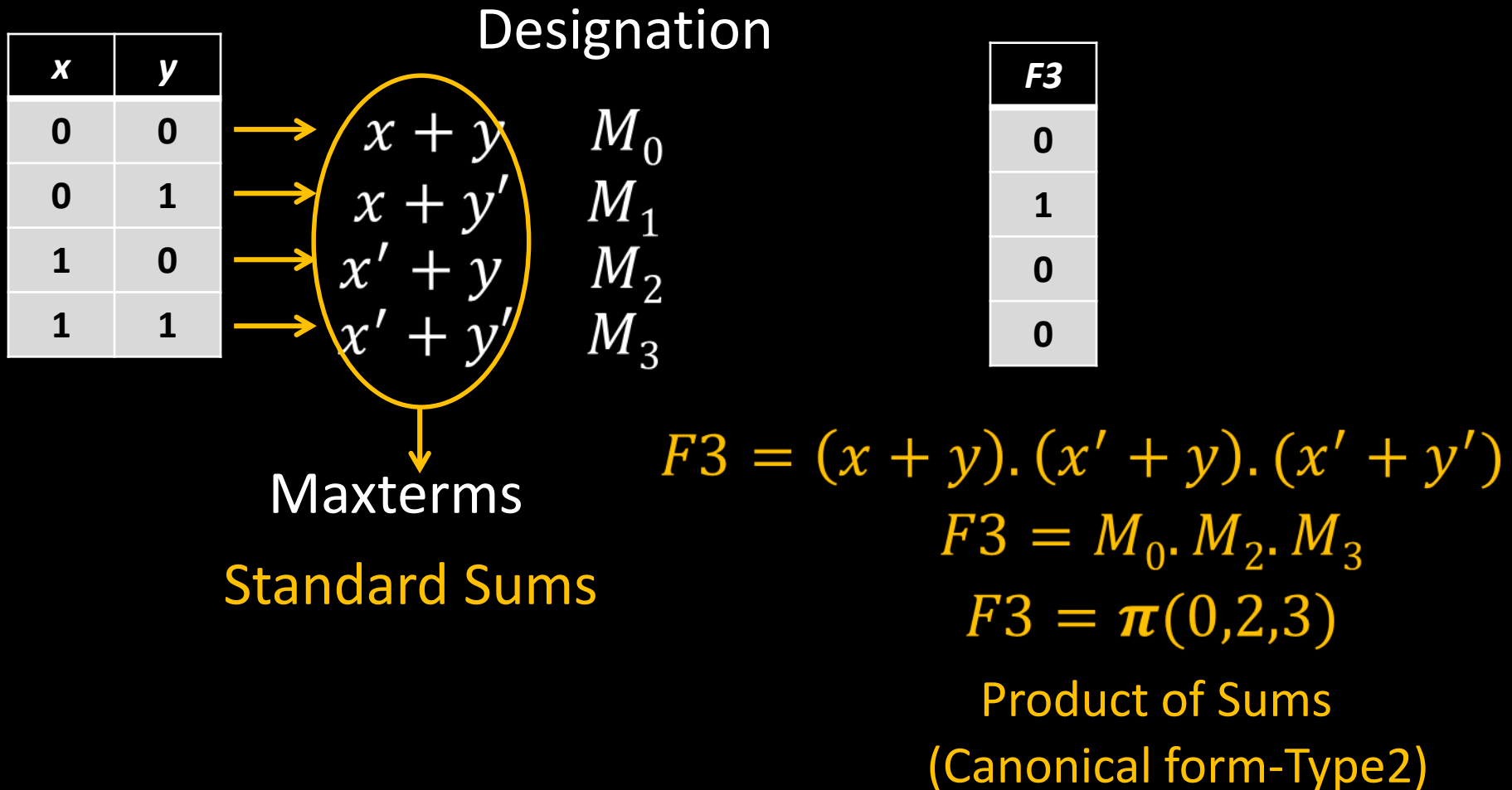
$$F3 = ??$$

$$F3 = (x + y) \cdot (x' + y')$$



# Boolean Functions

Each input entry will be represented by a term





# Boolean Algebra

## Some important Theorems

1.  $x + x = x$

$x \cdot x = x$

2.  $(x')' = x$

3.  $x + (y + z) = (x + y) + z$

$x \cdot (y \cdot z) = (x \cdot y) \cdot z$

4.  $(x + y)' = x'y'$  DeMorgan's Th.  $(x \cdot y)' = x' + y'$





# Boolean Algebra

Theorems can be proved in two ways

-using Truth tables

$$x + xy = x$$

$x$	$y$	$xy$	$x+xy$
0	0	0	0
0	1	0	0
1	0	0	1
1	1	1	1

=

$x$
0
0
1
1

*LHS* = *RHS*



# Boolean Functions

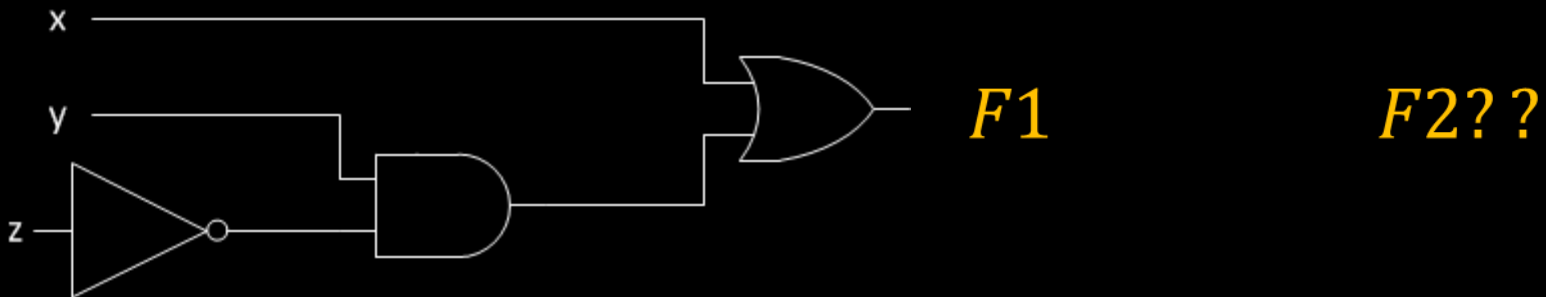
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$$F2 = (x + yz')'$$

Single variable in normal or complemented form - **literal**

Can be transformed into **circuit**





# Boolean Functions

Boolean function can be represented by truth table in only one way

<i>x</i>	<i>y</i>	<i>z</i>	<i>F1</i>	<i>F2</i>
0	0	0	0	1
0	0	1	0	1
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

$$F1 = x + yz'$$

$$F2 = (x + yz')'$$



# Boolean Functions

Boolean function can be represented in **many ways in algebraic form**

$$F2 = (x + yz')'$$

**Both represent same function**

$$F2 = x'y'z' + x'y'z + x'z$$

**Implementation of both same? Try it out**

**Need for simplification**



# Boolean Functions

Algebraic manipulation done using postulates and theorems

For example:  $F1 = x(x' + y)$       1 Not, 1 OR, 1 AND

$$F1 = xx' + xy$$

$$F1 = xy$$
      1 AND



# Boolean Functions

Before learning

-K Maps

-QM(Quine-McCluskey) Method

Canonical and standard forms of Boolean functions



# Boolean Functions

## Canonical Forms-Type1

x	y	F1
0	0	0
0	1	0
1	0	0
1	1	1

$$F1 = xy + F2 = x'y'$$

x	y	F2
0	0	1
0	1	0
1	0	0
1	1	0

x	y	F3
0	0	1
0	1	0
1	0	0
1	1	1

$$F3 = ??$$

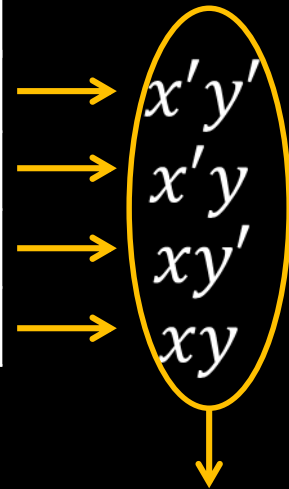
$$F3 = xy + x'y'$$



# Boolean Functions

Each input entry will be represented by a term

$x$	$y$
0	0
0	1
1	0
1	1



Minterms

Standard Product

Designation

$m_0$

$m_1$

$m_2$

$m_3$

$F3$
1
0
1
1

$$F3 = x'y' + xy' + xy$$

$$F3 = m_0 + m_2 + m_3$$

$$F3 = \sum(0,2,3)$$

Sum of Products

(Canonical form-Type1)





# Boolean Functions

## Canonical Forms-Type2

$x$	$y$	$F1$
0	0	0
0	1	1
1	0	1
1	1	1

$$F1 = x + y \quad . \quad F2 = x' + y'$$

$x$	$y$	$F2$
0	0	1
0	1	1
1	0	1
1	1	0

$x$	$y$	$F3$
0	0	0
0	1	1
1	0	1
1	1	0

$$F3 = ??$$

$$F3 = (x + y) \cdot (x' + y')$$



# Boolean Functions

Each input entry will be represented by a term

$x$	$y$		Designation	
0	0	→	$x + y$	$M_0$
0	1	→	$x + y'$	$M_1$
1	0	→	$x' + y$	$M_2$
1	1	→	$x' + y'$	$M_3$

$F3$
0
1
0
0

Maxterms

Standard Sums

$$F3 = (x + y). (x' + y). (x' + y')$$

$$F3 = M_0. M_2. M_3$$

$$F3 = \pi(0,2,3)$$

Product of Sums

(Canonical form-Type2)



# Boolean Functions

## Canonical forms from Truth table

x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Sum of Products Representation ?

Product of Sums Representation ?



# Boolean Functions

Can we convert from one canonical form to another ?

x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Using Truth Table

$$F = m_1 + m_2 + m_4 + m_7$$

*(Three variables – 0,1,2,3,4,5,6,7)*

$$F = M_0 \cdot M_3 \cdot M_5 \cdot M_6$$

**Minterm and Maxterm are Complementary to each other**



# Boolean Functions

x	y	z	F	F'
0	0	0	0	1
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	0

$$F = m_1 + m_2 + m_4 + m_7$$

$$F' = m_0 + m_3 + m_5 + m_6$$

$$F' = x'y'z' + x'yz + xy'z + xyz'$$

$$F = (F')'$$

$$F = (x'y'z' + x'yz + xy'z + xyz')'$$

$$F = (x'y'z')'.(x'yz)'.(xy'z)'.(xyz')'$$

$$F = (x + y + z).(x + y' + z').(x' + y + z').(x' + y' + z)$$

$$F = M_0.M_3.M_5.M_6$$



# Boolean Functions

*Express in minterms:*

**Example 1:**  $F = y + x'yz'$

$$F = y(x + x')(z + z') + x'yz' \quad \text{Simplify}$$

$$F = xyz + xyz' + x'yz + x'yz' + x'yz'$$

$$F = m_7 + m_6 + m_3 + m_2$$



# Boolean Functions

*Express in minterms:*

**Example 1:  $F = y + x'yz'$**

x	y	z	$x'yz'$	F
0	0	0	0	0
0	0	1	0	0
0	1	0	1	1
0	1	1	0	1
1	0	0	0	0
1	0	1	0	0
1	1	0	0	1
1	1	1	0	1

$$F = m_2 + m_3 + m_6 + m_7$$

$$F = M_0 \cdot M_1 \cdot M_4 \cdot M_5$$



# Boolean Functions

Standard Form – Terms that form the function may contain one, two or many literals

$$F1 = y' + xy + x'yz'$$

Sum of Products

$$F2 = x(y' + z).(x + y + z')$$

Product of Sums

How is it different from canonical form ??

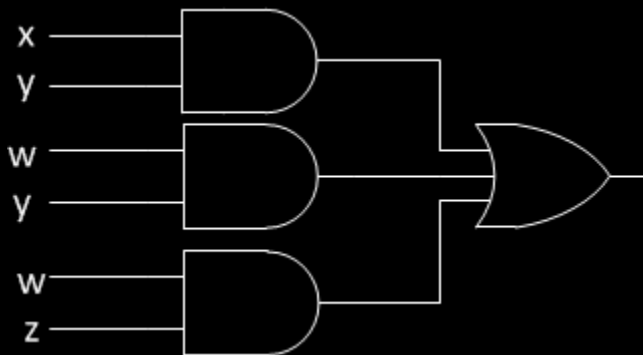
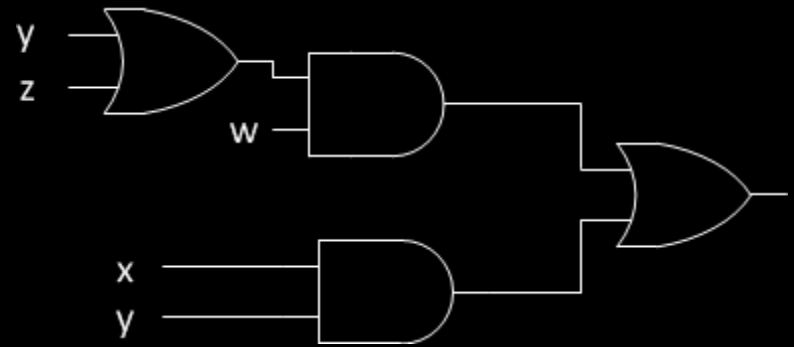


# Boolean Functions

Some times standard forms result in better implementation

What form ??

$$F = xy + w(y + z)$$



$$F = xy + wy + wz$$

Standard Form

SPOT THE DIFF ??



# Thank You