#### **Lecture 7: Instruction Set Architecture**



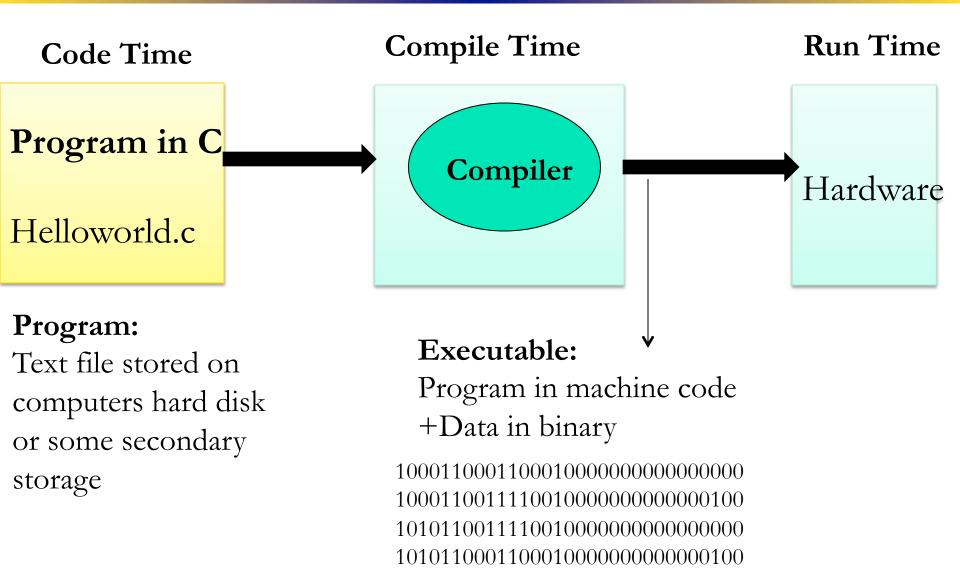
CSE 30: Computer Organization and Systems Programming Winter 2014 Diba Mirza Dept. of Computer Science and Engineering University of California, San Diego

# Outline

- 1. Steps in program translation
- 2. Hardware/Software Interface Preliminaries
  - 1. Instruction Set Architecture
    - 1. General ISA Design (Architecture)
    - 2. Architecture vs. Micro architecture
    - 3. Different types of ISA: RISC vs CISC
  - 2. Assembly programmer's view of the system
    - 1. Registers: Special and general purpose
    - 2. Assembly and machine code (program translation detail)
    - 3. Layout of ARM instructions in memory
  - 3. Steps in program execution
  - 4. Basic Types of ARM Assembly Instructions

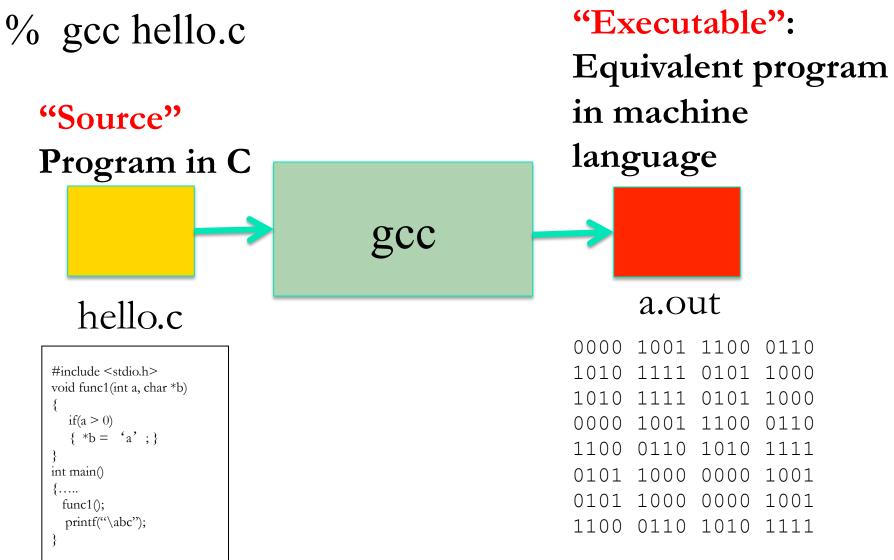


# Steps in program translation





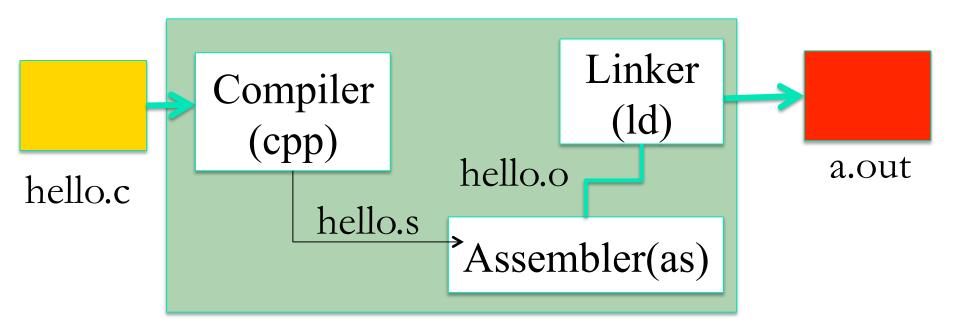
## Compile time: What does gcc do?





# Steps in gcc

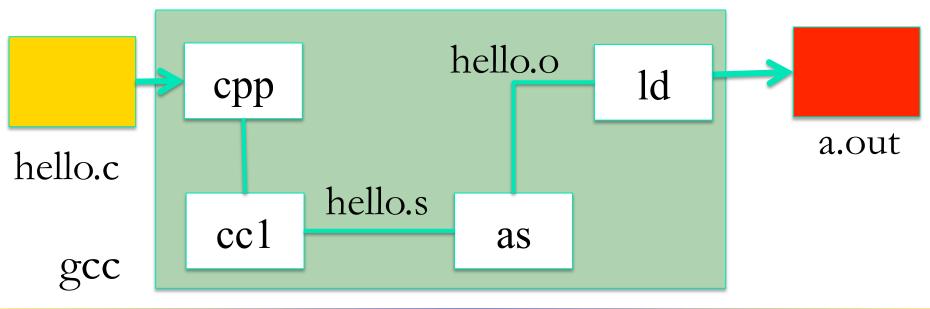
The translation is actually done in a number of steps





# Steps in gcc

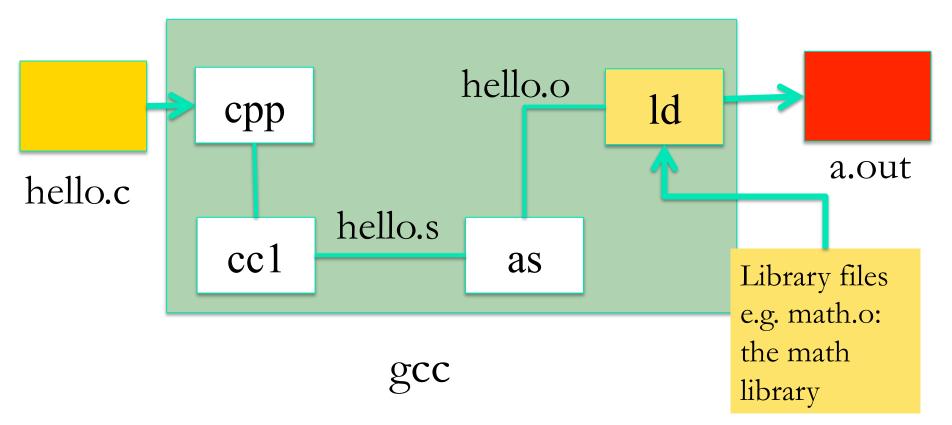
Ask compiler to show temporary files:
% gcc –S hello.c (gives hello.s – assembly code)
% gcc –c hello.c (gives hello.o – object module)
% gcc –o prog\_hello hello.c (gives prog\_hello.o - named executable)





# Include code written by others

- Code written by others (libraries) can be included
- Id (linkage editor) merges one or more object files with the relevant libraries to produce a single executable





# **Assembly Language**

A. Is the binary representation of a program.

B. A symbolic representation of machine instructions

c. A set of instructions that the machine can directly execute

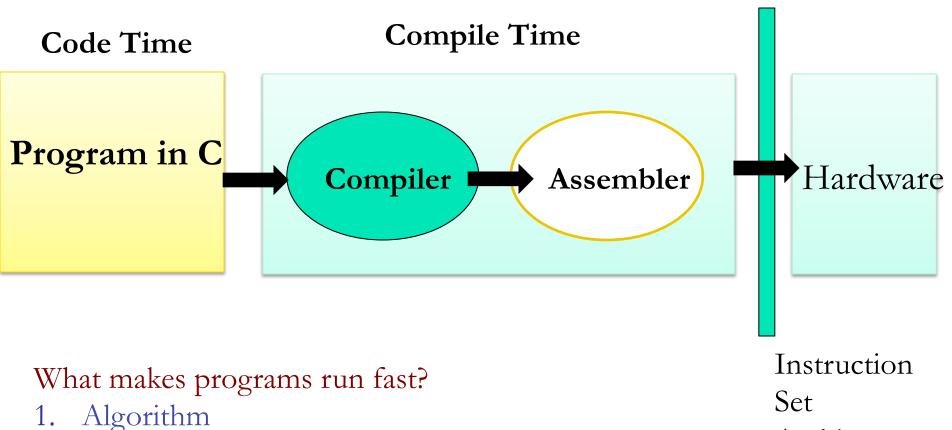


## Machine vs Assembly Language

- Machine Language: A particular set of instructions that the CPU can directly execute – but these are ones and zeros
- Assembly language is a symbolic version of the equivalent machine language
  - each statement (called an <u>Instruction</u>), executes exactly one of a short list of simple commands
  - Unlike in C (and most other High Level Languages), each line of assembly code contains at most 1 instruction
  - Instructions are related to operations (e.g. =, +, -, \*) in C or Java



# Steps in program translation



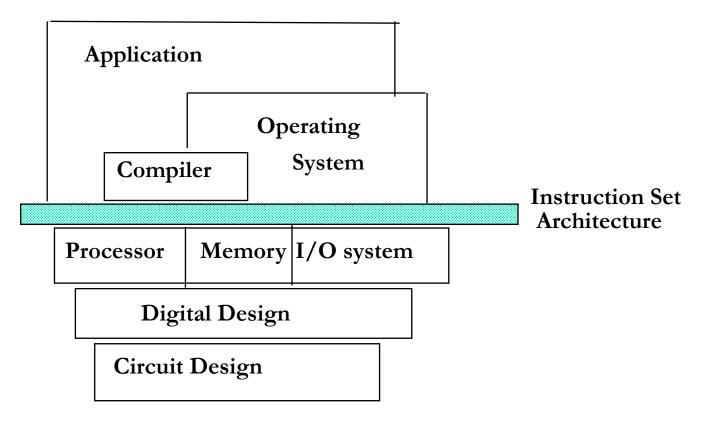
- Compiler Translation to Machine code 2.
- ISA and hardware implementation 3.

Architecture



#### What is the Instruction Set Architecture?

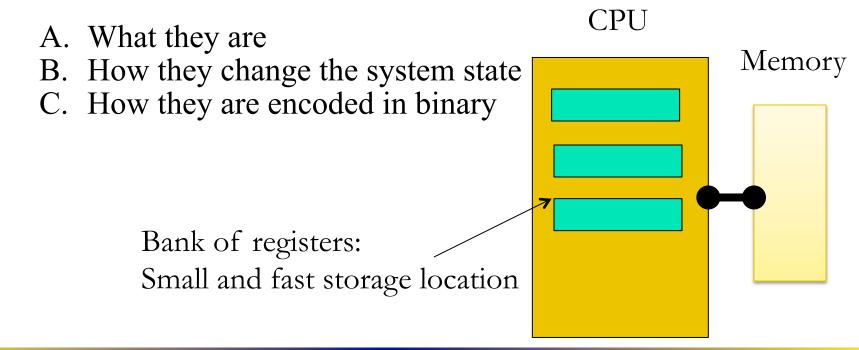
- Agreed-upon interface between all the software that runs on the machine and the hardware that executes it
- Primitive set of instructions a particular CPU implements





#### **General ISA Design Aspects**

- Everything about h/w that is visible to the s/w and can be manipulated by it via basic machine instructions.
   Example: Registers: How many? What size? Memory: How to access contents?
- 2. The set of basic machine instructions:





# Is the ISA different for different CPUs?

- Different CPUs implement different sets of instructions.
  - Examples: ARM, Intel x86, IBM/Motorola PowerPC (Macintosh), MIPS, Inter IA32 ...

- Two styles of CPU design:
  - RISC (Reduced Instruction Set Computing)
  - CISC (Complex Instruction Set Computing)



# **RISC versus CISC (Historically)**

- Complex Instruction Set Computing e.g x86
  - Larger instruction set
  - More complicated instructions built into hardware
  - Variable length
  - Multiple clock cycles per instruction
- Reduced Instruction Set Computing e.g. ARM
  - Small, highly optimized set of instructions
  - Memory accesses are specific instructions
  - One instruction per clock cycle
  - Instructions are of the same size and fixed format



 $A = A^*B$ 

#### RISC CISC

- LOAD A, eax MULT B, A
- LOAD B, ebx
- PROD eax, ebx
- STORE ebx, A



## **RISC vs CISC**

#### RISC

- More work for compiler
- More RAM used to store instructions
- Easier to debug: more reliable processors
- Easier to optimize
  - One clock cycle per instruction

#### CISC

- Less work for compiler
- Fewer instructions to store
- Harder to maintain and debug



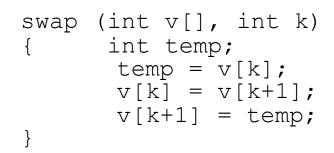
# For a program written in a high-level-language, the conversion from \_\_\_\_ to \_\_\_\_ will be \_\_\_\_\_ for different target processors

	High level Language to Assembly	Assembly to Machine Language
Α	One-to-Many	Many-to-Many
В	Many-to-Many	One-to-One
С	One-to-Many	One-to-One
D	One-to-One	One-to-One

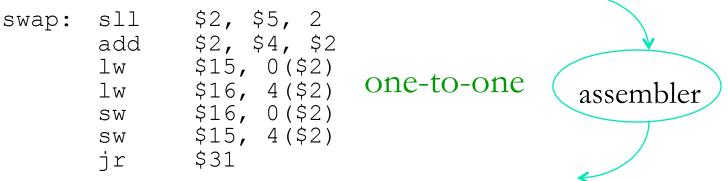


### Translations

#### High-level language program (in C)



Assembly language program (for MIPS)



one-to-many

C compiler

Machine (object, binary) code (for MIPS)

00000000000001010001000010000000000000010000010000100000100000



### Architecture vs Micro-architecture

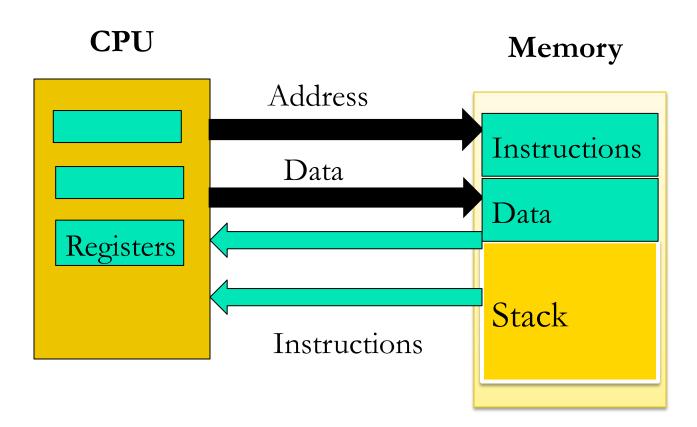
#### Architecture:

- Parts of processor design needed to write programs in assembly
- What is visible to s/w
   E.g Number of registers

#### Micro-Architecture:

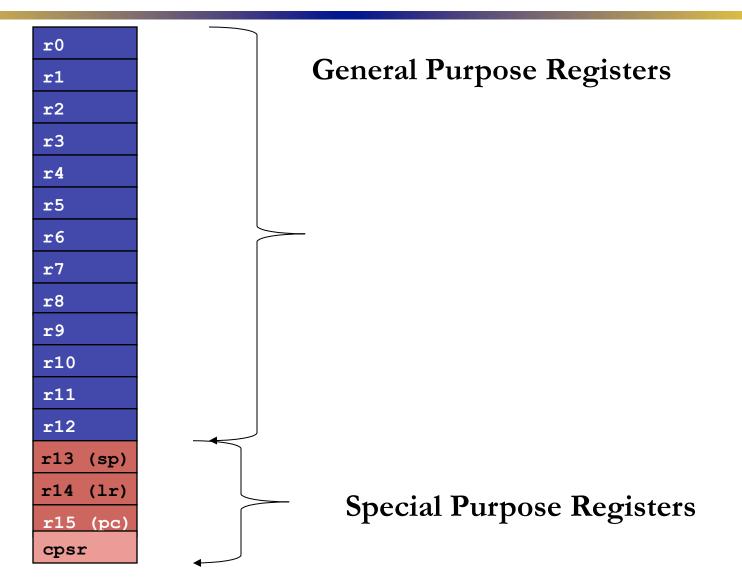


#### The Assembly Programmer's View of the machine



- Registers: (Very) Small amount of memory inside the CPU
- Each ARM register is 32 bits wide
  - Groups of 32 bits called a <u>word</u> in ARM
- ₹UCSD

#### **The ARM Register Set**





#### **Assembly Variables: Registers**

- Unlike HLL like C or Java, assembly cannot use variables
  - Why not? Keep Hardware Simple
- Data is put into a register before it is used for arithmetic, tested, etc.
- Manipulated data is then stored back in main memory.
- Benefit: Since registers are directly in hardware, they are very fast



#### C, Java Variables vs. Registers

- In C (and most High Level Languages) variables declared first and given a type
  - Example: int fahr, celsius; char a, b, c, d, e;
- Each variable can ONLY represent a value of the type it was declared as (cannot mix and match int and char variables).
- In Assembly Language, the registers have no type; operation determines how register contents are treated



# Which one of the following is an optimization that is typically done by ARM compilers?

- A. Put the address of frequently used data in registers
- B. Put the value of frequently used data in registers
- C. Put as much data as possible in the registers
- D. Put as few program data in the registers
- E. Registers are not available to the programmer/compiler they are only used internally by the processor to fetch and decode instructions



#### Layout of instructions in memory

```
swap (int v[], int k)
{ int temp;
   temp = v[k];
   v[k] = v[k+1];
   v[k+1] = temp;
}
```

0x4000 Swap:	mov	r2, r5
0x4004	add	r4, r2, #1
0x4008	ldr	r10, [r6,r2]
0x400c	ldr	r11, [r6,r4]
0x4010	mov	r1, r10
0x4014	str	r11,[r6,r2]
0x4018	str	r1, [r6,r4]
0x401c	bx	lr



### **Machine Instructions**

A. Can be thought of as numbers.

B. Are commands the computer performs.

C. Were what people originally wrote programs using.

D. All of the above.



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#### Steps in program execution

1						
	▼0x4000	Swap:	mov	r2, r5		
	0x4004		add	r4, r2, #1		
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	0x4014		str	r11,[r6,r2]		
	0x4018		str	r1, [r6,r4]		
	0x401c		bx	lr		
	L					
			Inst	ructions		
			Data	a		
		•	•			
	PC:0x4000		0.	1	Memory	
			Stac	K Memo		
	CPU					



#### **Assembly Language**

A. Is in binary.

B. Allows languages to be designed for their specific uses.

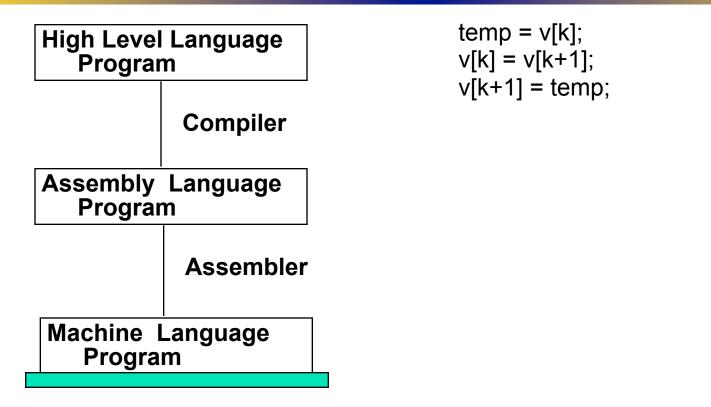
**C**. Has one line for every machine language instruction.



# **Basic Types of Instructions**

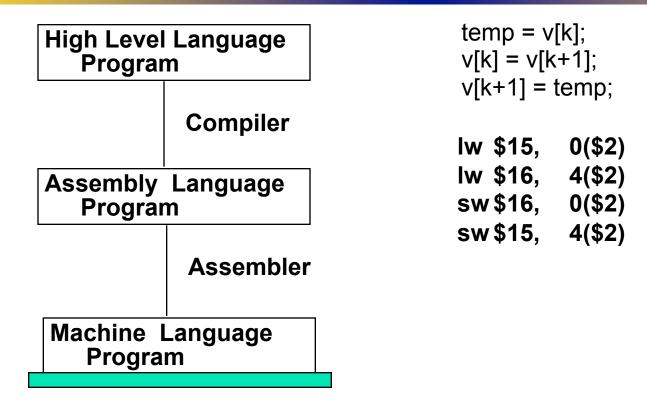
- 1. Arithmetic: Only processor and registers involved
  - 1. compute the sum (or difference) of two registers, store the result in a register
  - 2. move the contents of one register to another
- 2. Data Transfer Instructions: Interacts with memory
  - load a word from memory into a register
  - 2. store the contents of a register into a memory word
- 3. Control Transfer Instructions: Change flow of execution
  - 1. jump to another instruction
  - 2. conditional jump (e.g., branch if registeri == 0)
  - 3. jump to a subroutine





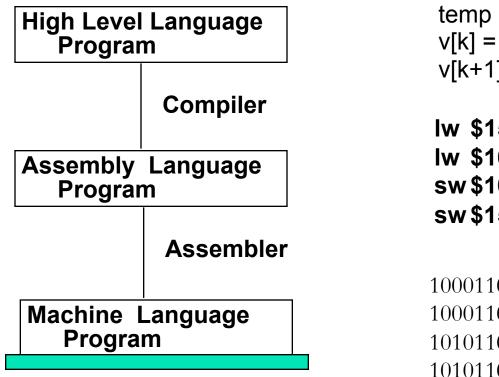
**Machine Interpretation** 



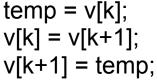


**Machine Interpretation** 



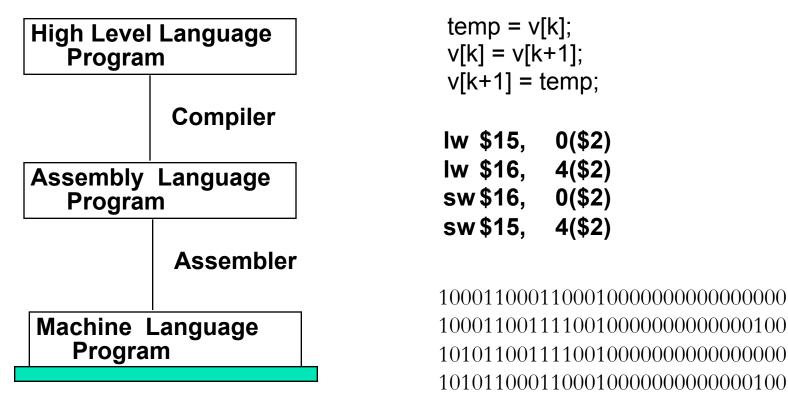


**Machine Interpretation** 



Iw \$15, 0(\$2) Iw \$16, 4(\$2) sw \$16, 0(\$2) sw \$15, 4(\$2)





**Machine Interpretation** 

Machine does something!

