Latency in DOCSIS Networks



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The various DOCSIS versions

From a latency perspective

- DOCSIS 1.0 ca. 1996, deployments ~1998
 - Fundamental request-grant upstream MAC layer definition
- DOCSIS 1.1 ca. 1999, deployments ~2001
 - Additions for configured Quality of Service
 - Packet classifiers
 - Flow independence
 - QoS configuration per flow
- DOCSIS 2.0 ca. 2001, deployments ~2003
 - TCP ACK suppression/prioritization
- DOCSIS 3.0 ca. 2006, deployments ~ 2007
 - Buffer Control feature added in 2011
- DOCSIS 3.1 ca. 2013, deployments ~2015
 - Light Sleep Mode
 - AQM Mandatory

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DOCSIS Request-Grant Upstream MAC

- Upstream channel scheduling is driven by "MAP" Intervals (typ. 2ms)
- Packet(s) arrive at the cable modem (CM)
- CM waits* for the next contention request opportunity
 - *typically less than 2ms
- CM sends request message (subject to rate shaping)
- CMTS** scheduler collects requests, then schedules and communicates future transmit opportunities (grants)
- Due to serialization, propagation and interleaver delays, as well as CMTS/CM processing delays, grant occurs 2 MAP Intervals after the request was sent
- Without congestion, typically 4-8ms access latency

**Cable Modem Termination System

Quality of Service (D1.1 and above)

- For known applications with known QoS requirements
- Operator configures packet classifiers and service flow QoS parameters
 - Token bucket rate shaping, priority, guaranteed rate, lowlatency scheduling, etc.
- Service Flows queue traffic and access channel independently
- Modems today support 16 or 32 service flows, each with an independent hardware queue



TCP ACK Suppression/Prioritization Dealing with buffer bloat before "Bufferbloat"

- Queue build-up from upstream TCP sessions delays upstream TCP ACKs downstream throughput suffers.
- All modem vendors implement proprietary mechanisms in D2.0 and above to move ACKs to the head of the queue and discard superfluous ACKs.

• TCP RTT depends on which side you measure from



Buffer Control (D3.0) Bufferbloat is everywhere!

 Amended specification in 2011 to allow operator to set perservice flow buffer sizes.

- Requires configuration by operator
- Interest was high, adoption has been slow.

Light Sleep Mode (D3.1)

Reducing Energy Consumption during "Idle"

- Response to Political Pressures such as:
 - More Efficient Modems, Routers Could Save Consumers \$330 Million Annually – NRDC
 - "These small, innocuous black boxes that never sleep consume enough electricity each year to
 power all 1.2 million homes in the Silicon Valley area, the hi-tech capital of the world," said NRDC
 senior scientist Noah Horowitz. "Small network devices suck roughly the same amount of energy
 around the clock, whether or not you are sending or receiving any data. But there are steps that
 manufacturers can and should take to make sure these devices are no longer energy vampires."
 - 88 million Internet consumers in US:
 - More Efficient Modems, Routers Could Save Each Consumer \$0.31 Monthly
 - EPA Energy Star Small Network Equipment Spec
 - California Energy Commission Consumer Electronics Efficiency Pre-Rulemaking
 - EU Lot 26 & Networked Standby Regulation
 - EU Broadband Code of Conduct

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Light Sleep Mode (D3.1)

- Modems in Light Sleep Mode will shut down receiver for periods of up to 200ms.
 - Interval set by CMTS, could be less.
 - Downstream packets queued at CMTS until wake interval

- Baseline latency measured during network idle conditions may not give you the results you expect.
- Latency under load might actually be better than "baseline"



Active Queue Management (D3.1)

- AQM will be mandatory for both CM and CMTS in D3.1
- On by default, can be disabled on a per service flow basis
- CMTS can implement an algo. of the vendor's choosing
- CM MUST implement single-queue PIE, but can also implement other algorithms
 - PIE chosen over CoDel, CoDel-DT, SFQ-CoDel, SFQ-PIE
- Currently investigating if existing D3.0 equipment can be upgraded to support AQM

Why no *FQ?

- Hardware complexity of 32 Service Flows x 32 queues
 - Or, operational complexity of Service Flows sharing a pool of N queues
- Tight deadlines between MAP & grant
 - *any* additional processing at dequeue time is hard
- Limited additional benefit compared to single queue AQM at 100Mbps+
- Concerns about VPN traffic
- Hash collisions not feasible to have 1024 queues (see above)

