

Linux centered heterogeneous multi-core architectures

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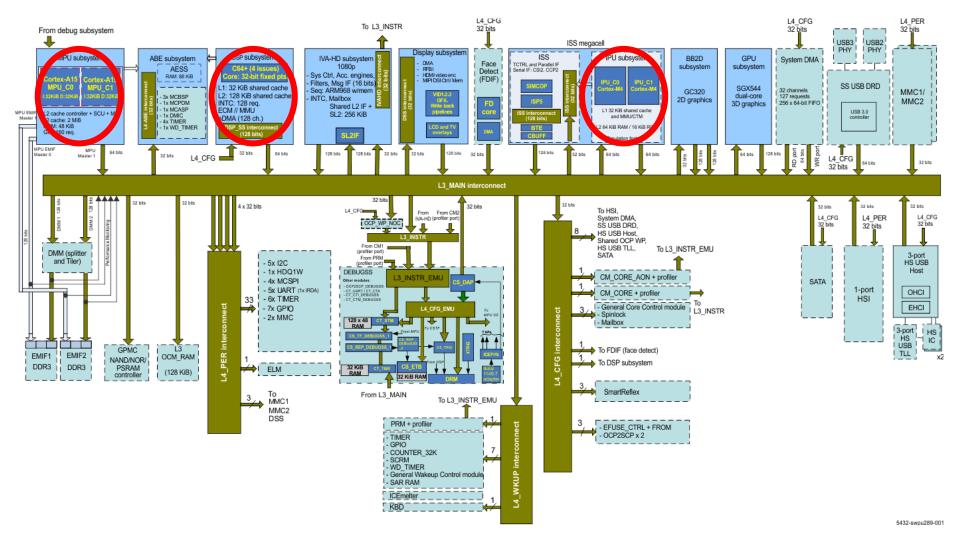
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- Evolution of HW platforms
- Rationale for heterogeneous multi-core architectures
- Complex SW platforms: problems and requirements
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- An rpsmsg based inter-kernels IPC service and its socket API
- Exporting the rpmsg bus interface to the user space
- Linux topics
 - Device model
 - Device Tree
 - Device driver API
 - Sockets
 - Platform and misc drivers

HW platforms: TI OMAP 5





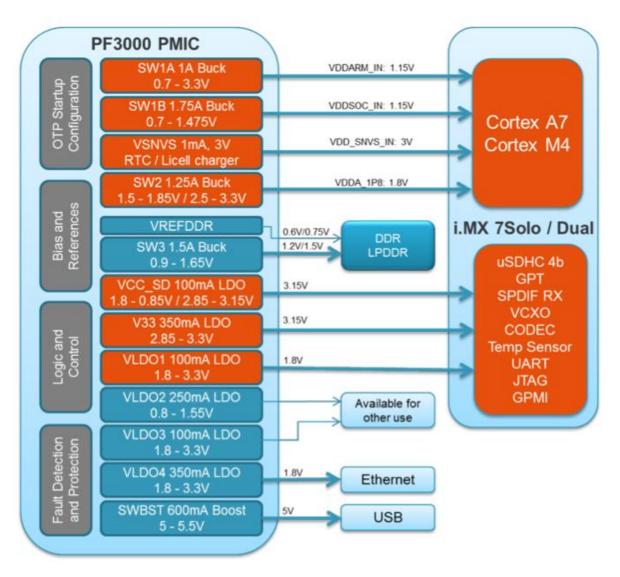
HW platforms: TI Sitara AM572x



AM572x				
MPU (2x ARM® Cortex™–A15)	IVA HD 1080p Video Co-Processor	Display Subs	LCD1 LCD2	
GPU (2x SGX544 3D)	BB2D (GC320 2D)	Blend / Scale	HDMI 1.4a	
DSP (2x C66x [™] Co-Processor) IPU1 (2x Cortex [™] -M4) IPU2 (2x Cortex [™] -M4)				
EDMA SDMA MMU x2 VIP x3 VPE				
High-Speed Interconnect				
System		Connec	Connectivity	
	ners x16 PWM SS x	USB 3.0 Dual Role FS/HS/SS	PCle SS x2	
	WDT HDQ	w/ PHYs	PRU-ICSS x2	
GPIO x8 R	TC SS KBD	USB 2.0 Dual Role FS/HS w/ PHY	GMAC_SW	
Serial Interfaces		Program/Data Storage		
UART x10	QSPI			
	CASP x8 Up to 2 OCMC w/ E	.5 MiB GPMC / ELM RAM (NAND/NOR/	EMIF x2 2x 32-bit DDR3(L) ECC	

HW platforms: NXP i.MX 7Solo/Dual

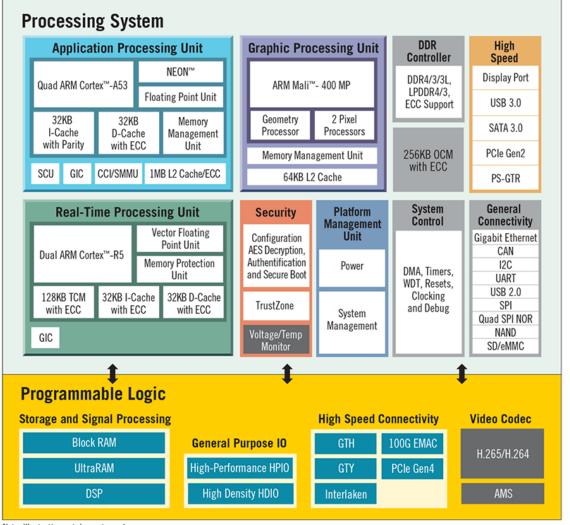




HW platforms: Xilinx UltraScale MPSoC



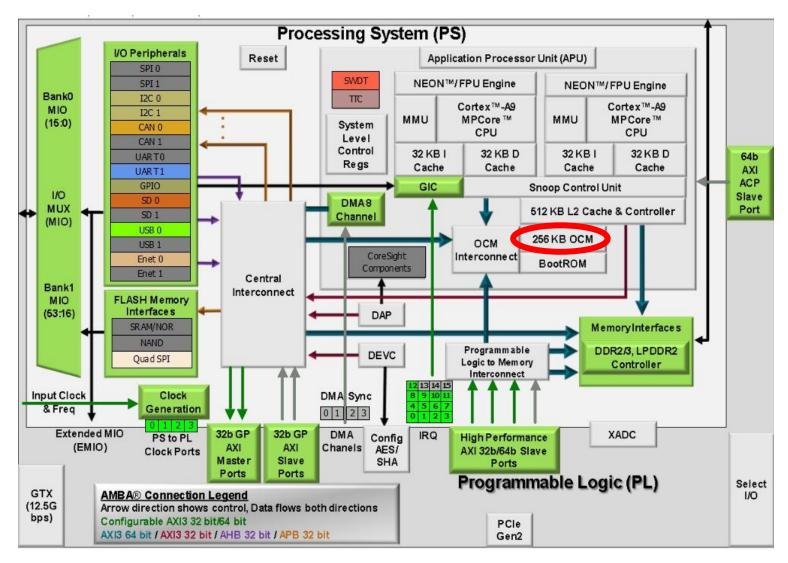
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Note: Illustration not drawn to scale.

HW platforms: Xilinx Zynq and SW generated heterogeneity





Why AMP? (Asymmetric Multi-Processing)



- Consolidation of applications
 - Reuse
 - Space / weight / power reduction
 - "A growing number of embedded use cases require concurrent execution of isolated SW environments within the system" (F. Baum, Mentor Graphics)
- Robustness / security
- Boot time
- Heterogeneous functional and performance requirements
 - > Real time

Heterogeneous requirements: Computing vs. controlling



Computing

- Large and complex applications
- Heavy computational requirements
- Real time / high throughput
- Complex arithmetic
- Large data movements

Controlling

- Real time / determinism
- Minimum latency

Addressing design challenges in heterogeneous multicore embedded systems (W. Kurisu, Mentor Graphics)



- Each device runs its own operating system or operating environment
- 2. Each device runs on its own discrete processor and those processors are typically different the type of application drives the processor selection, ranging from lowend microcontrollers to high-end application processors; each component of the system has full ownership of all the hardware available to the component. Examples of that hardware include the processors, graphic processing units, memory, I/O, cache, etc.
- 3. The discrete components of this system are typically loosely connected each component boots independently (?) and communicates with each other through messages over some physical connection.

Complex SW platforms



- Multiple kernels and multiple independent instances of a kernel on the same chip
 - Linux
 - RTOS (e.g. FreeRTOS is not multicore!)
- Partitioning of resources
- Boot & life cycle of processing cores and kernels
- Interprocessor/intercore communications
- Interprocess communications (IPC)
- Programming model

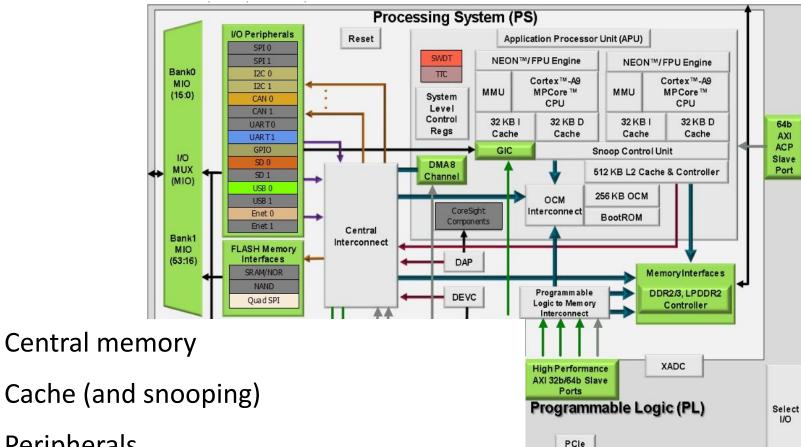
Why Linux centered?



- Because of silicon vendors' support (of the main processor of the chip)
- Because of Linux support of "computing" requirements
- Because Linux already supports, to some extent, heterogeneous architectures

Partitioning of resources





Gen2

- Cache (and snooping)
- Peripherals
 - Interrupts (and handling of PIC)
 - Virtual I/O

Boot



- Pin mux
 - Consistency with partitioning of resources
 - Implications on SW factory
- Loading of executable images and coordination with life cycle management
 - MMU to match relocation address of RTOS-based executable images

Life cycle



- Coordinated start/stop of different kernel instances
- remoteproc
 - Developed by TI
 - Master-slave architecture
 - Integrated with rpmsg support of interprocessor/intercore communications
 - Allocation and initialization of shared memory communication resources
 - 2 cores can communicate via rpmsg only if one is the remoteproc master of the other
 - Integrated in Linux main branch (master role only)
 - Implemented as a platform driver

Interprocessor/intercore communications: rpmsg



- Developed by TI
- Based on standard Linux components (virtio)
- Point-to-point architecture between remoteproc master and its remotes (host-device pattern)
- Message style communications, based on circular buffers in shared memory (2 uni-directional vrings per point-to-point connection)
- Cache configuration must guarantee that communicating cores have a coherent view of shared memory
- Integrated in Linux as a bus driver (an I/O subsystem)
 - The API offered by the rpmsg bus driver in Linux is in kernel space!
 - Several client drivers (network or character drivers) can support different transport/application dialogues on the same rpmsg bus 16

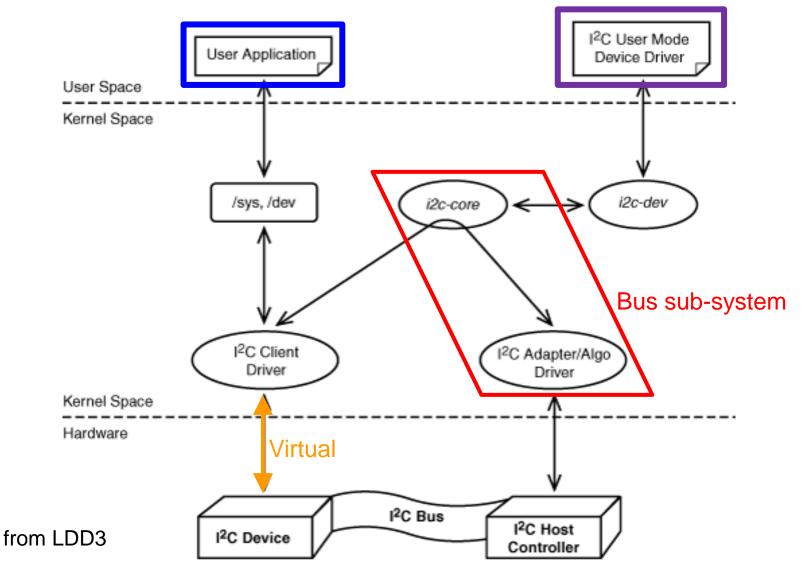
rpmsg communication topology



- Constrained by connection with remoteproc
- Allows only communications between a remoteproc master and each of its remotes
 - 2 uni-directional vrings are created for communications between the remoteproc master and each remote
- No support for routing (e.g. by remoteproc/rpmsg master)
- Supported communication topologies
 - Star
 - Tree (restricted to directly linked nodes)
- Linux support limited to rpmsg master side (center of star)

Linux I/O subsystem (device model)



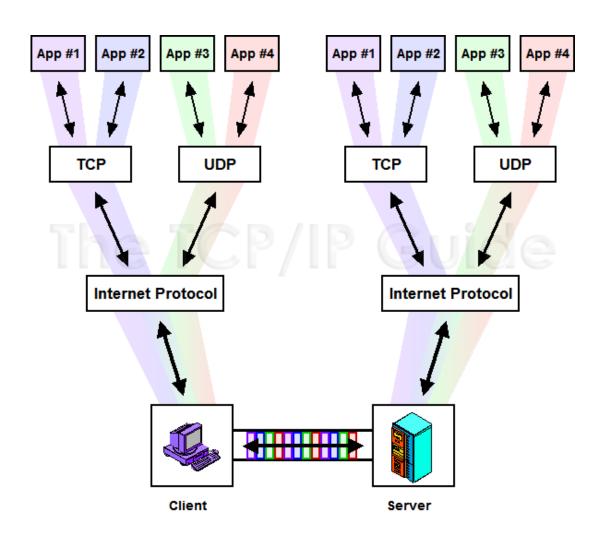




- Analogous to a Data Link layer service
- De/multiplexing of higher layer services (of rpmsg channels)
 - A remote creates an rmpsg channel by binding to a host provided service (identified by a string, the name of the channel)
 - The channel is then identified in the 2 directions by dynamically created numerical endpoints
 - The channel (host side) identifies also the remote we are communicating with
 - There may multiple active channels on a same vring pair
- Tx side provides reliable/flow-controlled and unreliable/best-effort services
- Rx side expects that when a message is received it is immediately extracted from the circular buffer (is dealt with by higher layer SW)
- Reliability of an rpmsg based transport (channel) service depends on the support of flow-control by the transport protocol!

De/multiplexing









```
struct rpmsg driver {
   struct device driver drv;
   const struct rpmsg device id *id table;
   int (*probe) (struct rpmsg channel *dev);
   void (*remove) (struct rpmsg channel *dev);
   void (*callback) (struct rpmsg channel *, void *data,
                   int len, void *priv, u32 src);
};
static struct rpmsg device id rpmsg ipcproto id table[] = {
   { .name = RPMSG PROTO CHANNEL ID }, // e.g. "rpmsg-ipcproto"
  { },
};
MODULE DEVICE TABLE (rpmsg, rpmsg ipcproto id table);
static struct rpmsg driver rpmsg ipcproto driver = {
   .drv.name = KBUILD MODNAME,
   .id table = rpmsg ipcproto id table,
   .probe = rpmsg ipcproto probe, // when channel created
   = rpmsg ipcproto remove,
   .remove
                                                           21
};
```



- Registers an rpmsg driver with the rpmsg bus.
- User should provide a pointer to an rpmsg_driver struct,
 which contains
 - the driver's ->probe() and ->remove() functions,
 - an rx callback, and
 - an id_table specifying the names of the channels this driver is interested to be probed with (e.g. "rpmsq-ipcproto").



- Sends a message across to the remote processor on a given channel.
- The caller should specify the channel, the data it wants to send and its length (in bytes).
- The message will be sent on the specified channel, i.e. its source and destination address fields will be set to the channel's src and dst addresses (endpoints).
- In case there are no TX buffers available, the function will block until
 one becomes available (i.e. until the remote processor consumes a tx
 buffer and puts it back on virtio's used descriptor ring), or a timeout of
 15 seconds elapses.
- When the latter happens, -ERESTARTSYS is returned.



- Sends a message across to the remote processor on a given channel.
- The caller should specify the channel, the data it wants to send, and its length (in bytes).
- The message will be sent on the specified channel, i.e. its source and destination address fields will be set to the channel's src and dst addresses.
- In case there are no TX buffers available, the function will immediately return -ENOMEM without waiting until one becomes available.



- Passes over to the client driver associated to the channel rpdev the data that have been received by the bus driver, along with opaque structure priv.
- The callback function must dispatch data so that following messages in the vring can be processed.

Message style IPC service



- Analogous to UDP
 - Users can create their own transport communication endpoints (service access points, analogous to UDP ports)
 - Service access point identified by a numbered port
 - Best-effort
 - Transport protocol rpmsg dgproto similar to UDP
- Implemented by a (client driver) protocol module
- On rpmsg channel "rpmsg-ipcproto"
- Interfaced via the socket (system call) API
 - Address family AF RPMSG
 - Socket type SOCK_DGRAM

Message style IPC service



Socket layer				
AF_INET protocol stack	AF_SYSTEM protocol stack	AF_RPMSG protocol stack		
TCP, UDP,		<pre>IPC messaging rpmsg_dgproto</pre>		
subnetworks		rpmsg		

Why an additional transport layer?



- Multiple application dialogues between two cores
- But why not simply using different rpmsg channels?
 - Users must be able to create their communication end-points
 - Different dialogues may have different requirements: e.g. tcp supports reliable, stream based communications while udp supports best-effort, message based communications
 - Without an additional transport layer protocol rpmsg supports only best-effort, message based communications
 - If we want a SOCK_SEQPACKET semantic we need a complex connection oriented protocol that implements flow control

rpmsg IPC: core address



- In our IPC autonomous cores on a chip are addressed via a integer identifier
- When the remote requests the creation of a channel the client driver gets a reference to the channel, and the channel allows to locate the description of the remote in the Device Tree
- The definition of an alias in the Device Tree allows us to associate an rpmsg core address to a remote
- The transport protocol rpmsg_dgproto keeps the association rpmsg channel ↔ core address
 - Tx side: core address → rpmsg channel
 - Rx side: rpmsg channel → core address

Device Tree



```
aliases {
        ethernet0 = &gem0;
        serial0 = &uart1;
        spi0 = &qspi;
        rproc1 = &remoteproc1;  // remote core #1
    };
remoteproc1: remoteproc@1F000000 {
        compatible = "xlnx,zynq remoteproc";
        reg = < 0x1F000000 0x1000000 >;
        interrupt-parent = <&intc>;
        interrupts = < 0 37 0 0 38 0 >;
        firmware = "firmware.elf";
        ipino = <0>;
        vring0 = <15>;
        vring1 = <14>;
    };
};
```

rpmsg IPC messaging: open issues



- Performance / functionalities
 - Reliable communications
 - Priority and guaranteed bandwidth
 - Max size of messages
 - 0-copy (now 2 copies Linux side)
 - OCM vs. DDR (what shared memory?)
 - Caching?
 - Routing
- API
 - Implementing SOCK_SEQPACKET communications
 - Extended semantics

rpmsg on RTOSs



- Focus on FreeRTOS but other RTOSs may be relevant
 - > SYSBIOS for TI chips
- Portable implementation provided by OpenAMP
 - Port available for FreeRTOS Xilinx/Freescale
 - Only rpmsg bus
 - No link with upstream Linux community (port of bus driver on Linux is in user space!)
- How can we work upstream for FreeRTOS?

Multicore Association



- http://www.multicore-association.org/index.php
- MCAPI: the Multicore Communications API specification defines an API and a semantic for communication and synchronization between processing cores in embedded systems
- OpenAMP: an open source framework that allows operating systems to interact within a broad range of complex homogeneous and heterogeneous architectures and allows asymmetric multiprocessing applications to leverage parallelism offered by the multicore configuration

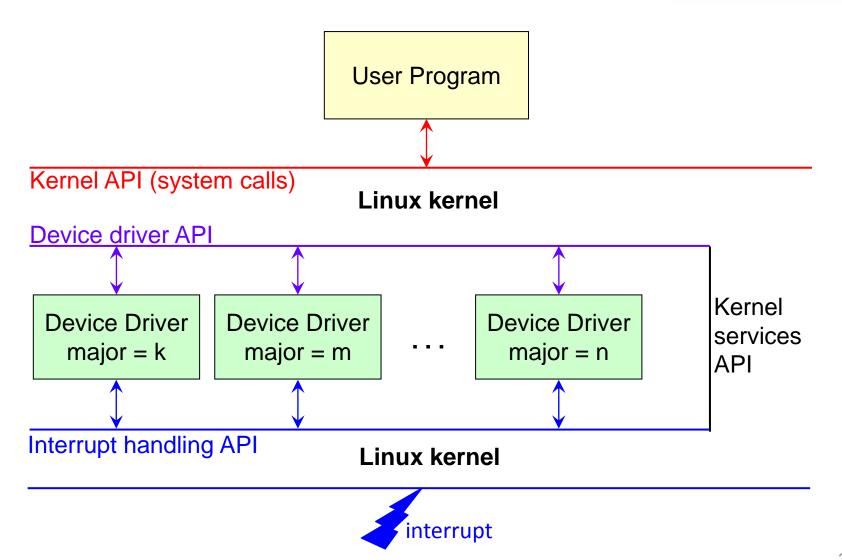
rpmsg API in user space



- Master side
- Via a character driver
- 1 major number and N minor numbers
- Each minor number associated to a pair [remote core, rpmsg service]
- Pair [remote core, rpmsg service] associated to a filename in /dev
- Service is best effort

User space, kernel & drivers





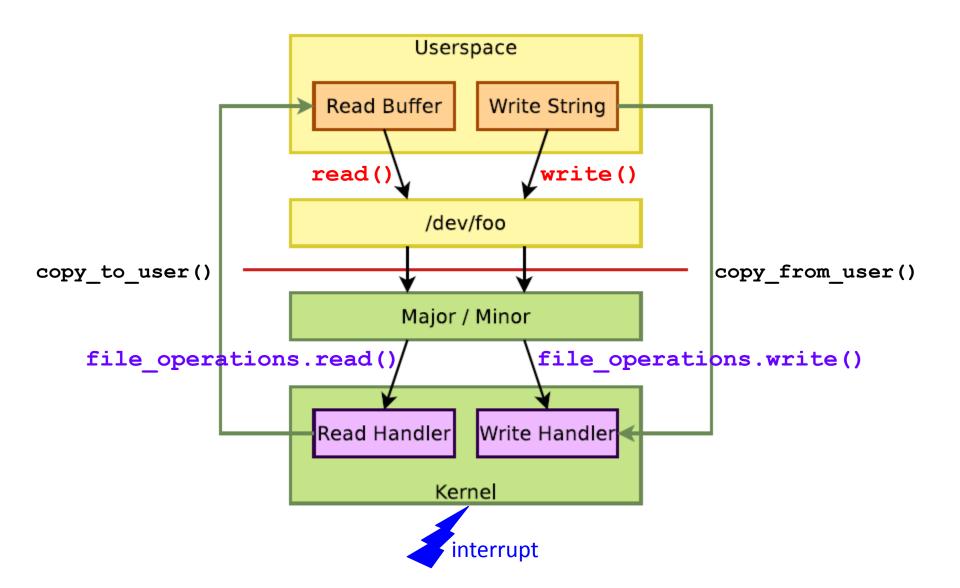
Character drivers



```
struct file operations {
    struct module *owner;
   loff t (*llseek) (struct file *, loff t, int);
   ssize t (*read) (struct file *, char user *, size t, loff t *);
   ssize t (*write) (struct file *, const char user *, size t, loff t *);
   ssize t (*aio read) (struct kiocb *, const struct iovec *, unsigned long, loff t);
   ssize t (*aio write) (struct kiocb *, const struct iovec *, unsigned long, loff t);
   int (*iterate) (struct file *, struct dir context *);
   unsigned int (*poll) (struct file *, struct poll table struct *);
   long (*unlocked ioctl) (struct file *, unsigned int, unsigned long);
   long (*compat ioctl) (struct file *, unsigned int, unsigned long);
   int (*mmap) (struct file *, struct vm area struct *);
   int (*open) (struct inode *, struct file *);
   int (*flush) (struct file *, fl owner t id);
   int (*release) (struct inode *, struct file *);
   int (*fsync) (struct file *, loff t, loff t, int datasync);
   int (*aio fsync) (struct kiocb *, int datasync);
   int (*fasync) (int, struct file *, int);
   int (*lock) (struct file *, int, struct file lock *);
   ssize t (*sendpage) (struct file *, struct page *, int, size t, loff t *, int);
   unsigned long (*get unmapped area) (struct file *, unsigned long, unsigned long,
                                       unsigned long, unsigned long);
   int (*check flags)(int);
   int (*flock) (struct file *, int, struct file lock *);
    ssize t (*splice write) (struct pipe inode info *, struct file *, loff t *, size t,
                            unsigned int);
   ssize t (*splice read) (struct file *, loff t *, struct pipe inode info *, size t,
                           unsigned int);
   int (*setlease)(struct file *, long, struct file lock **);
   long (*fallocate)(struct file *file, int mode, loff t offset, loff t len);
   int (*show fdinfo) (struct seq file *m, struct file *f);
};
```

From user space to char drivers





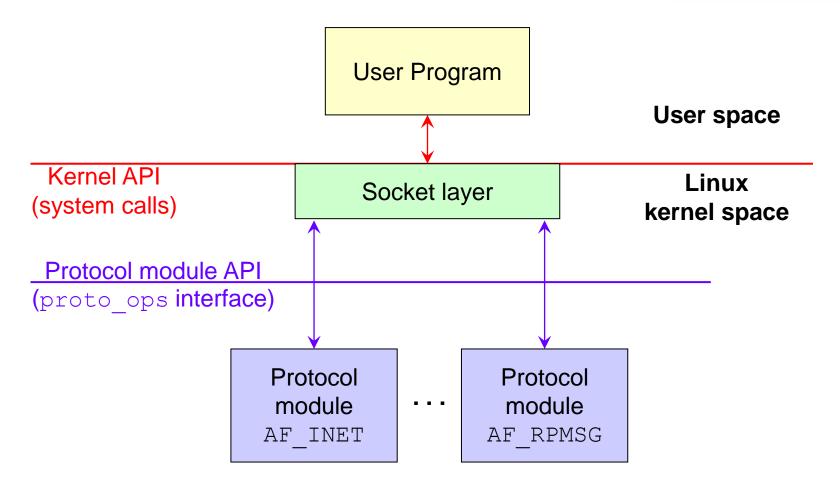
Network protocol



```
static const struct proto ops rpmsg sock ops = {
                   = PF RPMSG,
    .family
                   = THIS MODULE,
    .owner
    .release
                   = rpmsg sock release,
                   = rpmsq sock connect,
    .connect
                   = rpmsg sock getname,
    .getname
    .sendmsq
                   = rpmsg sock sendmsg,
    .recvmsq
                   = rpmsg sock recvmsg,
    .bind
                   = rpmsq sock bind,
                   = sock no poll,
    .poll
    .listen
                   = sock no listen,
    .accept
                   = sock no accept,
    .ioctl
                   = sock no ioctl,
                   = sock no mmap,
    .mmap
    .socketpair
                   = sock no socketpair,
    .shutdown
                   = sock no shutdown,
    .setsockopt
                   = sock no setsockopt,
    .getsockopt
                   = sock no getsockopt
};
```

Network protocol





Programming model



Pthreads

- E.g. manager-worker model: manager on Linux, workers on RTOSs
- pthread_attr_setaffinity_np() allows to control on what core a thread is created/run
- How can we share data between threads working on different cores?
- RPC
- OpenMP

Platform driver / device



- On embedded systems, devices are often not connected through a bus allowing enumeration, hotplugging, and providing unique identifiers for devices.
- However, we still want the devices to be part of the device model.
- The solution to this is the platform driver / platform device infrastructure.
- The platform devices are the devices that are directly connected to the CPU (e.g. memory mapped devices), without any kind of bus.
- https://www.kernel.org/doc/Documentation/drivermodel/platform.txt

Platform driver / device



```
struct platform_device {
                *name;
   const char
   u32
                     id;
    struct device dev;
   u32
                     num resources;
    struct resource *resource;
};
struct platform driver {
    int (*probe) (struct platform device *);
    int (*remove) (struct platform device *);
    void (*shutdown) (struct platform device *);
    int (*suspend) (struct platform device *, pm message t state);
    int (*suspend late) (struct platform device *,
                        pm message t state);
    int (*resume early) (struct platform device *);
    int (*resume) (struct platform device *);
    struct device driver driver;
};
```

.1/9



• From S. A. Edwards, Device Drivers, Columbia University, http://www.cs.columbia.edu/~sedwards/classes/2014/4840/device-drivers.pdf

Module's API:

```
#ifndef VGA LED H
#define VGA LED H
#include <linux/ioctl.h>
#define VGA LED DIGITS 8
typedef struct {
 unsigned char digit; // 0, 1, .. , VGA LED DIGITS-1
 unsigned char segments; // LSB: segment a; MSB: decimal point
} vga led arg t;
#define VGA LED MAGIC 'q'
// ioctls and their arguments
#define VGA LED WRITE DIGIT IOW(VGA LED MAGIC, 1, vga led arg t*)
#define VGA LED READ DIGIT IOWR(VGA LED MAGIC, 2, vga led arg t*)
#endif
```

.2/9



• Excerpt of the Device Tree:

```
lightweight_bridge: bridge@0xff200000 {
   compatible = "simplebus";
   #address-cells = <1>;
   #size-cells = <1>;
   ranges = < 0x0 0xff200000 0x200000 >;
   vga_led: vga_led@0 {
      compatible = "altr,vga_led";
      reg = <0x0 0x8>;
   };
};
```

.3/9



• <u>Driver source code – part 1:</u>

```
#include <linux/module.h>
#include <linux/init.h>
#include <linux/errno.h>
#include <linux/version.h>
#include <linux/platform device.h>
#include <linux/miscdevice.h>
#include <linux/io.h>
#include <linux/of.h>
#include <linux/of address.h>
#include <linux/fs.h>
#include <linux/uaccess.h>
#include "vga led.h"
#define DRIVER NAME "vga led"
struct vga led dev {
    struct resource res;
                                 // Resource: our registers
   void iomem *virtbase;
                                 // Pointer to registers
   u8 segments[VGA LED DIGITS];
} dev;
static void write_digit(int digit, u8 segments) {
    iowrite8(segments, dev.virtbase + digit);
    dev.segments[digit] = segments;
```

.4/9



• <u>Driver source code – part 2:</u>

```
static long vga led ioctl (struct file *f, unsigned int cmd,
                          unsigned long arg) {
 vga led arg t vla;
  switch (cmd) {
    case VGA LED WRITE DIGIT:
      if (copy from user(&vla, (vga_led_arg_t *) arg,
                         sizeof(vga led arg t))) return -EACCES;
      if (vla.digit > 8) return -EINVAL;
      write digit(vla.digit, vla.segments);
      break;
    case VGA LED READ DIGIT:
      if (copy from user(&vla, (vga led arg t *) arg,
                         sizeof(vga led arg t))) return -EACCES;
      if (vla.digit > 8) return -EINVAL;
      vla.segments = dev.segments[vla.digit];
      if (copy to user((vga led arg t *) arg, &vla,
                       sizeof(vga led arg t))) return -EACCES;
     break;
    default: return EINVAL;
  return 0;
```

.5/9



Driver source code – part 3:

```
static const struct file operations vga led fops = {
    .owner = THIS MODULE,
    .unlocked ioctl = vga_led_ioctl,
};
// we define our module as a misc device, with its minor
// number dynamically assigned
static struct miscdevice vga_led_misc_device = {
    .minor = MISC DYNAMIC MINOR,
    .name = DRIVER NAME,
    .fops = &vga led fops,
};
static int vga led remove(struct platform device *pdev) {
  iounmap(dev.virtbase);
  release mem region (dev.res.start, resource size (&dev.res));
 misc deregister (&vga led misc device);
  return 0;
```

.6/9



Driver source code – part 4:

```
static int init vga led probe(struct platform device *pdev) {
  static unsigned char welcome message[VGA LED DIGITS] = {
                0x3E, 0x7D, 0x77, 0x08, 0x38, 0x79, 0x5E, 0x00};
  int i, ret;
  // Register ourselves as a misc device: creates /dev/vga led
  ret = misc register(&vga led misc device);
  // Find our registers in device tree; verify availability
  ret = of address to resource (pdev->dev.of node, 0, &dev.res);
  if (ret) {
   ret = -ENOENT;
   goto out deregister;
  if (request mem region (dev.res.start, resource size (&dev.res),
                         DRIVER NAME) == NULL) {
    ret = EBUSY;
    goto out deregister;
  // vga led probe() continues
```

.7/9



Driver source code – part 5:

```
// Arrange access to our registers (calls ioremap)
  dev.virtbase = of iomap(pdev->dev.of node, 0);
  if (dev.virtbase == NULL) {
    ret = -ENOMEM;
   goto out release mem region;
  // Display a welcome message
  for (i = 0; i < VGA LED DIGITS; i++) {
   write digit(i, welcome message[i]);
  return 0;
out release mem region:
  release mem region (dev.res.start, resource size (&dev.res));
out deregister:
 misc deregister (&vga led misc device);
 return ret;
```

.8/9



Driver source code – part 6:

```
static const struct of device id vga led of match[] = {
    { .compatible = "altr,vga led" },
    { },
};
MODULE DEVICE TABLE (of, vga led of match);
static struct platform driver vga led driver = {
    .driver = {
               .name = DRIVER NAME,
               .owner = THIS MODULE,
               .of match table = of match ptr(vga led of match),
    },
    .probe = vga led probe,
    .remove = exit p(vga led remove),
};
static int init vga led init(void) {
  pr info(DRIVER NAME ": init\n");
  return return platform driver register (&vga led driver);
static void exit vga led exit(void) {
  platform driver unregister (&vga led driver);
  pr info(DRIVER NAME ": exit\n");
                                                                50
```

.9/9



Driver source code – part 7:

```
module_init(vga_led_init);
module_exit(vga_led_exit);

MODULE_LICENSE("GPL");
MODULE_AUTHOR("Stephen A. Edwards, Columbia University");
MODULE_DESCRIPTION("VGA 7-segment LED Emulator");
```

notes:

- 1. The init function of the module is vga led init().
- 2. Function vga_led_init() registers a platform driver with
 .compatible=="altr,vga_led" and probe function vga_led_probe().
- 3. Because a compatible match is found with device tree node vga_led, function vga_led_probe() is activated (a platform_device struct is filled with data extracted from the device tree and passed to the probe function).
- 4. Function vga_led_probe() registers a miscDevice with the misc subsystem (by invoking misc_register()), thus completing the initialization of the vga_led driver module.
- 5. The only file operation supported by the vga_led module is ioctl().

Miscellaneous Character Drivers



- MiscDevice is a thin layer around character devices.
- The misc driver and all miscDevices are assigned major number 10.
- Minor numbers may be assigned dynamically to miscDevices.
- The misc subsystem automatically creates the special file representing a miscDevice in /dev directory.
- The misc driver exports two functions for user modules to register and unregister (minor numbers must be unique among miscDevices):

```
#include <linux/miscdevice.h>
struct miscdevice {
                     // MISC DYNAMIC MINOR assigns it dynamically
  int minor;
  const char *name; // for humans, will appear in /proc/misc file
  const struct file operations *fops;
  struct miscdevice *next, *prev;
                      // must be cleared before registering
};
int misc register(struct miscdevice *misc);
int misc deregister(struct miscdevice *misc);
                                                             52
```

remoteproc: device tree excerpt



```
remoteproc1: remoteproc@1 {
    compatible = "xlnx,zynq_remoteproc";
    reg = < 0x1F000000 0x1000000 >;
    //reg = < 0x18000000 0x08000000 >;
    interrupt-parent = <&intc>;
    interrupts = < 0 37 0 0 38 0 >;
    firmware = "firmware.elf";
    //status = "disabled";
    ipino = <0>;
    vring0 = <15>;
    vring1 = <14>;
};
```

remoteproc: driver excerpt .1/4



```
/*
* Zyng Remote Processor driver
* Copyright (C) 2012 Michal Simek <monstr@monstr.eu>
* Copyright (C) 2012 PetaLogix
 *
  Based on origin OMAP Remote Processor driver
* Copyright (C) 2011 Texas Instruments, Inc.
* Copyright (C) 2011 Google, Inc.
 *
  This program is free software; you can redistribute it and/or
 * modify it under the terms of the GNU General Public License
* version 2 as published by the Free Software Foundation.
* This program is distributed in the hope that it will be useful,
* but WITHOUT ANY WARRANTY; without even the implied warranty of
* MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE. See the
* GNU General Public License for more details.
 */
```

remoteproc: driver excerpt .2/4



```
/* Match table for OF platform binding */
static const struct of device id zyng remoteproc match[] = {
             { .compatible = "xlnx,zynq remoteproc", },
             { /* end of list */ },
};
MODULE DEVICE TABLE (of, zynq remoteproc match);
static struct platform driver zynq remoteproc driver = {
             .probe = zynq remoteproc probe,
             .remove = zynq remoteproc remove,
             .driver = {
               .name = "zynq remoteproc",
               .of match table = zyng remoteproc match,
             },
};
```

remoteproc: driver excerpt .3/4



remoteproc: driver excerpt .3/4



module platform driver(zynq remoteproc driver); // espande: // module driver(zynq remoteproc driver, platform driver register, // platform driver unregister); // espande: // static int init zynq remoteproc driver init(void) { return platform driver register(&zynq_remoteproc_driver); // **//** } // module init(zynq remoteproc driver init); // // static void exit zynq remoteproc driver exit(void) { // platform driver unregister(&zynq remoteproc driver); **//** } // module exit(zynq remoteproc driver exit); module param(firmware, charp, 0); MODULE PARM DESC (firmware, "Override the firmware image name. Default value in DTS."); MODULE AUTHOR ("Michal Simek <monstr@monstr.eu"); MODULE LICENSE ("GPL v2"); 57 MODULE DESCRIPTION("Zynq remote processor control driver");

Acknowledgements



OpenNext is a project co-funded by



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