

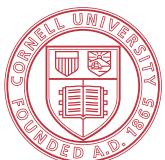


ECE 5775

High-Level Digital Design Automation

Fall 2022

Tutorial on C-Based HLS



Cornell University



Announcements

- ▶ Lab 1 will be released today
 - Friday 09/09
- ▶ Complete the [paper](#) reading before next lecture
- ▶ TA office hours posted on Ed
 - Wednesday 4:30-5:30pm, Online
 - Friday 3:30-4:30pm, Rhodes 380

Agenda

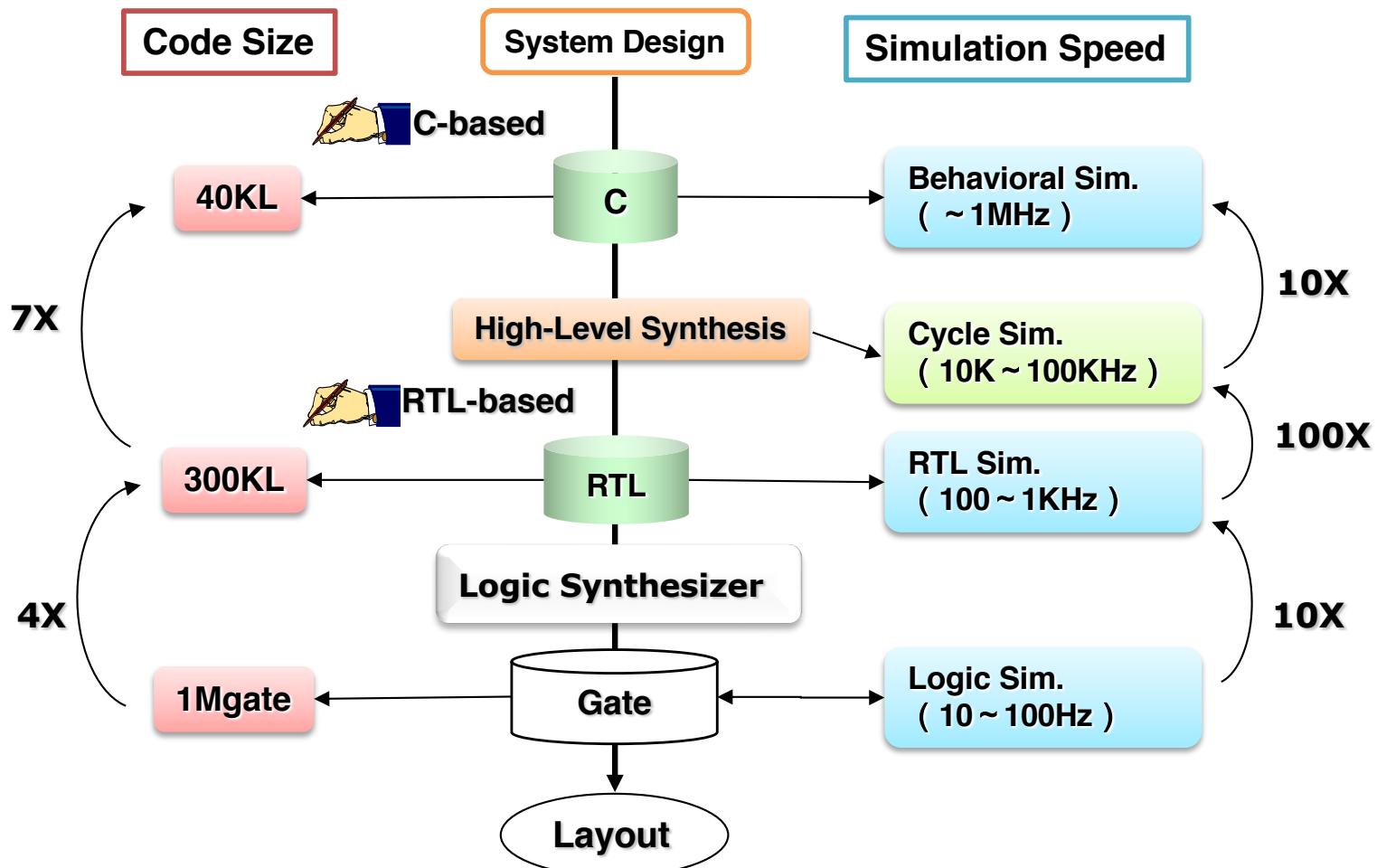
- ▶ Introduction to high-level synthesis (HLS)
 - C-based synthesis
 - Common HLS optimizations
- ▶ Matrix-vector multiplication using Vivado HLS
 - Led by Niansong Zhang (PhD TA)

High-Level Synthesis (HLS)

- ▶ **What**
 - An automated design process that transforms **high-level functional specifications** into **optimized register-transfer level (RTL)** descriptions for efficient hardware implementation
 - Input spec. to HLS is typically untimed or partially timed
- ▶ **Why**
 - **Productivity:** Lower design complexity & faster simulation speed
 - **Portability:** Single (untimed) source → multiple implementations
 - **Quality:** Quicker design space exploration → higher quality

HLS Raises the Level of Design Abstraction

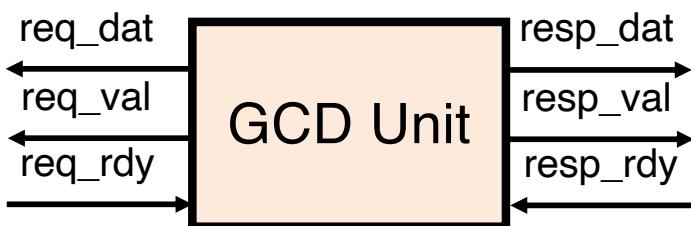
- ▶ Significant code size reduction
- ▶ Shorter simulation/verification cycle



[source: Wakabayashi, DAC'05 tutorial]

A Simple Example: RTL vs. HLS

- ▶ A GCD unit with handshake



HLS Code

```
void GCD ( msg& req,
           msg& resp ) {
    short a = req.msg_a;
    short b = req.msg_b;
    while ( a != b ) {
        if (a > b)
            a = a - b;
        else
            b = b - a;
    }
    resp.msg = a;
}
```

Manual RTL (partial)

```
module GcdUnitRTL
(
    input wire [ 0:0] clk,
    input wire [ 31:0] req_dat,
    output wire [ 0:0] req_rdy,
    input wire [ 0:0] req_val,
    input wire [ 0:0] reset,
    output wire [ 15:0] resp_dat,
    input wire [ 0:0] resp_rdy,
    output wire [ 0:0] resp_val
);
```

```
always @ (*) begin
    if ((curr_state_0 == STATE_IDLE))
        if (req_val)
            next_state_0 = STATE_CALC;

    if ((curr_state_0 == STATE_CALC))
        if (!is_a_lt_b&&is_b_zero)
            next_state_0 = STATE_DONE;

    if ((curr_state_0 == STATE_DONE))
        if ((resp_val&&resp_rdy))
            next_state_0 = STATE_IDLE;
end
```

Module declaration

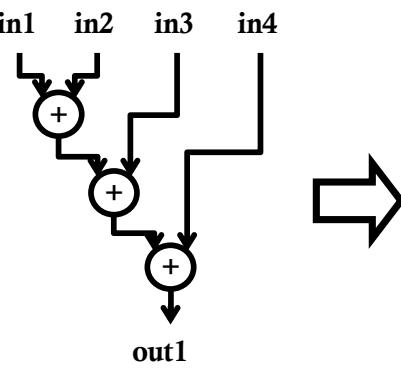
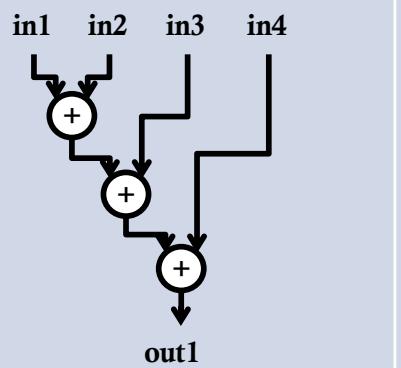
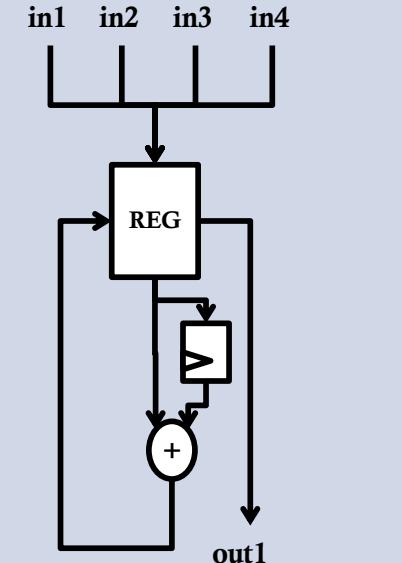
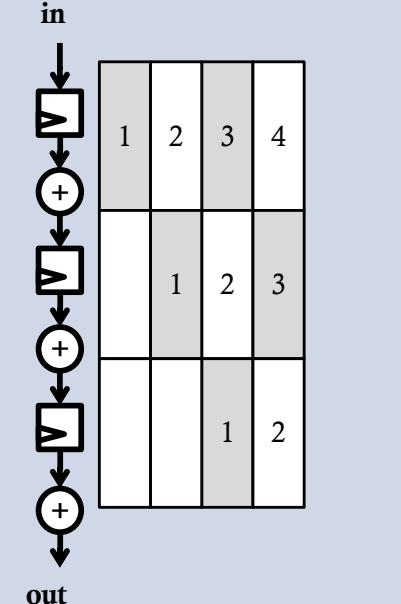
```
always @ (*) begin
    if ((current_state_1 == STATE_IDLE))
        req_rdy = 1; resp_val = 0;
        a_mux_sel = A_MUX_SEL_IN; b_mux_sel = B_MUX_SEL_IN;
        a_reg_en = 1; b_reg_en = 1;

    if ((current_state_1 == STATE_CALC))
        do_swap = is_a_lt_b; do_sub = ~is_b_zero;
        req_rdy = 0; resp_val = 0;
        a_mux_sel = do_swap ? A_MUX_SEL_B : A_MUX_SEL_SUB;
        a_reg_en = 1; b_reg_en = do_swap;
        b_mux_sel = B_MUX_SEL_A;

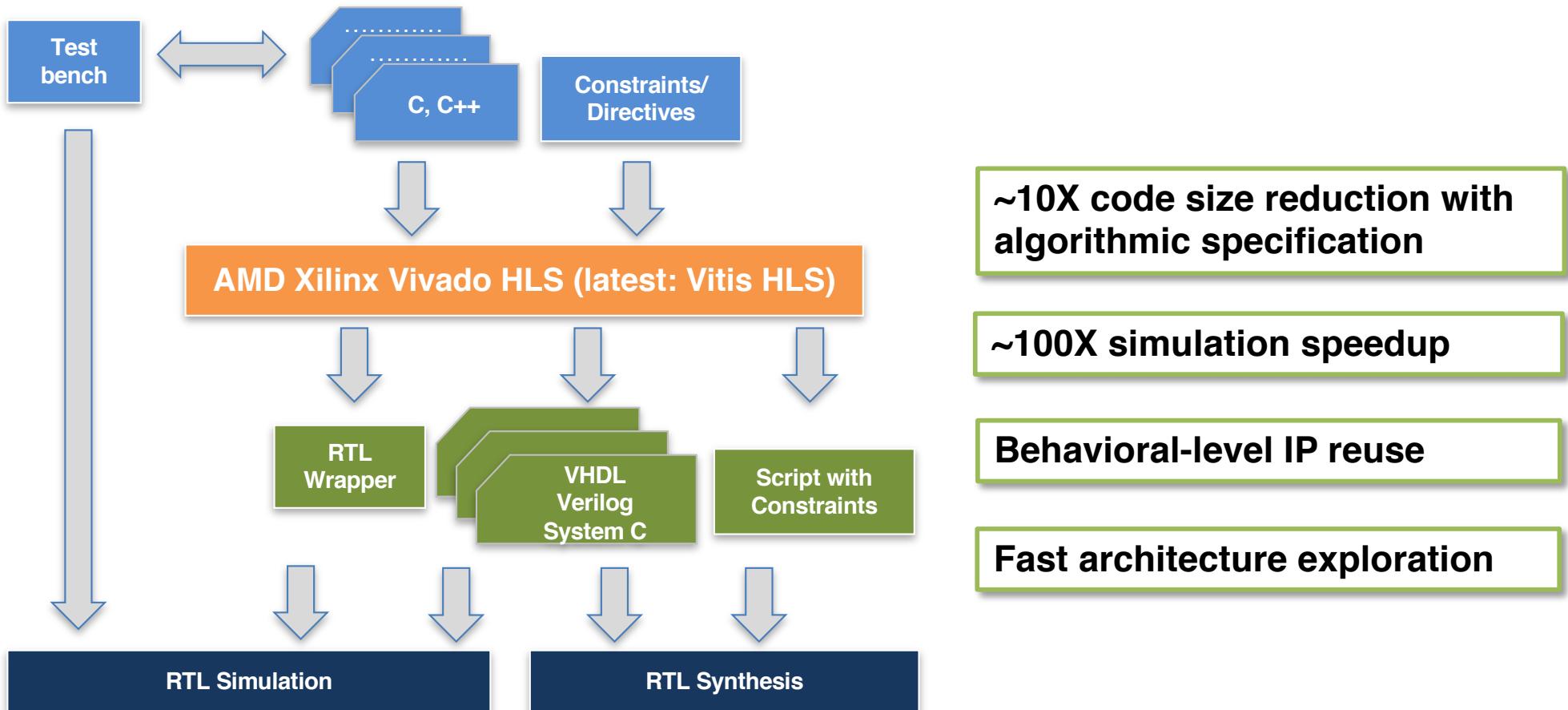
    else
        if ((current_state_1 == STATE_DONE))
            req_rdy = 0; resp_val = 1;
            a_mux_sel = A_MUX_SEL_X; b_mux_sel = B_MUX_SEL_X;
            a_reg_en = 0; b_reg_en = 0;
end
```

State transition

Single Untimed Source to Multiple Implementations

Untimed	Combinational for Latency	Sequential for Area	Pipelined for Throughput
 <p>Control-Data Flow Graph (CDFG)</p> $out1 = f(in1, in2, in3, in4)$			
	$t_{clk} \approx 3 * d_{add}$ $T_1 = 1 / t_{clk}$ $A_1 = 3 * A_{add}$	$t_{clk} \approx d_{add} + d_{setup}$ $T_2 = 1 / (3 * t_{clk})$ $A_2 = A_{add} + 2 * A_{reg}$	$t_{clk} \approx d_{add} + d_{setup}$ $T_3 = 1 / t_{clk}$ $A_3 = 3 * A_{add} + 6 * A_{reg}$

A Representative C-Based HLS Tool



Typical C/C++ Synthesizable Subset

- ▶ Data types:
 - **Primitive types:** (u)char, (u)short , (u)int, (u)long, float, double
 - **Arbitrary precision** integer or fixed-point types
 - **Composite types:** array, struct, class
 - **Templated types:** template< >
 - **Statically determinable pointers**
- ▶ No (or very limited) support for dynamic memory allocations and recursive function calls

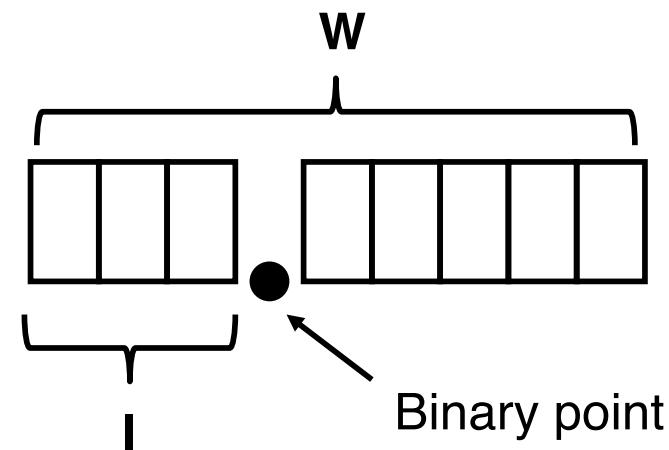
Arbitrary Precision Integer

- ▶ C/C++ only provides a limited set of native integer types
 - char (8b), short (16b), int (32b), long (??), long long (64b)
 - Byte aligned: efficient in processors
- ▶ Arbitrary precision integer in Vivado HLS
 - Signed: **ap_int**; Unsigned **ap_uint**
 - Two's complement representation for signed integer
 - Templatized class **ap_int<W>** or **ap_uint<W>**
 - W is the user-specified bitwidth

```
#include <ap_int>
...
ap_int<9>    x; // 9-bit
ap_uint<24>   y; // 24-bit unsigned
ap_uint<512>  z; // 512-bit unsigned
```

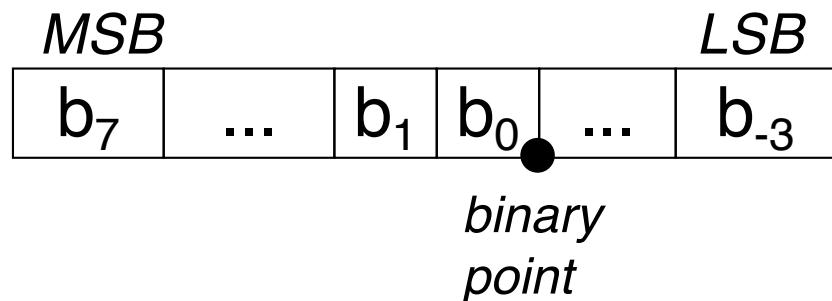
Fixed-Point Type in Vivado HLS

- ▶ Arbitrary precision fixed-point type
 - Signed: **ap_fixed**; Unsigned **ap_ufixed**
 - Templatized class **ap_fixed<W, I, Q, O>**
 - W: total bitwidth
 - I: integer bitwidth
 - Q: quantization mode
 - O: overflow mode



Example: Fixed-Point Modeling

- ▶ `ap_ufixed<11, 8, AP_TRN, AP_WRAP> x;`



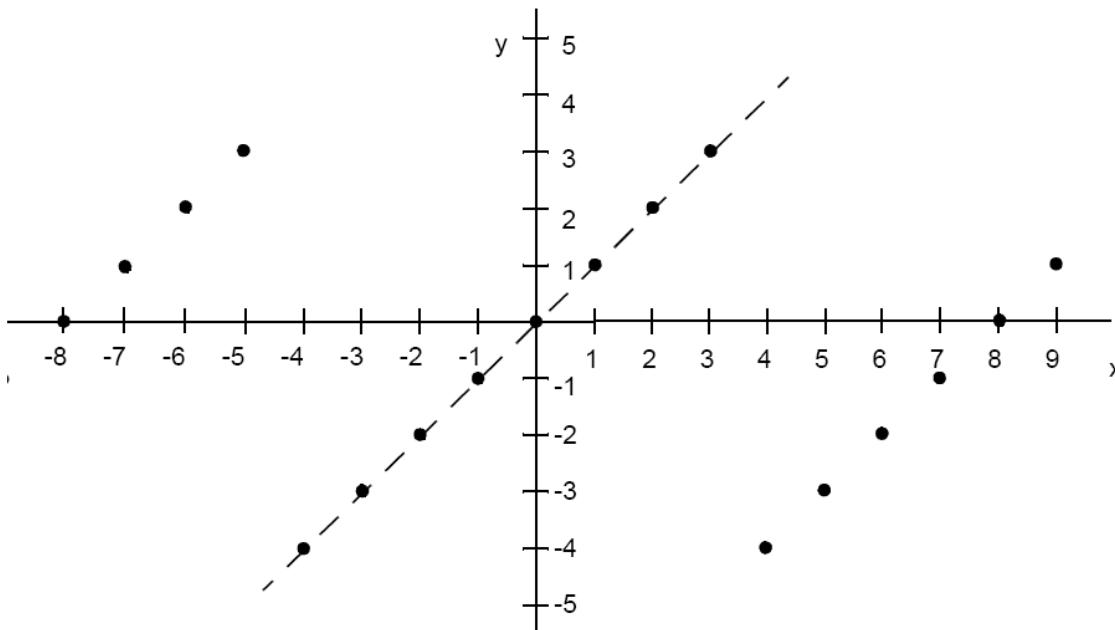
- 11 is the total number of bits in the type
- 8 bits to the left of the decimal point
- AP_WRAP defines wrapping behavior for overflow
- AP_TRN defines truncation behavior for quantization

Fixed-Point Type: Overflow Behavior

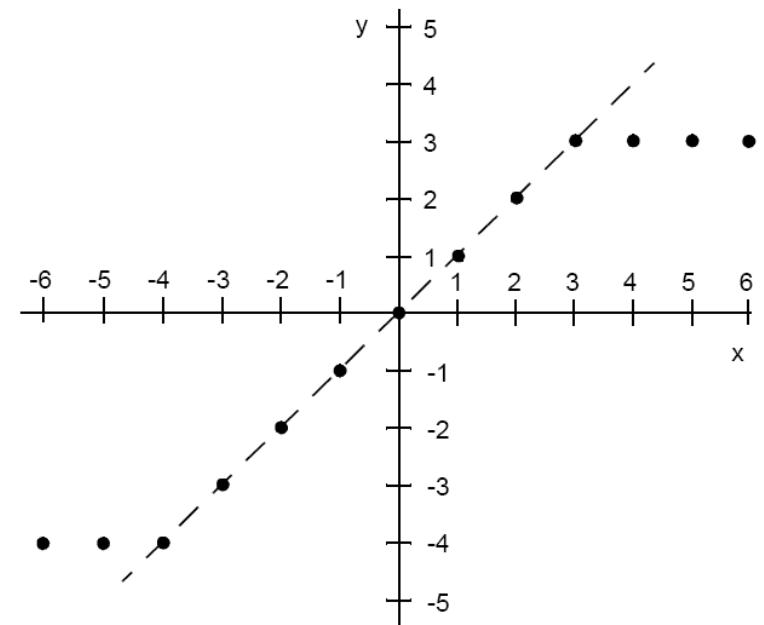
► ap_fixed overflow mode

- Determines the behavior of the fixed-point type when the result of an operation generates more precision in the **MSBs** than is available

```
ap_fixed<W, IW_X> x;  
ap_fixed<W, IW_Y> y = x; // assuming IW_Y < IW_X
```



Default: AP_WRAP (wrapping mode)



AP_SAT (saturation mode)

Fixed-Point Type: Quantization Behavior

► ap_fixed quantization mode

- Determines the behavior of the fixed-point type when the result of an operation generates more precision in the **LSBs** than is available
- Default mode: AP_TRN (truncation)
- Other rounding modes: AP_RND, AP_RND_ZERO, AP_RND_INF, ...

ap_fixed<4, 2, AP_TRN> x = 1.25; (b'01.01)

ap_fixed<3, 2, AP_TRN> y = x;
 └→ 1.0 (b'01.0)

ap_fixed<4, 2, AP_TRN> x = -1.25; (b'10.11)

ap_fixed<3, 2, AP_TRN> y = x;
 └→ -1.5 (b'10.1)

Typical C/C++ Constructs to RTL Mapping

<u>C/C++</u> Constructs	→	<u>RTL</u> Components
Functions	→	Modules
Arguments	→	Input/output ports
Operators	→	Functional units
Scalars	→	Wires or registers
Arrays	→	Memories
Control flows	→	Control logics

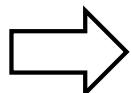
Functions and Design Hierarchy

- ▶ Each function is usually translated into an RTL module
 - Function arguments become ports on the RTL blocks
 - Functions may be inlined to dissolve their hierarchy

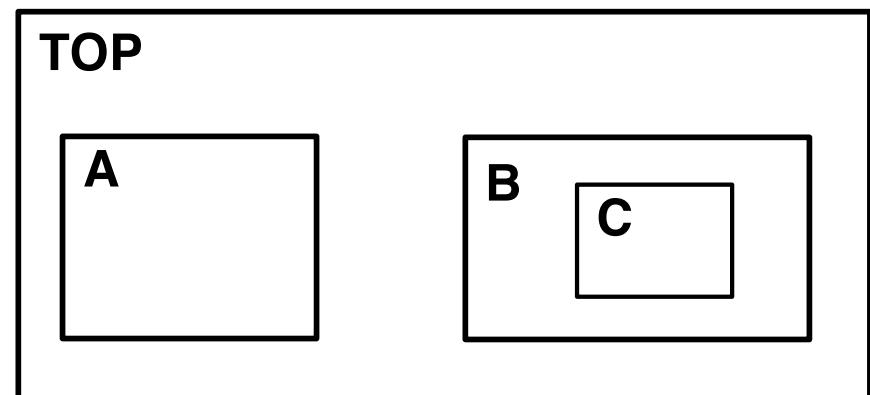
Source code

```
void A() { ... /* body of A */ }
void C() { ... /* body of C */ }
void B() {
    C();
}

void TOP(int x, int* y) {
    A(); B();
}
```



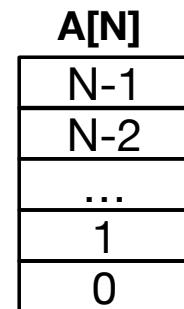
RTL module hierarchy



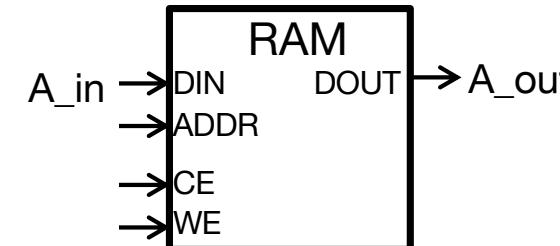
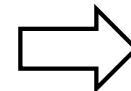
Arrays

- ▶ By default, an array in C code is implemented by a memory module in RTL
 - Read & write array -> RAM; Constant array -> ROM

```
void TOP() {  
    int A[N];  
    for (i = 0; i < N; i++)  
        A[i] = A[i] + i;  
}
```



Array



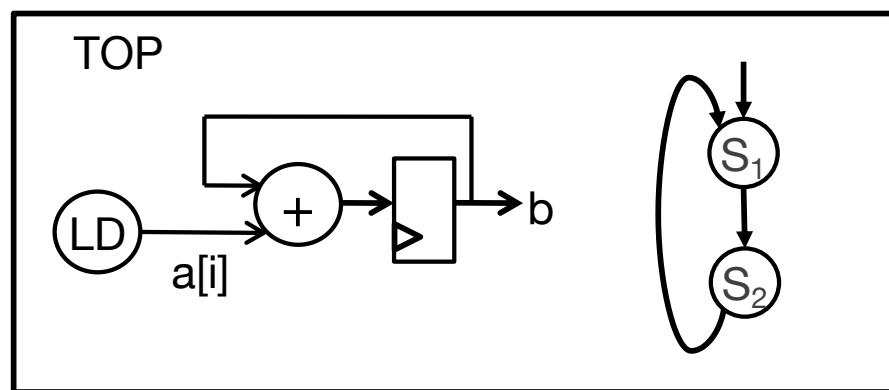
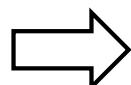
Memory Module (in SRAM)

- ▶ An array can be partitioned and implemented with multiple RAMs
 - Extreme case: completely partitioned into individual elements that map to discrete registers
- ▶ Multiple arrays can be merged and mapped to one RAM

Loops

- ▶ By default, loops are “rolled”
 - Each loop iteration corresponds to a “sequence” of states (more generally, an FSM)
 - This state sequence will be repeated multiple times based on the loop trip count (or loop bound)

```
void TOP() {  
    ...  
    for (i = 0; i < N; i++)  
        sum += A[i];  
}
```

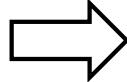


Loop Unrolling

- ▶ Unrolling can expose more parallelism to achieve shorter latency or higher throughput
 - (+) Decreased loop control overhead
 - (+) Increased parallelism for scheduling
 - (-) Increased operation count, which may negatively impact area, timing, and power

```
for (int i = 0; i < 8; i++)  
    A[i] = C[i] + D[i];
```

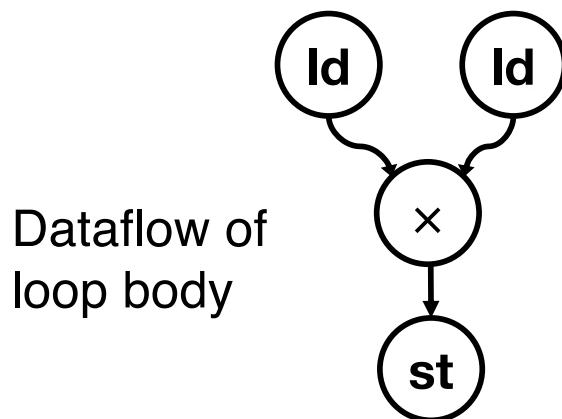
Unrolling


$$\begin{aligned} A[0] &= C[0] + D[0]; \\ A[1] &= C[1] + D[1]; \\ A[2] &= C[2] + D[2]; \\ &\dots \\ A[7] &= C[7] + D[7]; \end{aligned}$$

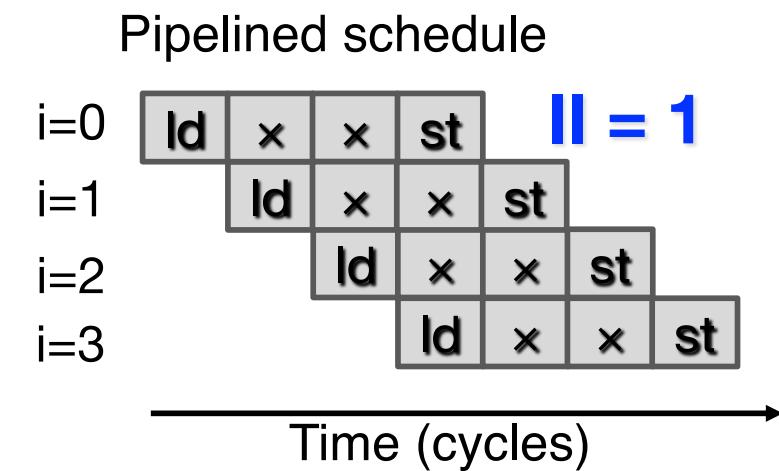
Loop Pipelining

- ▶ Pipelining is one of the most important optimizations for HLS
 - Key factor: **Initiation Interval (II)**
 - Allows a new iteration to begin processing, II cycles after the start of the previous iteration (II=1 means the loop is fully pipelined)

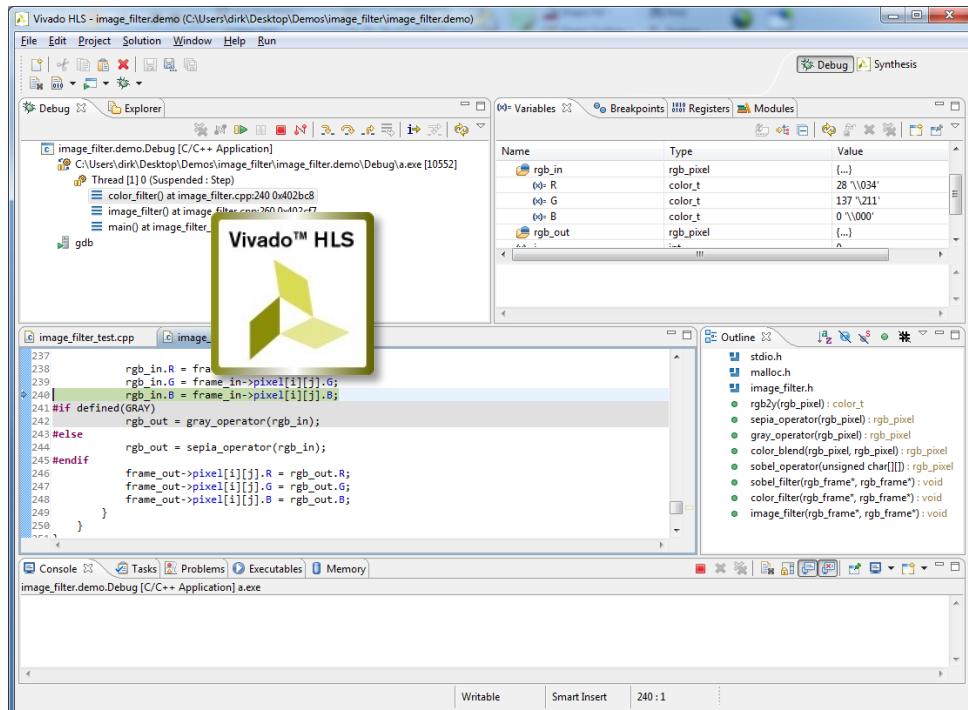
```
for (i = 0; i < N; ++i)
    p[i] = x[i] * y[i];
```



Id – Load (memory read)
st – Store (memory write)



A Tutorial on Vivado HLS



```
(vivado19) nz264@brg-zhang-xcel:~/shared/ece5997/mvmul-tutorial$ vivado_hls -f run.tcl
***** Vivado(TM) HLS - High-Level Synthesis from C, C++ and SystemC v2019.2.1 (64-bit)
**** IP Build 2729494 on Thu Dec 5 07:38:25 MST 2019
** Copyright 1986-2019 Xilinx, Inc. All Rights Reserved.

source /opt/xilinx/Xilinx_Vivado_vitis_2019.2/Vivado/2019.2/scripts/vivado_hls/hls.tcl -notrace
INFO: Applying HLS Y2K22 patch v1.2 for IP revision
INFO: [HLS 200-10] Running '/opt/xilinx/Xilinx_Vivado_vitis_2019.2/Vivado/2019.2/bin/unwrapped/lnx64.o/vivado_hls'
INFO: [HLS 200-10] For user 'nz264' on host 'en-ec-brg-zhang-xcel.coecis.cornell.edu' (Linux_x86_64 version 3.10.0-1160.71.1.e17.x86_64) on Mon Aug 22 11:07:33 EDT 2022
INFO: [HLS 200-10] On os "CentOS Linux release 7.9.2009 (Core)"
INFO: [HLS 200-10] In directory '/work/shared/users/phd/nz264/ece5997/mvmul-tutorial'
Sourcing Tcl script 'run.tcl'
INFO: [HLS 200-10] Opening and resetting project '/work/shared/users/phd/nz264/ece5997/mvmul-tutorial/mvmul_vitis.prj'.
INFO: [HLS 200-10] Adding design file 'mvmul_unroll.c' to the project
INFO: [HLS 200-10] Adding test bench file 'mvmul_top.c' to the project
INFO: [HLS 200-10] Opening and resetting solution '/work/shared/users/phd/nz264/ece5997/mvmul-tutorial/mvmul_vitis.prj/solution1'.
INFO: [HLS 200-10] Cleaning up the solution database.
INFO: [HLS 200-10] Setting target device to 'xc7z020-clg484-1'
INFO: [SYN 201-201] Setting up clock 'default' with a period of 10ns.
INFO: [SCHED 204-61] Option 'relax_ihi_for_timing' is enabled, will increase II to preserve clock frequency constraints.
INFO: [HLS 200-10] Analyzing design file 'mvmul_unroll.c' ...
INFO: [HLS 200-111] Finished Linking Time (s): cpu = 00:00:11 ; elapsed = 00:00:18 . Memory (MB): peak = 1057.715 ; gain = 527.219 ; free physical = 97063 ; free virtual = 219377
INFO: [HLS 200-111] Finished Checking Pragmas Time (s): cpu = 00:00:11 ; elapsed = 00:00:18 . Memory (MB): peak = 1057.715 ; gain = 527.219 ; free physical = 97063 ; free virtual = 219377
INFO: [HLS 200-10] Starting code transformations ...
INFO: [HLS 200-111] Finished Standard Transformations Time (s): cpu = 00:00:12 ; elapsed = 00:00:19 . Memory (MB): peak = 1057.715 ; gain = 527.219 ; free physical = 97048 ; free virtual = 219377
INFO: [HLS 200-10] Checking synthesizability ...
INFO: [HLS 200-111] Finished Checking Synthesizability Time (s): cpu = 00:00:12 ; elapsed = 00:00:19 . Memory (MB): peak = 1057.715 ; gain = 527.219 ; free physical = 97063 ; free virtual = 219377
INFO: [HLS 200-489] Unrolling loop 'ACC_LOOP' ('mvmul_unroll.c:17') in function 'mvmul' completely with a factor of 16.
INFO: [XFROM 293-111] Balancing expressions in function 'mvmul' ('mvmul_unroll.c:6)...15 expression(s) balanced.
INFO: [HLS 200-111] Finished Pre-synthesis Time (s): cpu = 00:00:12 ; elapsed = 00:00:19 . Memory (MB): peak = 1057.715 ; gain = 527.219 ; free physical = 97044 ; free virtual = 219358
INFO: [HLS 200-111] Finished Architecture Synthesis Time (s): cpu = 00:00:12 ; elapsed = 00:00:19 . Memory (MB): peak = 1057.715 ; gain = 527.219 ; free physical = 97043 ; free virtual = 219357
INFO: [HLS 200-10] Starting hardware synthesis ...
INFO: [HLS 200-10] Synthesizing 'mvmul' ...
WARNING: [SYN 201-107] Renaming port name 'mvmul/output' to 'mvmul/output_r' to avoid the conflict with HDL keywords or other object names.
INFO: [HLS 200-10] --
INFO: [HLS 200-42] -- Implementing module 'mvmul'
INFO: [HLS 200-10] --
INFO: [SCHED 204-11] Starting scheduling ...
INFO: [SCHED 204-11] Finished scheduling.
```

AMD Xilinx Vivado HLS (v2019.2)

Matrix-Vector Multiplication

$$\mathbf{y} = A \mathbf{x} = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1N} \\ a_{21} & a_{22} & \cdots & a_{2N} \\ \vdots & \vdots & \ddots & \vdots \\ a_{N1} & a_{N2} & \cdots & a_{NN} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} a_{11}x_1 + a_{12}x_2 + \cdots + a_{1N}x_N \\ a_{21}x_1 + a_{22}x_2 + \cdots + a_{2N}x_N \\ \vdots \\ a_{N1}x_1 + a_{N2}x_2 + \cdots + a_{NN}x_N \end{bmatrix}$$

A: input matrix
x : input vector
y : output vector

```
// original, non-optimized version of MV

#define N 16 // dimension of matrix/vector

void MV ( int A[N][N], int x[N], int y[N] ) {
    // iterate over rows of matrix A
    for ( int i = 0; i < N; i++ ) {
        int acc = 0;
        // inner product
        for ( int j = 0; j < N; j++ ) {
            acc += A[i][j] * x[j];
        }
        y[i] = acc;
    }
}
```

Setup on ECE Linux Server

- ▶ Log into ecelinux server
 - Host name: ecelinux.ece.cornell.edu
 - User name and password: [Your NetID credentials]
- ▶ Setup tools for this class
 - Source class setup script to setup Vivado HLS

```
> source /classes/ece5775/setup-ece5775.sh
```
- ▶ Test Vivado HLS
 - Open Vivado HLS interactive environment

```
> vivado_hls -i
```
 - List the available commands

```
> help
```

Copy MV Example to Your Home Directory

```
> cd ~  
> cp -r /classes/ece5775/mv-tutorial/ .  
> ls
```

- ▶ Design files
 - mv.h: function prototypes
 - mv_*.cpp: function definitions
- ▶ Testbench files
 - testbench.cpp: function used to test the design
- ▶ Synthesis configuration files
 - run.tcl: script for configuring and running Vivado HLS

Project Tcl Script

```
#=====
# run.tcl for matrix-vector (MV) multiplication
#=====

# open the HLS project mv.prj
open_project -reset mv.prj

# set the top-level function of the design to mv
set_top mv

# add design and testbench files
add_files mv_initial.cpp
add_files -tb testbench.cpp

open_solution "solution1"

# use Zynq device
set_part xc7z020clg484-1

# target clock period is 10 ns
create_clock -period 10
```

```
# do a c simulation
csim_design
```

```
# synthesize the design
csynth_design
```

```
# do a co-simulation
cosim_design
```

```
# export design
export_design
```

```
# exit Vivado HLS
exit
```

You can use multiple Tcl scripts to automate different runs with different configurations.

Synthesize and Simulate the Design

```
> vivado_hls -f run.tcl
```

```
Generating csim.exe  
128/128 correct values!  
INFO: [SIM 211-1] CSim done with 0 errors.
```

```
INFO: [HLS 200-10]  
INFO: [HLS 200-10] -- Scheduling module 'mv'  
INFO: [HLS 200-10]
```

```
INFO: [HLS 200-10]  
INFO: [HLS 200-10] -- Exploring micro-arch for module 'mv'  
INFO: [HLS 200-10]
```

```
INFO: [HLS 200-10]  
INFO: [HLS 200-10] -- Generating RTL for module 'mv'  
INFO: [HLS 200-10]
```

```
INFO: [COSIM 212-47] Using XSIM for RTL simulation.  
INFO: [COSIM 212-14] Instrumenting C test bench ...
```

```
INFO: [COSIM 212-12] Generating RTL test bench ...  
INFO: [COSIM 212-323] Starting verilog simulation.  
INFO: [COSIM 212-15] Starting XSIM ...
```

```
INFO: [COSIM 212-316] Starting C post checking ...  
128/128 correct values!
```

```
INFO: [COSIM 212-1000] *** C/RTL co-simulation finished: PASS ***
```

SW simulation only

Same as simply running a software program

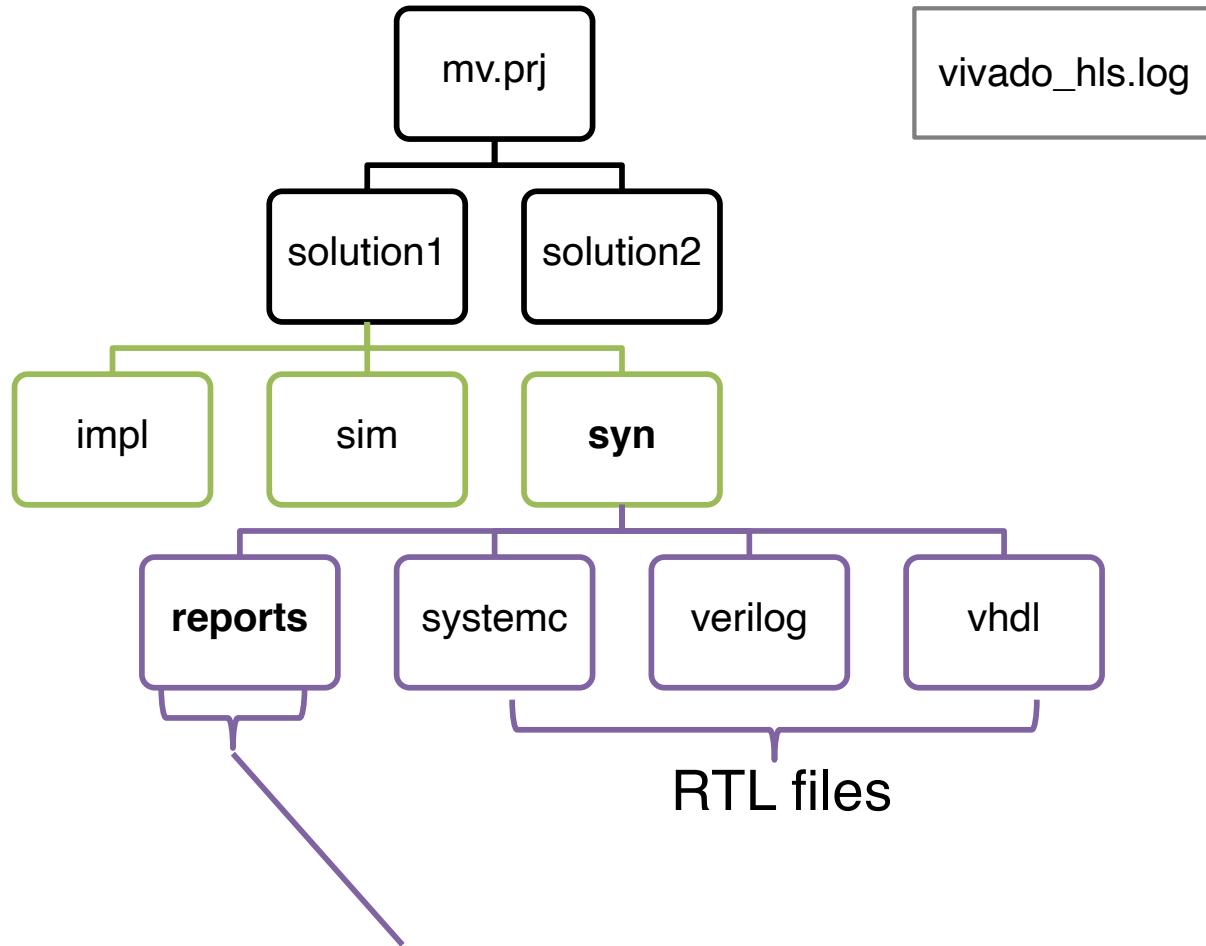
Synthesis

Compiling C to RTL

HW-SW co-simulation

RTL simulation driven by SW test bench

Synthesis Directory Structure



Synthesis reports of each function
in the design, except those inlined.

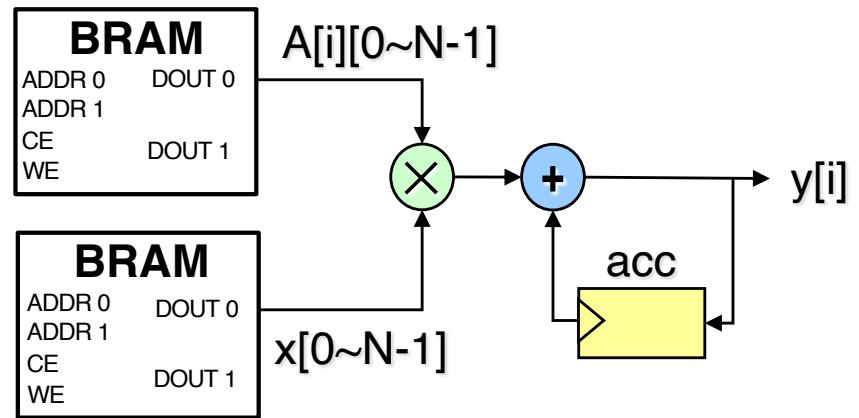
Default Microarchitecture

```
// original, non-optimized version of MV

#define N 16 // dimension of matrix/vector

void MV ( int A[N][N], int x[N], int y[N] ) {
    // iterate over rows of matrix A
    for ( int i = 0; i < N; i++ ) {
        int acc = 0;
        // inner product
        for ( int j = 0; j < N; j++ ) {
            acc += A[i][j] * x[j];
        }
        y[i] = acc;
    }
}
```

$$y[i] = \sum_{j=0}^N A[i][j] \times x[j]$$



Assumption: A and x are stored in two discrete **block RAMs** (**BRAMs**); each has 2 read/write ports

Latency: ??

Examine Overall Latency

```
// original, non-optimized version of MV

#define N 16 // dimension of matrix/vector

void MV ( int A[N][N], int x[N], int y[N] ) {
    // iterate over rows of matrix A
    for ( int i = 0; i < N; i++ ) {
        int acc = 0;
        // inner product
        for ( int j = 0; j < N; j++ ) {
            acc += A[i][j] * x[j];
        }
        y[i] = acc;
    }
}
```

4 cycles

$4 \times 16 + 2 = 66 \text{ cycles}$

Latency: $\sim 1056 = 66 \times 16$

Possible optimizations

- Unrolling
- Pipelining
- Array partitioning

Unroll Inner Loop

```
// MV with inner loop unroll
#define N 16

void MV ( int A[N][N], int x[N], int y[N] ) {
    // iterate over rows of matrix A
    for (int i = 0; i < N; i++) {
        int acc = 0;
        // inner product
        for (int j = 0; j < N; j++) {
            #pragma HLS unroll
            acc += A[i][j] * x[j];
        }
        y[i] = acc;
    }
}
```

Complete
unrolling

```
// unrolled multiply-accumulate
acc += A[i][0] * x[0];
acc += A[i][1] * x[1];
acc += A[i][2] * x[2];
...
acc += A[i][15] * x[15];
```

Latency: ??

Overall Latency after Unrolling

```
// MV with inner loop unroll
#define N 16

void MV ( int A[N][N], int x[N], int y[N] ) {
    // iterate over rows of matrix A
    for (int i = 0; i < N; i++) {
        int acc = 0;
        // inner product
        acc += A[i][0] * x[0];
        acc += A[i][1] * x[1];
        acc += A[i][2] * x[2];
        ...
        acc += A[i][15] * x[15];
        y[i] = acc;
    }
}
```

12 cycles

Latency: $\sim 192 = 12 \times N = 12 \times 16$

Pipeline Outer Loop

```
// MV with outer loop pipeline
#define N 16

void MV ( int A[N][N], int x[N], int y[N] ) {
    // iterate over rows of matrix A
    for (int i = 0; i < N; i++) {
        #pragma HLS pipeline
        int acc = 0;
        // inner product
        for (int j = 0; j < N; j++ ) {
            #pragma HLS unroll
            acc += A[i][j] * x[j];
        }
        output[i] = acc;
    }
}
```

Pipeline the outer loop

Inner loops automatically
unrolled when pipelining
the outer loop

|| = ??

Overall Latency after Pipelining

```
// MV with outer loop pipeline
#define N 16

void MV ( int A[N][N], int x[N], int y[N] ) {
    // iterate over rows of matrix A
    for (int i = 0; i < N; i++) {
        #pragma HLS pipeline
        int acc = 0;
        // inner product
        acc += A[i][0] * x[0];
        acc += A[i][1] * x[1];
        acc += A[i][2] * x[2];
        ...
        acc += A[i][15] * x[15];
        y[i] = acc;
    }
}
```

12 cycles

Target II = 1
Achieved II = 8

Pipeline rate limited
by BRAM ports
16 reads on A (or x),
but only two ports
available

$$\text{Latency: } \sim 132 = 12 + (N - 1) \times 8 = 12 + 15 \times 8$$

Partition Arrays

```
// MV with outer loop pipeline and array partition
#define N 16
void MV(int A[N][N], int x[N], int y[N]) {
    #pragma HLS array_partition variable=A dim=2
    #pragma HLS array_partition variable=x
    // iterate over rows of matrix A
    for (int i = 0; i < N; i++) {
        #pragma HLS pipeline
        // inner product
        for (int j = 0; j < N; j++) {
            ...
        }
    }
}
```

Latency: ??

// matrix col arrays

```
int A_col_0[16];
int A_col_1[16];
int A_col_2[16];
...
int A_col_15[16];
```

// vector elements

```
int x_reg_0;
int x_reg_1;
int x_reg_2;
...
int x_reg_15;
```

Partition along second axis
(column)

Complete array partitioning

Overall Latency after Partitioning

```
// MV with outer loop pipeline
#define N 16

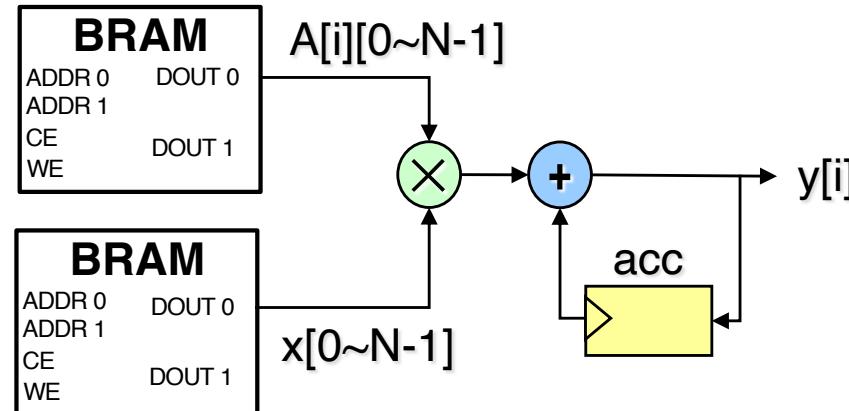
void MV ( int A[N][N], int x[N], int y[N] ) {
    #pragma HLS array_partition variable=A dim=2
    #pragma HLS array_partition variable=x
    // iterate over rows of matrix A
    for (int i = 0; i < N; i++) {
        #pragma HLS pipeline
        int acc = 0;
        // inner product
        acc += A[i][0] * x[0];
        acc += A[i][1] * x[1];
        acc += A[i][2] * x[2];
        ...
        acc += A[i][15] * x[15];
        y[i] = acc;
    }
}
```

6 cycles

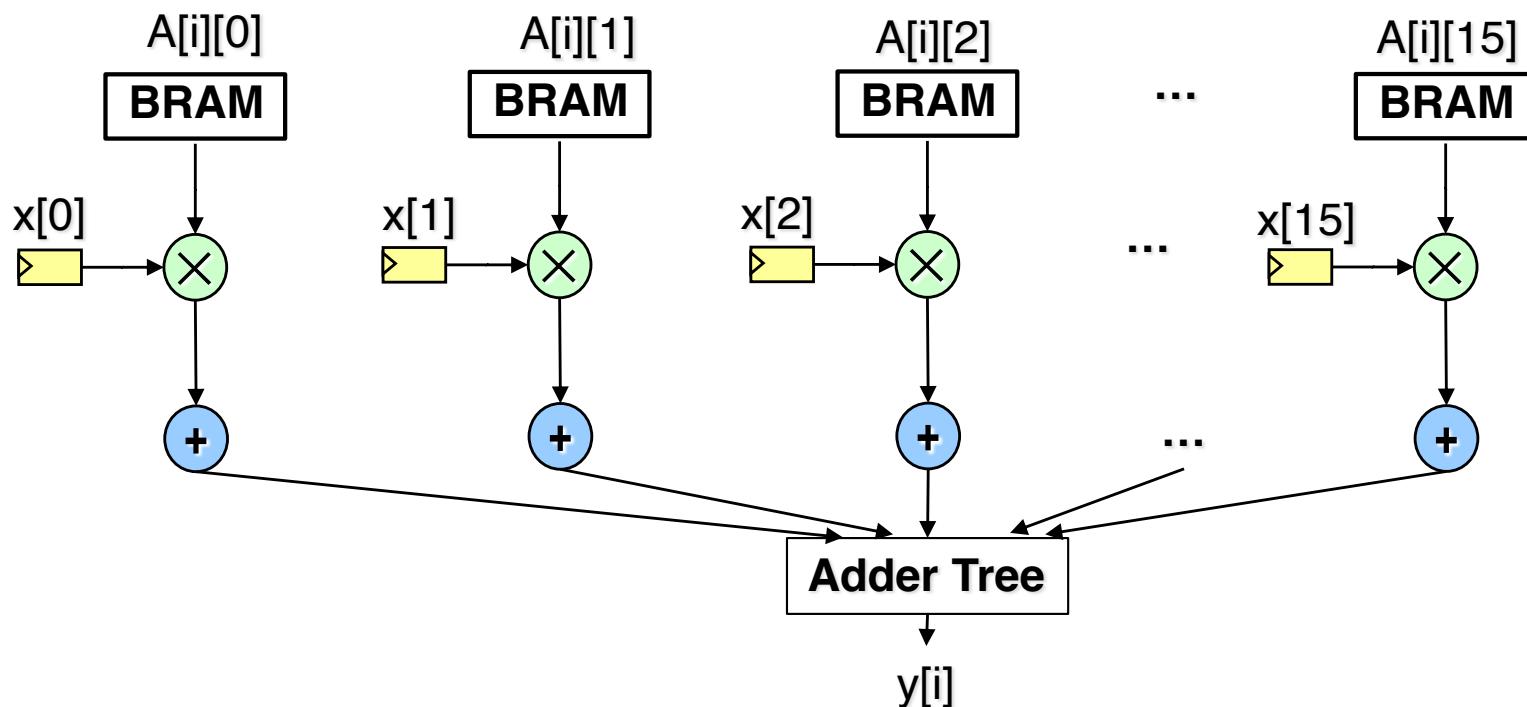
Target II = 1
Achieved II = 1

$$\text{Latency: } \sim 21 = 6 + (16 - 1) \times 1$$

Microarchitecture of the Optimized Design

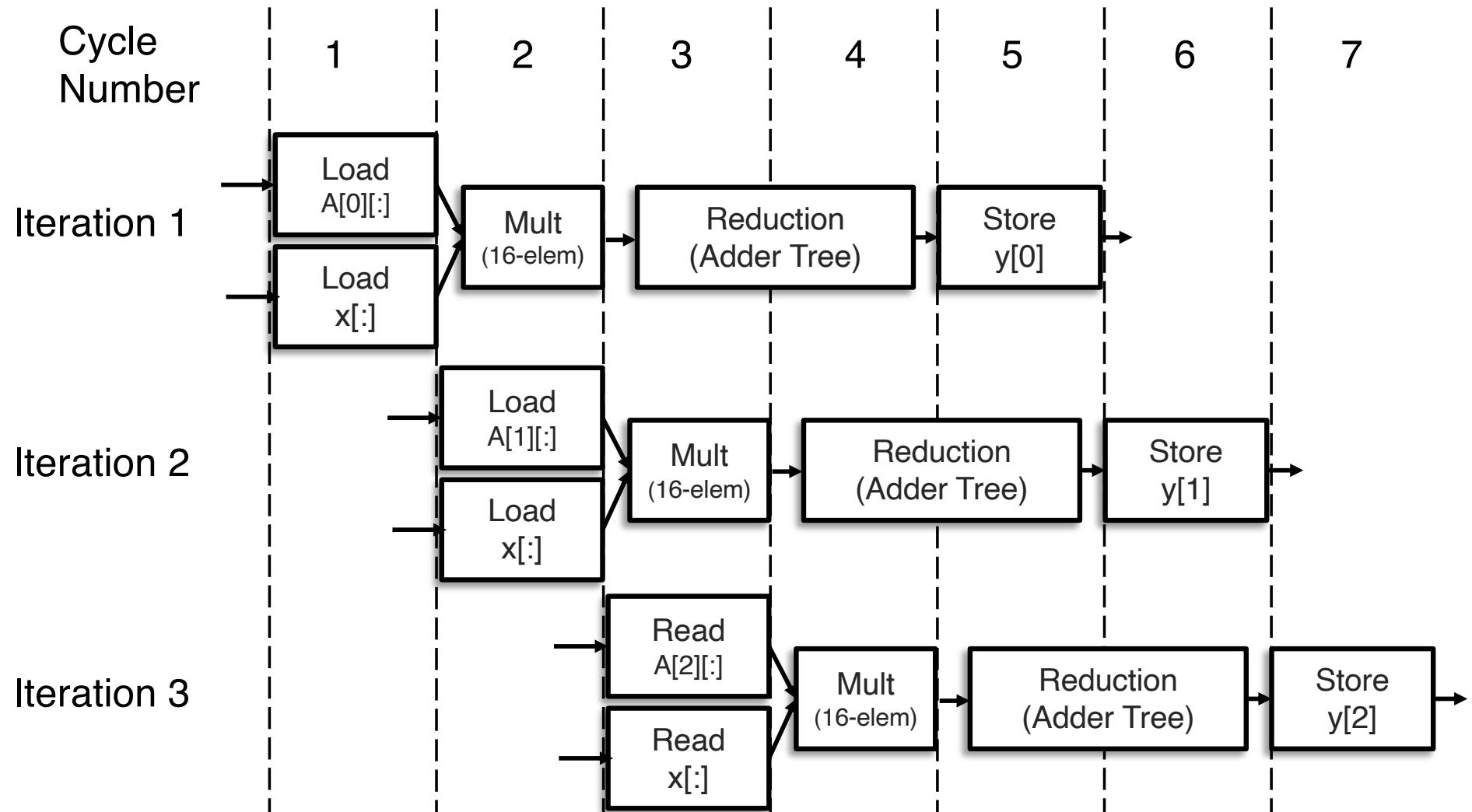


Default



Unrolled
+
Pipelined
+
Partitioned

Pipeline Schedule of the Optimized Design



[Figure credit] Ryan Kastner et. al., *Parallel Programming for FPGAs*, page 44.

Next Lecture

- ▶ Field-programmable gate arrays (FPGAs)