Introduction to MPLS

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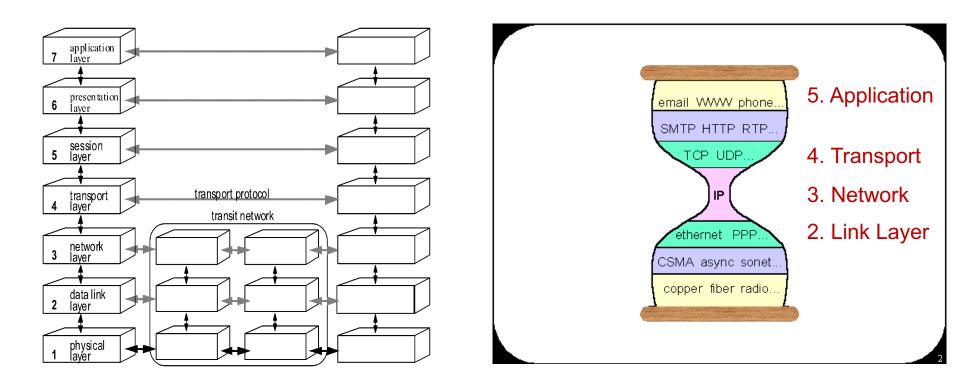
This presentation is based solely on Santanu Dasgupta's slide series that describe the Multiprotocol Label Switching (MPLS) concept Santanu Dasgupta is a Distinguished Architect at Cisco



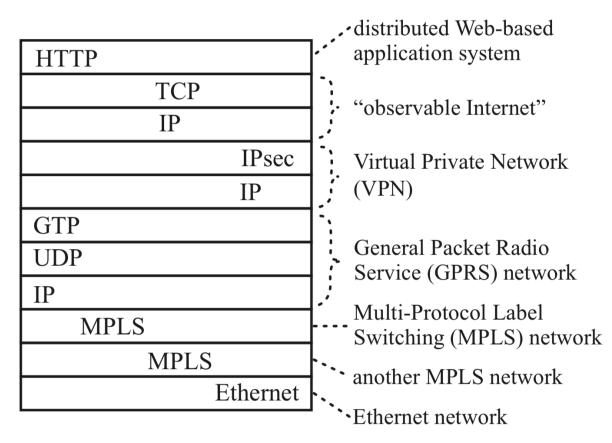


ISO/OSI model









Typical packet from AT&T backbone network

Pamela Zave and Jennifer Rexford. 2019. The compositional architecture of the internet. Commun. ACM 62, 3 (February 2019), 78-87. DOI: https://doi.org/10.1145/3226588 ·ı|ı.ı|ı. cısco

Introduction to MPLS

Santanu Dasgupta

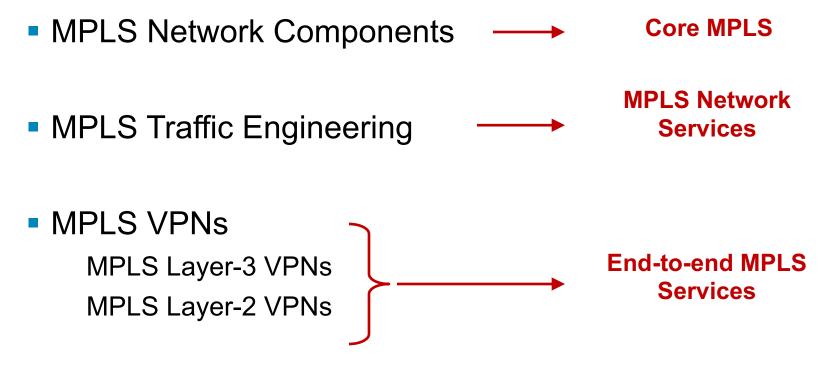
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Goals of this Lecture

- Understand the business drivers for MPLS
- Learn about MPLS customer and market segments
- Understand the problems MPLS is addressing
- Understand benefits of deploying MPLS
- Understand the major MPLS technology components
- Learn the basics of MPLS technology
- Understand typical applications of MPLS

Agenda

Introduction



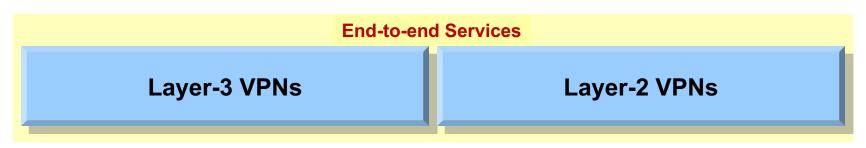
Summary

Introduction

The business drivers for MPLS



The Big Picture



	MPLS Network Services	
MPLS QoS	MPLS TE	MPLS OAM/MIBs
	Core MPLS	

MPLS Signaling and Forwarding

Network Infrastructure

Why Multi Protocol Label Switching?

SP/Carrier perspective

Reduce costs (CAPEX); consolidate networks

Consolidated network for multiple Layer-2/3 services

Support increasingly stringent SLAs

Handle increasing scale/complexity of IP-based services

Enterprise/end-user perspective

Campus/LAN

Need for network segmentation (users, applications, etc.)

WAN connectivity (connecting enterprise networks)

Need for easier configuration of site-to-site WAN connectivity

What Is MPLS Technology?

- It's all about labels ...
- Use the best of both worlds

Layer-2 (ATM/FR): efficient forwarding and traffic engineering Layer-3 (IP): flexible and scalable

MPLS forwarding plane

Use of labels for forwarding Layer-2/3 data traffic

Labeled packets are being switched instead of routed

Leverage layer-2 forwarding efficiency

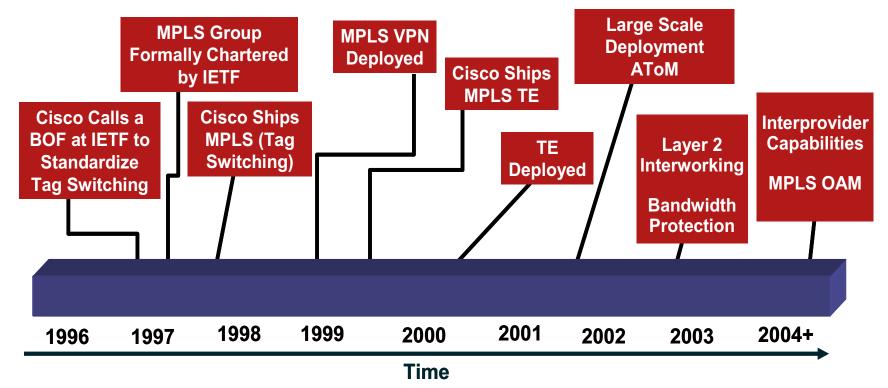
MPLS control/signaling plane

Use of existing IP control protocols extensions + new protocols to exchange label information

Leverage layer-3 control protocol flexibility and scalability

Evolution of MPLS

- Evolved from tag switching in 1996 to full IETF standard, covering over 130 RFCs
- Key application initially were Layer-3 VPNs, followed by Traffic Engineering (TE), and Layer-2 VPNs



MPLS Applications



Service Providers	Enterprise Data Center	Data center interconnects	EWAN Edge
L2/L3VPN's TE/FRR QoS High Availability	VPN's TE/FRR High Availability	VPN's / VRF's VRF-Aware Security High Availability	VPN's / VRF's VRF Aware Security High Availability
Hosted Data centers Data center interconnect Segmentation for IT Mergers, Acquisitions, spinoffs	Departmental segmentation Service multiplexing Security Mergers, Acquisitions, spinoffs	Disaster Recovery Vmotion support Branch Interconnects	Internet Access Branch Connectivity

- Network Consolidation Merging Multiple parallel network into a shared infrastructure
- Network segmentation By user groups or business function
- Service and policy centralization Security policies and appliances at a central location
- New applications readiness Converged multi-service network
- Increased network security User groups segmentation with VPNs

Enterprise MPLS Customers

- Two types of enterprise customers for MPLS technology
- MPLS indirectly used as subscribed WAN service Enterprise subscribes to WAN connectivity data service offered by external Service Provider
 Data connectivity service implemented by Service Provider via MPLS VPN technology (e.g., layer-2 and layer-3 VPNs)
 VPN Service can be managed or unmanaged
- MPLS used as part of self managed network
 Enterprise deploys MPLS in it's own network
 Enterprise manages it's own MPLS-based network

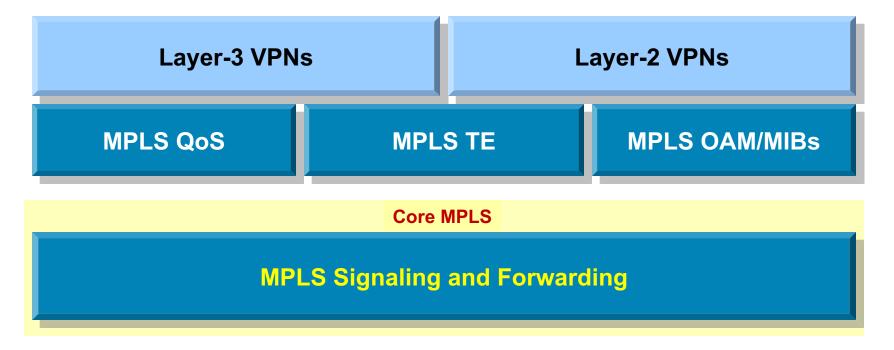
MPLS Technology Components

Basic building blocks of MPLS



MPLS Forwarding and Signaling

MPLS label forwarding and signaling mechanisms



Network Infrastructure

Basic Building Blocks

The big picture

MPLS-enabled network devices

Label Switched Paths (LSPs)

The internals

MPLS labels

Processing of MPLS labels

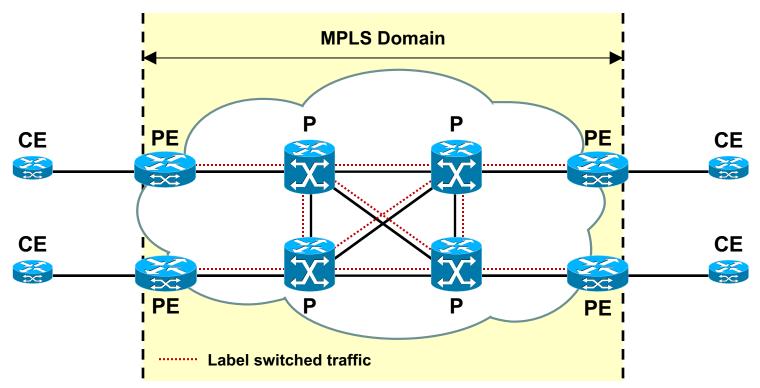
Exchange of label mapping information

Forwarding of labeled packets

 Other related protocols and protocols to exchange label information

Between MPLS-enabled devices

MPLS Network Overview



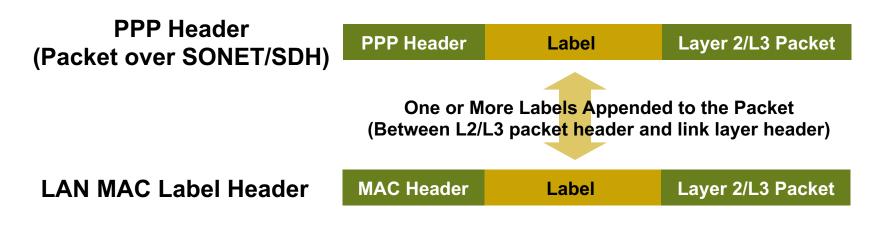
- P (Provider) router = label switching router = core router (LSR) Switches MPLS-labeled packets
- PE (Provider Edge) router = edge router (LSR) Imposes and removes MPLS labels
- CE (Customer Edge) router
 Connects customer network to MPLS network

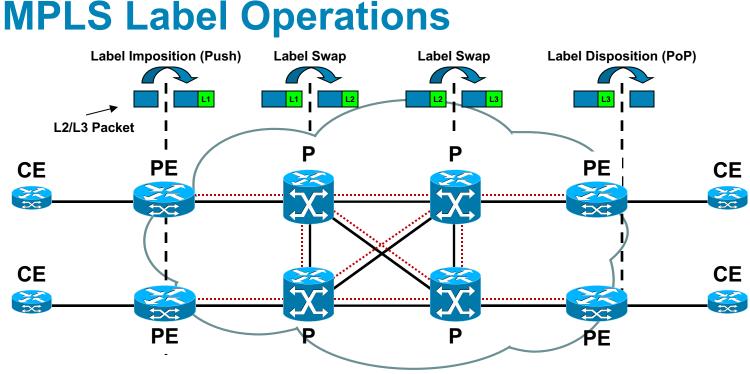
MPLS Label and Label Encapsulation MPLS Label



COS/EXP = Class of Service: 3 Bits; S = Bottom of Stack; TTL = Time to Live

MPLS Label Encapsulation





Label imposition (Push)

By ingress PE router; classify and label packets

Label swapping or switching

By P router; forward packets using labels; indicates service class & destination

Label disposition (PoP)

By egress PE router; remove label and forward original packet to destination CE

Forwarding Equivalence Class

 Mechanism to map ingress layer-2/3 packets onto a Label Switched Path (LSP) by ingress PE router

Part of label imposition (Push) operation

- Variety of FEC mappings possible
 - IP prefix/host address

Groups of addresses/sites (VPN x)

Used for L3VPNs

Layer 2 circuit ID (ATM, FR, PPP, HDLC, Ethernet)

Used for Pseudowires (L2VPNs)

Tunnel interface

Used for MPLS traffic engineering (TE)

Label Distribution Protocol

 MPLS nodes need to exchange label information with each other Ingress PE node (Push operation)

Needs to know what label to use for a given FEC to send packet to neighbor

Core P node (Swap operation)

Needs to know what label to use for swap operation for incoming labeled packets

Egress PE node (Pop operation)

Needs to tell upstream neighbor what label to use for specific FEC type LDP used for exchange of label (mapping) information

Label Distribution Protocol (LDP)

Defined in RFC 3035 and RFC3036; updated by RFC5036

LDP is a superset of the Cisco-specific Tag Distribution Protocol

 Note that, in addition LDP, also other protocols are being used for label information exchange

Will be discussed later

LDP Operations

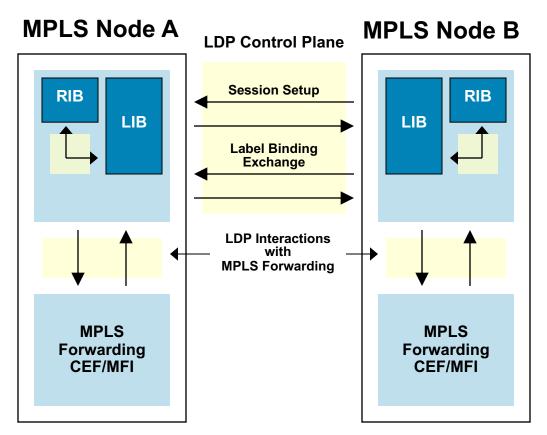
LDP startup

Local labels assigned to RIB prefixes and stored in LIB

Peer discovery and session setup

Exchange of MPLS label bindings

 Programming of MPLS forwarding
 Based on LIB info
 CEF/MFI updates



MPLS Control and Forwarding Plane

MPLS control plane

Used for distributing labels and building label-switched paths (LSPs)

Typically supported by LDP; also supported via RSVP and BGP

Labels define destination and service

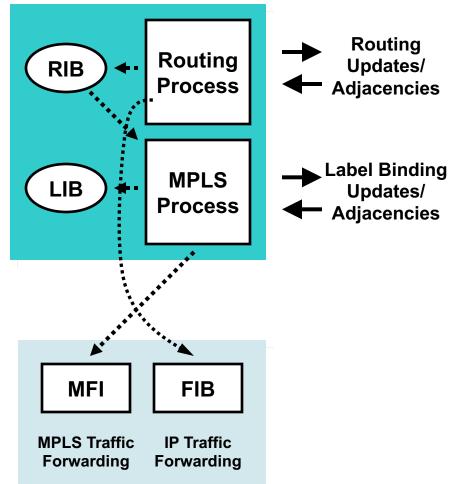
MPLS forwarding plane

Used for label imposition, swapping, and disposition

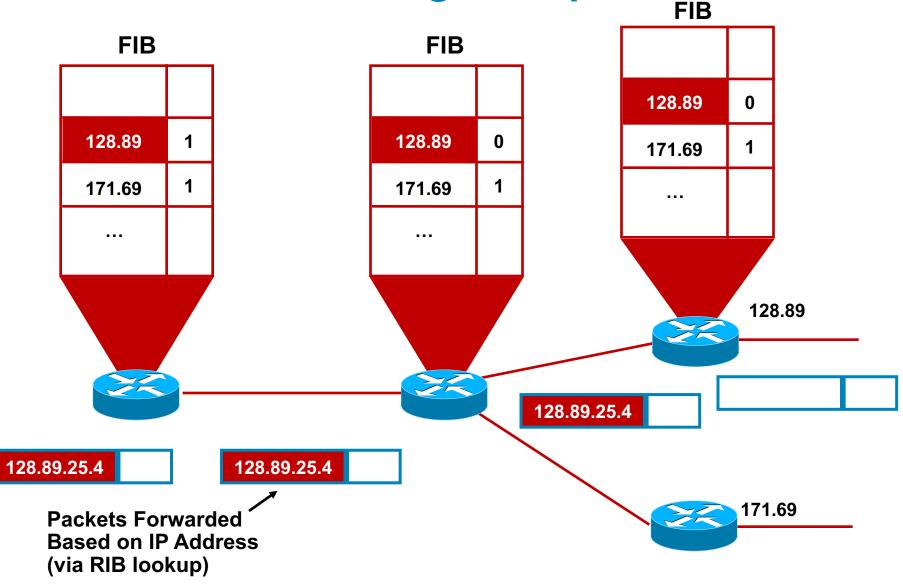
Independent of type of control plane

Labels separate forwarding from IP address-based routing

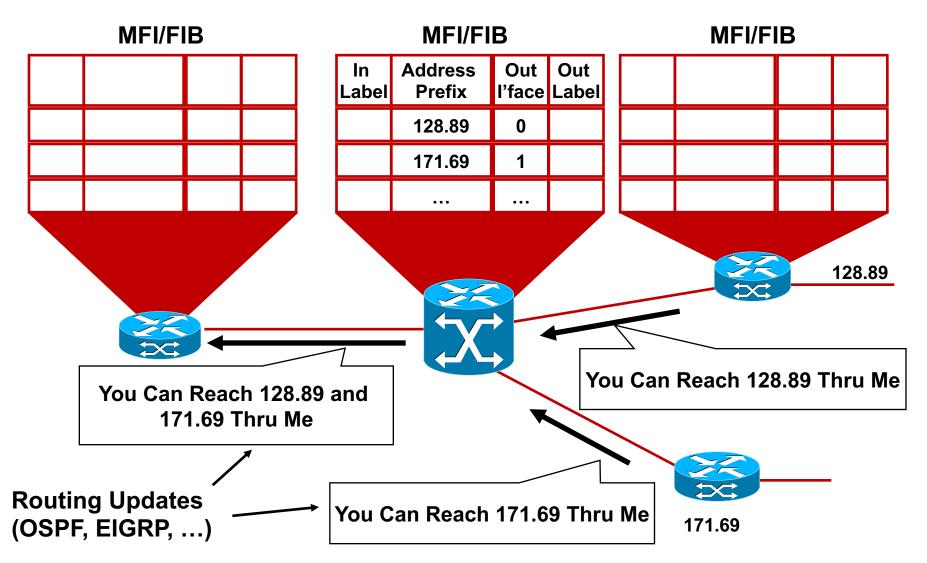
MFI – MPLS Forwarding Infrastructure



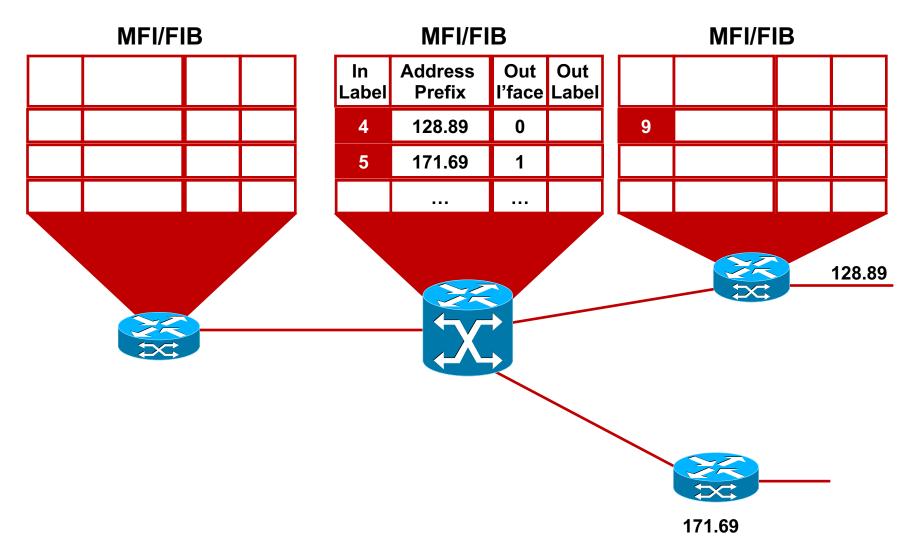
IP Packet Forwarding Example



Step 1: IP Routing (IGP) Convergence

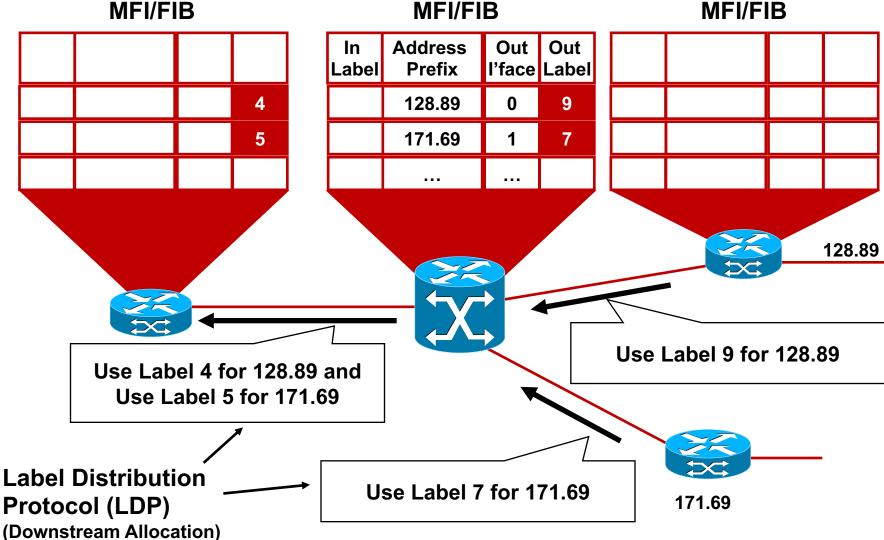


Step 2a: LDP Assigns Local Labels

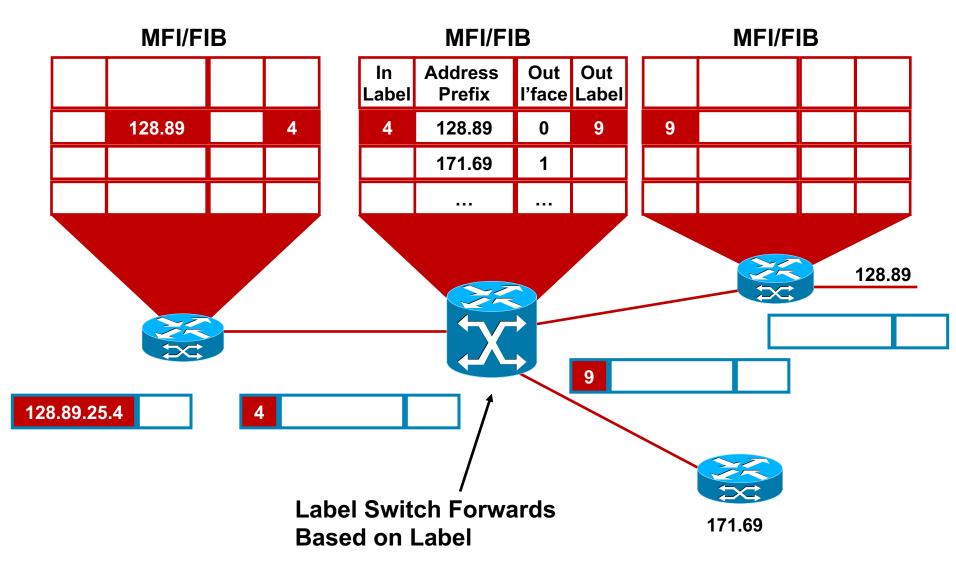


Step 2b: LDP Assigns Remote Labels

MFI/FIB



Step 3: Forwarding MPLS Packets

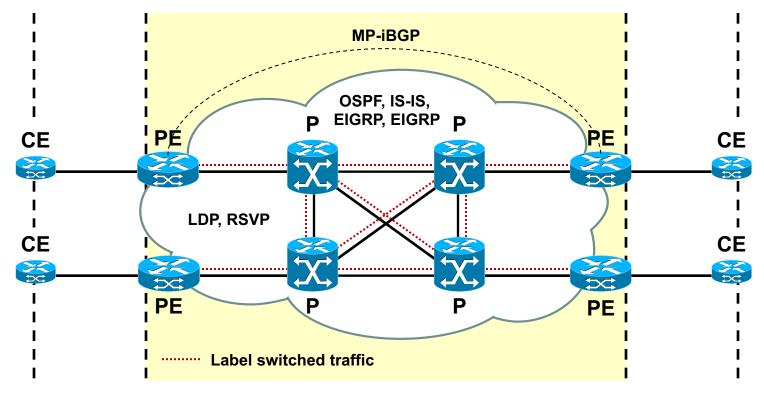


Summary Steps For MPLS Forwarding

- Each node maintains IP routing information via IGP (Interior Gateway Protocol)
 - IP routing table (RIB) and IP forwarding table (FIB)
- LDP leverages IGP routing information
- LDP label mapping exchange (between MPLS nodes) takes place after IGP has converged
 - LDP depends on IGP convergence
 - Label binding information stored in LIB
- Once LDP has received remote label binding information MPLS forwarding is updated

Label bindings are received from remote LDP peers MPLS forwarding via MFI (MPLS Forwarding Infrastructure)

MPLS Network Protocols



- IGP: OSPF, EIGRP, IS-IS on core facing and core links
- RSVP and/or LDP on core and/or core facing links
- MP-iBGP on PE devices (for MPLS services), MP-BGP: Multiprotocol Border Gateway Protocol, used for MPLS L3 VPN

Label Stacking

- More than one label can be used for MPLS packet encapsulation Creation of a label stack
- Recap: labels correspond to Forwarding Equivalence Class (FEC)

Each label in stack used for different purposes

- Outer label always used for switching MPLS packets in network
- Remaining inner labels used to specific services/FECs, etc.
- Last label in stack marked with EOS bit
- Allows building services such as MPLS VPNs; LDP + VPN label Traffic engineering (FRR): LDP + TE label VPNs over TE core: LDP + TE + VPN label Any transport over MPLS: LDP + PW label
 Any transport over MPLS: LDP + PW label
 Layer 2/3 Packet Header

Summary

- MPLS uses labels to forward traffic
- More than one label can be used for traffic encapsulation; multiple labels make up a label stack
- Traffic is encapsulated with label(s) at ingress and at egress labels are removed in MPLS network
- MPLS network consists of PE router at ingress/egress and P routers in the core
- MPLS control plane used for signaling label mapping information to set up end-to-end Label Switched Paths
- MPLS forwarding plane used for label imposition (PUSH), swapping, and disposition (POP) operation

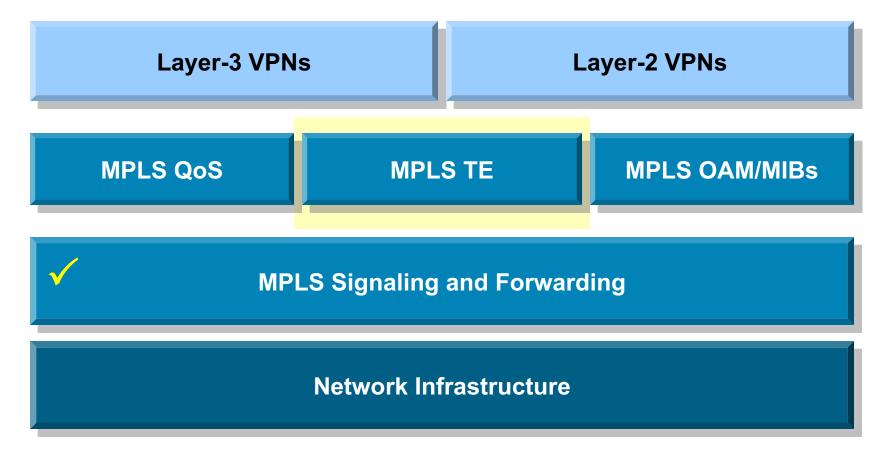
MPLS Traffic Engineering

Technology Overview and Applications



MPLS Technology Framework

 Traffic engineering capabilities for bandwidth management and network failure protection

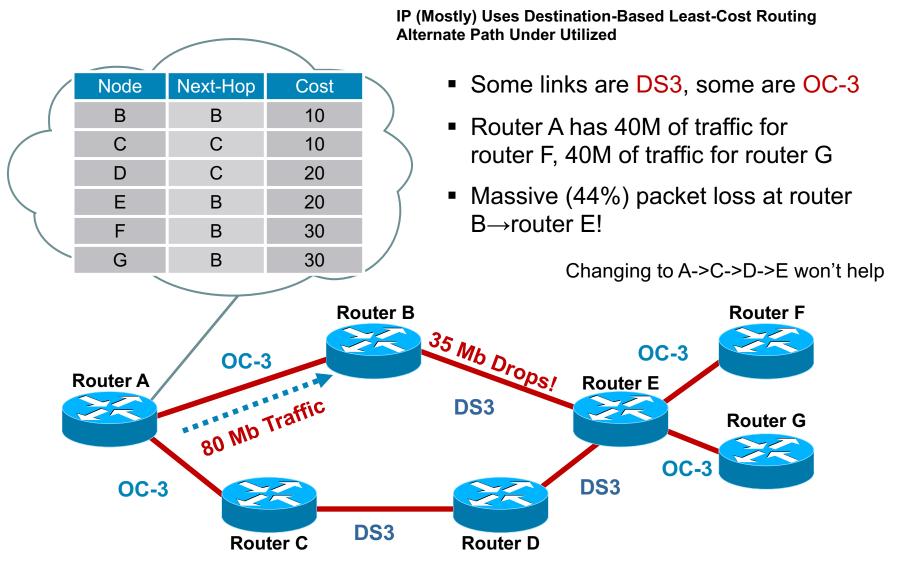


Why Traffic Engineering?

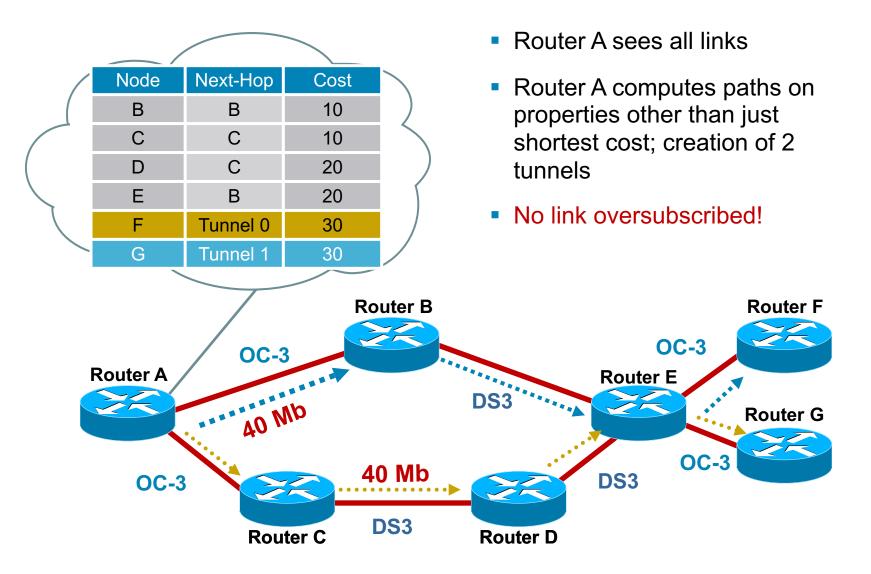
- Congestion in the network due to changing traffic patterns Election news, online trading, major sports events
- Better utilization of available bandwidth Route on the non-shortest path
- Route around failed links/nodes
 Fast rerouting around failures, transparently to users
 Like SONET APS (Automatic Protection Switching)
- Build new services—virtual leased line services
 VoIP toll-bypass applications, point-to-point bandwidth guarantees
- Capacity planning

TE improves aggregate availability of the network

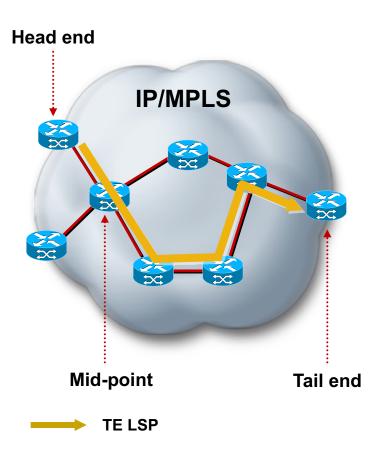
The Problem with Shortest-Path



How MPLS TE Solves the Problem

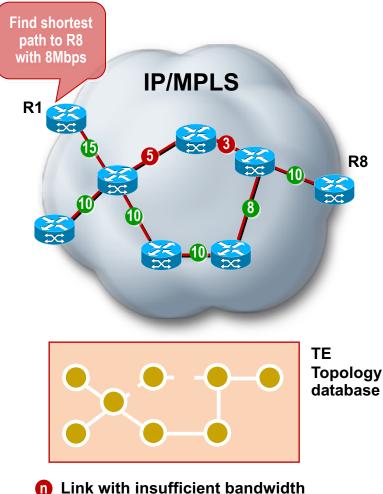


How MPLS TE Works



- Link information Distribution* ISIS-TE OSPF-TE
- Path Calculation (CSPF)*
- Path Setup (RSVP-TE)
- Forwarding Traffic down Tunnel Auto-route Static
 - Static
 - PBR
 - CBTS / PBTS
 - Forwarding Adjacency
 - **Tunnel select**

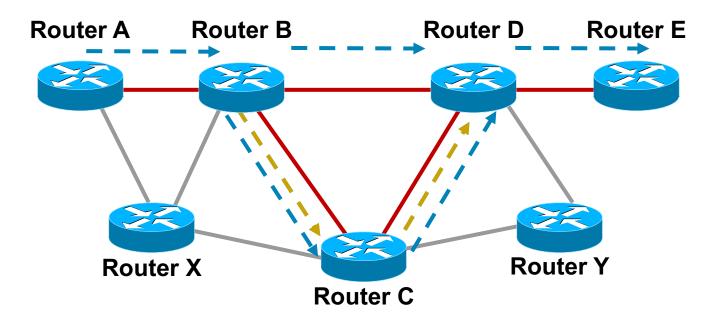
Path Calculation



Link with sufficient bandwidth

- TE nodes can perform constraint-based routing
- Constraints and topology database as input to path computation
- Shortest-path-first algorithm ignores links not meeting constraints
- Tunnel can be signaled once a path is found
- Not required if using offline path computation

MPLS TE FRR - Link Protection



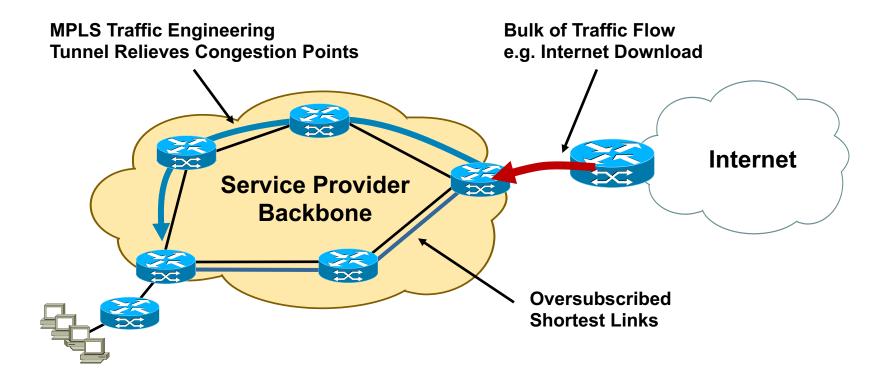
- Primary tunnel: $A \rightarrow B \rightarrow D \rightarrow E$
- Backup tunnel: $B \rightarrow C \rightarrow D$ (preprovisioned) – \rightarrow
- Recovery = ~ 50 ms

*Actual Time Varies—Well Below 50 ms in Lab Tests, Can Also Be Higher

Use Case 1: Tactical TE Deployment

Requirement: Need to Handle Scattered Congestion Points in the Network

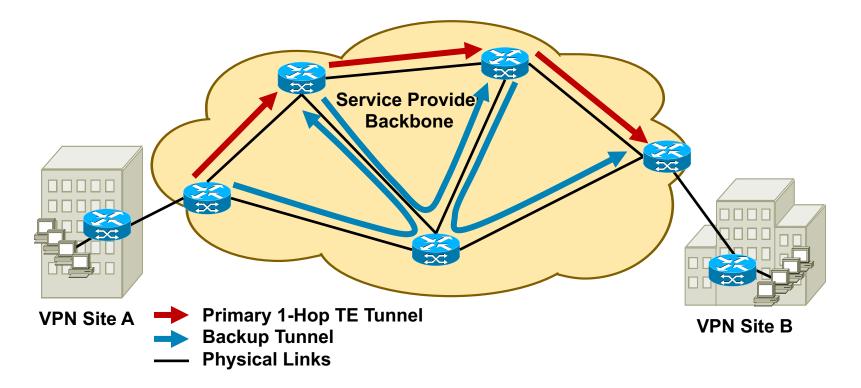
Solution: Deploy MPLS TE on Only Those Nodes that Face Congestion



Use Case 2: 1-Hop Tunnel Deployment

Requirement: Need Protection Only — Minimize Packet Loss of Bandwidth in the Core

Solution: Deploy MPLS Fast Reroute for Less than 50ms Failover Time with 1-Hop Primary TE Tunnels and Backup Tunnel for Each



MPLS TE Summary

- MPLS TE can be used to implement traffic engineering to enable enhanced network availability, utilization, and performance
- Enhanced network availability can be implemented via MPLS TE Fast Re-Route (FRR)

Link, node, and path protection

Automatically route around failed links/nodes; like SONET APS

 Better network bandwidth utilization can be implemented via creation of MPLS TE tunnels using explicit routes

Route on the non-shortest path

 MPLS TE can be used for capacity planning by creation of bandwidth-specific tunnels with explicit paths through the network Bandwidth management across links and end-to-end paths

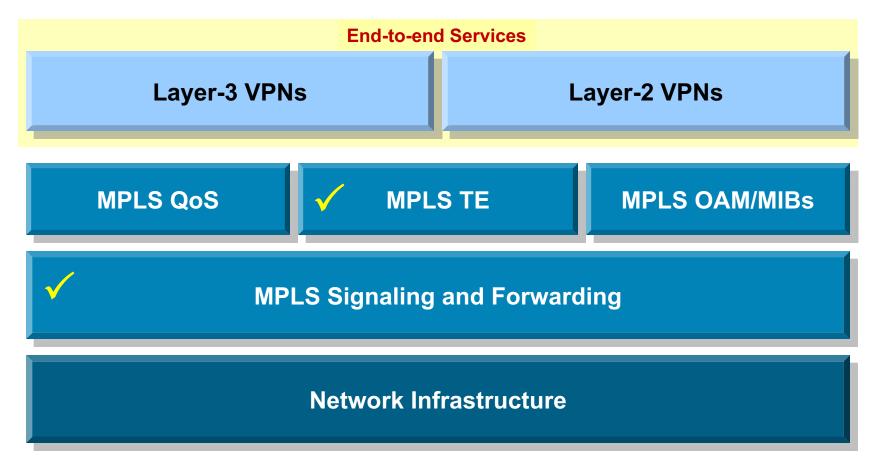
MPLS VPNs

Overviews



MPLS Technology Framework

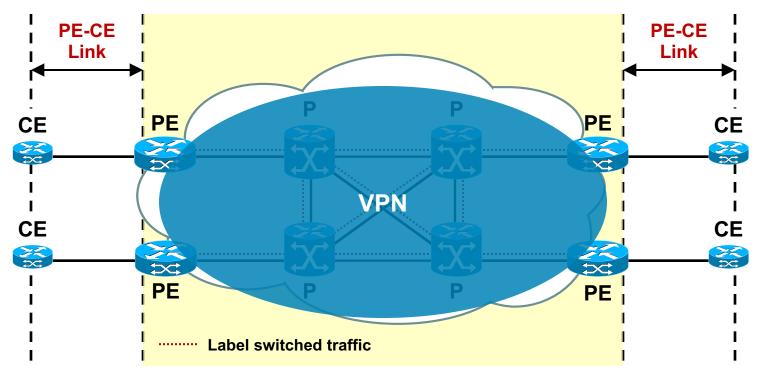
 End-to-end data connectivity services across MPLS networks (from PE to PE)



What Is a Virtual Private Network?

- VPN is a set of sites or groups which are allowed to communicate with each other in a secure way
 Typically over a shared public or private network infrastructure
- VPN is defined by a set of administrative policies Policies established by VPN customers themselves (DIY) Policies implemented by VPN service provider (managed/unmanaged)
- Different inter-site connectivity schemes possible Ranging from complete to partial mesh, hub-and-spoke
- Sites may be either within the same or in different organizations
 VPN can be either intranet or extranet
- Site may be in more than one VPN VPNs may overlap
- Not all sites have to be connected to the same service provider VPN can span multiple providers

MPLS VPN Example



PE-CE link

Connect customer network to SP network; layer-2 or layer-3

VPN

Dedicated secure connectivity over shared infrastructure

MPLS VPN Benefits

SP/Carrier perspective

Reduce costs (CAPEX)

Leverage same network for multiple services and customers

Migrate legacy networks onto single converged network Reduce costs (OPEX)

Easier service enablement; only edge node configuration

Enterprise/end-user perspective

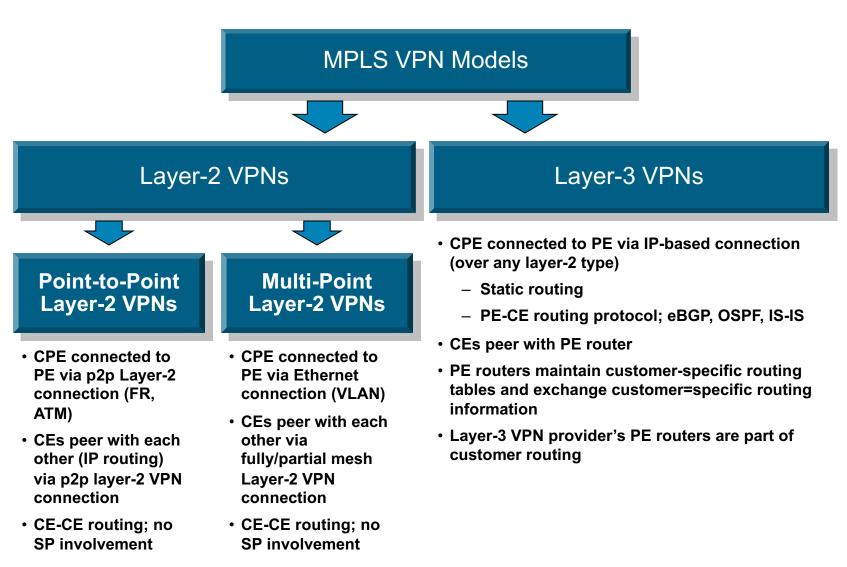
Enables site/campus network segmentation

Allows for dedicated connectivity for users, applications, etc.

Enables easier setup of WAN connectivity

Easier configuration of site-to-site WAN connectivity (for L3VPN); only one WAN connection needed

MPLS VPN Options



MPLS Layer-3 VPNs

Technology Overview and Applications

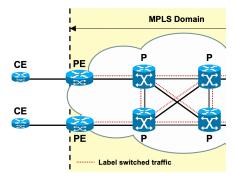


MPLS L3 VPN Technology Components

PE-CE link

Can be any type of layer-2 connection (e.g., FR, Ethernet)

CE configured to route IP traffic to/from adjacent PE Variety of routing options; static routes, eBGP, OSPF, IS-IS



- P (Provider) router = label switching router = core router (LSR) Switches MPLS-labeled packets
- PE (Provider Edge) router = edge router (LSR) Imposes and removes MPLS labels
- CE (Customer Edge) router Connects customer network to MPLS network

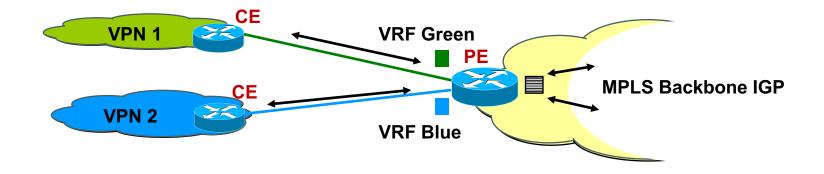
MPLS L3VPN Control Plane

Separation of customer routing via virtual VPN routing table In PE router: customer I/Fs connected to virtual routing table Between PE routers: customer routes exchanged via BGP

MPLS L3VPN Forwarding Plane

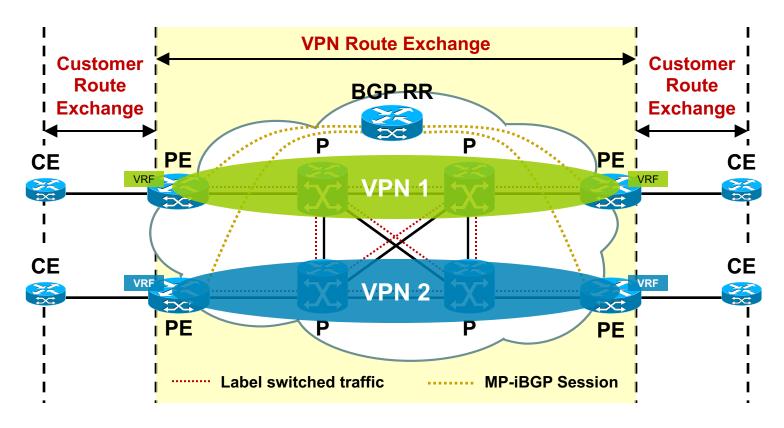
Separation of customer VPN traffic via additional VPN label VPN label used by receiving PE to identify VPN routing table

Virtual Routing and Forwarding Instance



- Virtual Routing and Forwarding Instance (VRF)
- Typically one VRF created for each customer VPN on PE router
- VRF associated with one or more customer interfaces
- VRF has its own instance of routing table (RIB) and forwarding table (CEF)
- VRF has its own instance for PE-CE configured routing protocols

VPN Route Distribution

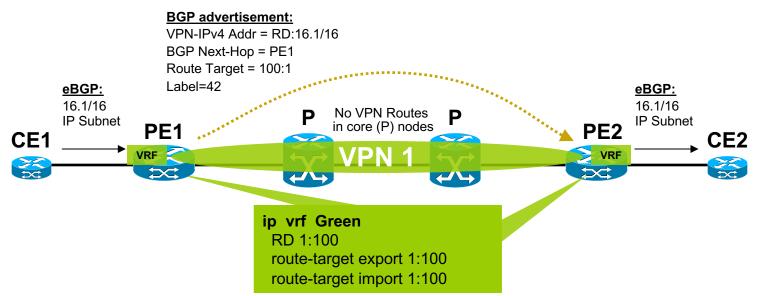


Full mesh of BGP sessions among all PE routers

Multi-Protocol BGP extensions (MP-iBGP)

Typically BGP Route Reflector (RR) used for improved scalability

VPN Control Plane Processing



Make customer routes unique:

- Route Distinguisher (RD): 8-byte field, VRF parameters; unique value assigned by a provider to each VPN to make different VPN routes unique
- VPNv4 address: RD+VPN IP prefix

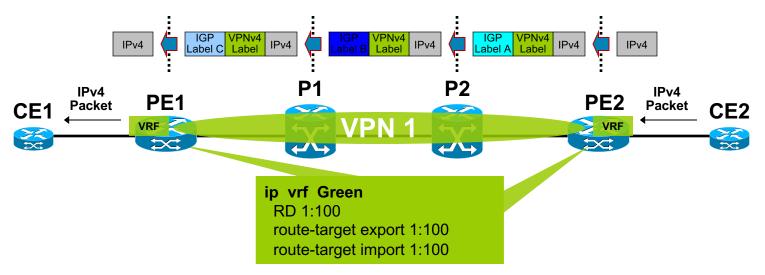
Selective distribute customer routes:

- Route Target (RT): 8-byte field, VRF parameter, unique value to define the import/export rules for VPNv4 routes
- MP-iBGP: advertises VPNv4* prefixes + labels

Processing Steps:

- 1. CE1 redistribute IPv4 route to PE1 via eBGP.
- 2. PE1 allocates VPN label for prefix learnt from CE1 to create unique VPNv4 route
- 3. PE1 redistributes VPNv4 route into MP-iBGP, it sets itself as a next hop and relays VPN site routes to PE2
- PE2 receives VPNv4 route and, via processing in local VRF (green), it redistributes original IPv4 route to CE2.

VPN Forwarding Plane Processing



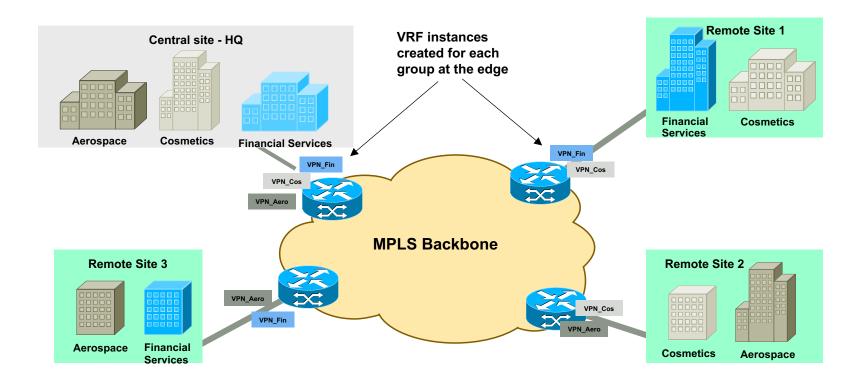
Processing Steps:

- 1. CE2 forwards IPv4 packet to PE2.
- 2. PE2 imposes pre-allocated VPN label (learned via MP-IBGP) to IPv4 packet received from CE2.
- 3. PE2 imposes outer IGP label (learned via LDP) and forwards labeled packet to next-hop P-router P2.
- 4. P-routers P1 and P2 swap outer IGP label and forward label packet to PE1.
- 5. Router PE1 strips VPN label and forwards IPv4 packet to CE1.

Use Case 1: Traffic Separation

<u>Requirement:</u> Need to ensure data separation between Aerospace, Cosmetics and Financial Services, while leveraging a shared infrastructure

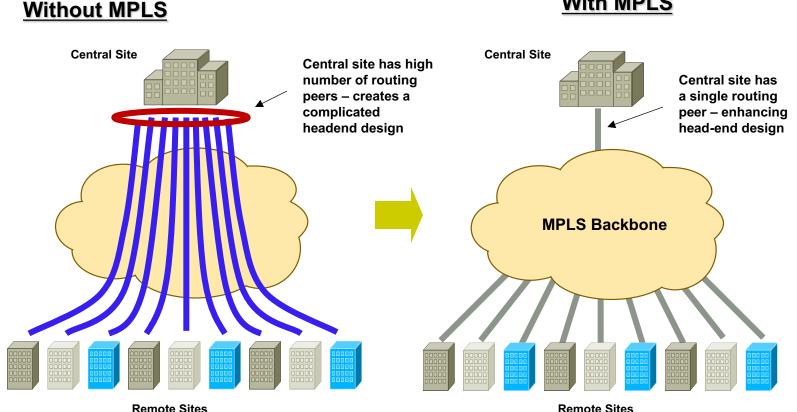
Solution: Create MPLS VPN for each group



Use Case 2: Simplify Hub Site Design

Requirement: To ease the scale and design of head-end site

Solution: Implement MPLS Layer 3 VPNs, which reduces the number of routing peers of the central site



With MPLS

MPLS Layer-3 VPN Summary

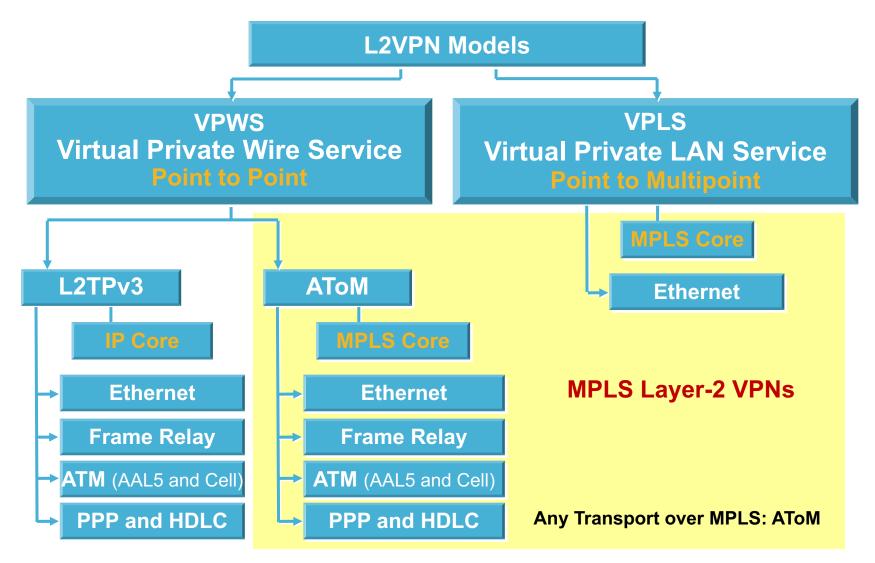
- Provide layer-3 connectivity among CE sites via IP peering (across PE-CE link)
- Implemented via VRFs on edge/PE nodes providing customer route and forwarding segmentation
- BGP used for control plane to exchange customer VPN (VPNv4) routes between PE routers
- MPLS VPNs enable full-mesh, hub-and-spoke, and hybrid IP connectivity among connected CE sites

MPLS Layer-2 VPNs

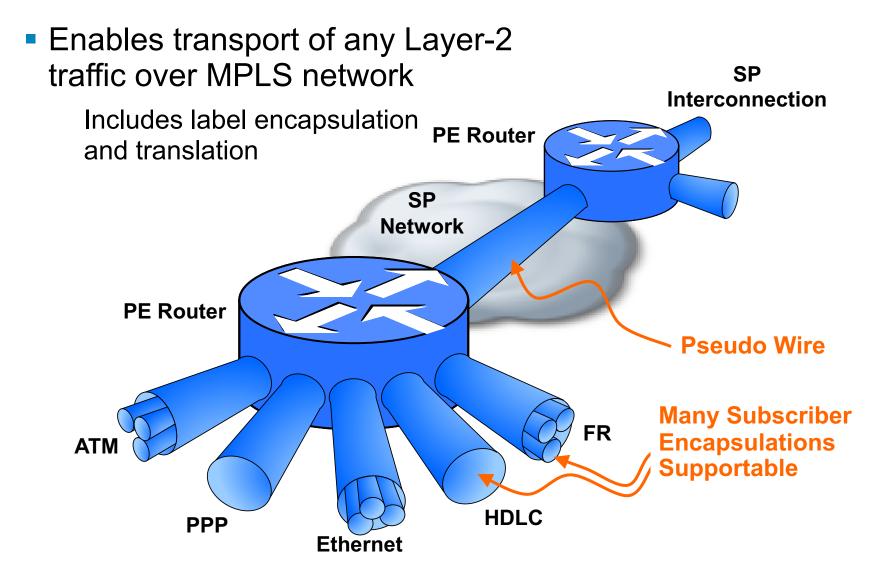
Technology Overview and Applications



L2VPN Options



Layer-2 VPN Overview



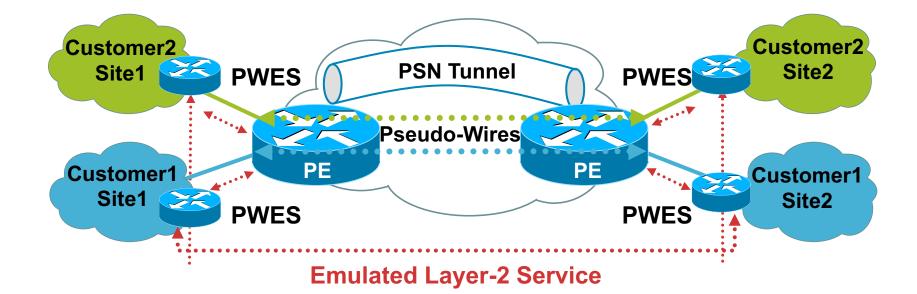
Any Transport over MPLS Architecture

- Based on IETF's Pseudo-Wire (PW) Reference Model
- PW is a connection (tunnel) between 2 PE Devices, which connects 2 PW End-Services

PW connects 2 Attachment Circuits (ACs)

Bi-directional (for p2p connections)

Use of PW/VC label for encapsulation



AToM Technology Components

PE-CE link

Referred to as Attachment Circuit (AC)

Can be any type of layer-2 connection (e.g., FR, Ethernet)

AToM Control Plane

Targeted LDP (Label Distribution Protocol) Session

Virtual Connection (VC)-label negotiation, withdrawal, error notification

AToM Forwarding Plane

2 labels used for encapsulation + control word

Outer tunnel (LDP) label

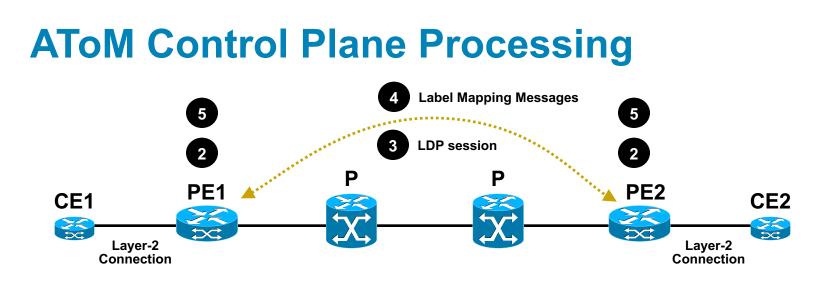
To get from ingress to egress PE using MPLS LSP

Inner de-multiplexer (VC) label

To identify L2 circuit (packet) encapsulated within tunnel label

Control word

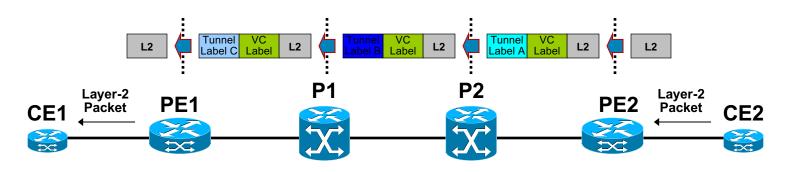
Replaces layer-2 header at ingress; used to rebuild layer-2 header at egress



Processing Steps (for both P1 and P2):

- 1. CE1 and CE2 are connected to PE routers via layer-2 connections
- 2. Via CLI, a new virtual circuit cross-connect is configured, connecting customer interface to manually provided VC ID with target remote PE
- 3. New targeted LDP session between PE routers established, in case one does not already exist
- 4. PE binds VC label with customer layer-2 interface and sends labelmapping message to remote PE over LDP session
- 5. Remote PE receives LDP label binding message and matches VC ID with local configured cross-connect

AToM Forwarding Plane Processing



Processing Steps:

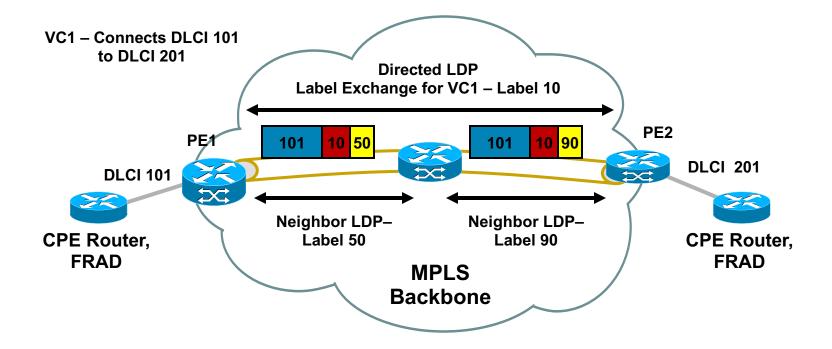
- 1. CE2 forwards layer-2 packet to PE2.
- 2. PE2 imposes VC (inner) label to layer-2 packet received from CE2 and optionally a control word as well (not shown).
- 3. PE2 imposes Tunnel outer label and forwards packet to P2.
- 4. P2 and P1 router forwards packet using outer (tunnel) label.
- Router PE2 strips Tunnel label and, based on VC label, layer-2 packet is forwarded to customer interface to CE1, after VC label is removed

In case control word is used, new layer-2 header is generated first.

Use Case: L2 Network Interconnect

<u>Requirement:</u> Need to create connectivity between remote customer sites, currently interconnected via Frame Relay WAN connectivity. Only point-to-point connectivity required.

Solution: Interconnect AToM PW between sites, enabling transparent Frame Relay WAN connectivity.



Layer-2 VPN Summary

- Enables transport of any Layer-2 traffic over MPLS network
- Two types of L2 VPNs; AToM for point-to-point and VPLS point-to-multipoint layer-2 connectivity
- Layer-2 VPN forwarding based on Pseudo Wires (PW), which use VC label for L2 packet encapsulation

LDP used for PW signaling

 AToM PWs suited for implementing transparent point-topoint connectivity between Layer-2 circuits

#