Module 7: Introduction to Queueing Theory (Notation, Single Queues, Little's Result)

(Slides based on Daniel A. Reed, ECE/CS 441 Notes, Fall 1995, used with permission)

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Outline of Section on Queueing Theory

- 1. Notation
- 2. Little's Result
- 3. Single Queues
- 4. Solutions for networks of queues Product Form Results (on blackboard, not slides)
- 5. Mean value analysis (if time permits)

Queueing Theory Notation

- Queuing characteristics
 - Arrival process
 - Service time distribution
 - Number of servers
 - System capacity
 - Population size
 - Service discipline
- Each of these is described mathematically
 - Descriptions determine tractability of (efficient) analytic solution
 - Only a small set of possibilities are solvable using standard queueing theory



Arrival Processes

- Suppose jobs arrive at times t_1, t_2, \dots, t_j
 - Random variables $\tau_i = t_i t_{i-1}$ are *inter-arrival times*
 - There are many possible assumptions for the distribution of the τ_i
 - Typical assumptions for the τ_i :
 - Independent
 - Identically distributed
 - Many other possible assumptions:
 - Bulk arrivals
 - Balking
 - Correlated arrivals
- For *Poisson* arrival, the inter-arrival times are:
 - IID (independent and identically distributed)
 - exponentially distributed (i.e., CDF $F(x) = 1 e^{-x/a}$)
- Other common arrival time distributions include
 - Erlang, Hyper-exponential, Deterministic, General (with a specified mean and variance)

Other Queue Features

- Service time
 - Interval spent actually receiving service (exclusive of waiting time)
 - As with arrival processes, there are many possible assumptions
 - Most common assumptions are
 - IID random variables
 - exponential service time distribution
- Number of servers
 - Servers may or may not be identical
 - Service discipline determines allocation of customers to servers
- System capacity
 - Maximum number of customers in the system (including those in service)
 - May be finite or infinite
- Population size
 - Total number of potential customers
 - May be finite or infinite

Other Queue Features (Continued)

- Service discipline
 - The order waiting customers are serviced
 - Many possibilities, including
 - First-come-first-serve (FCFS), the most common
 - Last-come-first-serve (LCFS)
 - Last-come-first-serve preempt resume (LCFS-PR)
 - Round robin (RR) with finite quantum size
 - Processor sharing (PS) --- RR with infinitesimal quantum size
 - Infinite server (IS)
 - Almost anything might be used, depending on the the total state of the queue
 - As expected, service discipline affects the nature of the stochastic process that represents the behavior of the queueing system

Queueing Discipline Specification

- Queueing discipline is typically specified using Kendall's notation (A/S/m/B/K/SD), where
 - Letters correspond to six queue attributes
 - *A*: interarrival time distribution
 - *S*: service time distribution
 - *m*: number of servers
 - *B*: number of buffers (system capacity)
 - *K*: population size
 - *SD*: service discipline
- Interarrival and service time specifiers
 - *M* exponential
 - E_k Erlang with parameter k
 - H_k hyperexponential with parameter k
 - *D* deterministic
 - *G* general (any distribution, mean and variance used in the solution)
- Bulk arrivals or service are denoted by a superscript
 - $M^{[x]}$ denotes exponential arrivals with group size x
 - -x is generally a random variable with separately specified distribution
- Omitted specifiers assume certain defaults
 - infinite buffer capacity
 - infinite population size
 - FCFS service discipline

Example Queueing Discipline Specifications

- *M/D/5/40/200/FCFS*
 - Exponentially distributed interarrival times
 - Deterministic service times
 - Five servers
 - Forty buffers (35 for waiting)
 - Total population of 200 customers
 - First-come-first-serve service discipline
- *M/M/*1
 - Exponentially distributed interarrival times
 - Exponentially distributed service times
 - One server
 - Infinite number of buffers
 - Infinite population size
 - First-come-first-serve service discipline

An Introductory Example

- Given these descriptions, what are examples of their application?
- Consider a typical bank
 - 5 tellers
 - Customers form a single line and are serviced FCFS
 - Excluding a run on the bank, the waiting room is effectively infinite
 - For a large bank, the population is effectively infinite
 - Bulk arrivals are possible if friends arrive together for service
- What about service time and inter-arrival time distributions?
 - We can go measure them with a watch at the bank
 - Or, we can make mathematically simplifying assumptions
 - Latter is most common and exponential distribution is typical
- Combining these facts and assumptions
 - M/M/1 queue
 - As we shall see, the mean queue length (including one in service) for an M/M/1 queue is

$$\frac{\lambda}{\mu - \lambda}$$

– Where μ

- λ is the mean inter-arrival time
- μ is the mean service time

Notation and Basic "Facts"

- Standard variable names
 - τ is job interarrival time
 - $\lambda = 1/E[\tau]$ mean job arrival rate
 - *s* is service time per customer (job)
 - *m* is number of servers
 - $\mu = 1/E[s]$ is mean service rate per server
 - $n = n_q + n_s$ is number of jobs in the system
 - n_q is number of jobs waiting to receive service
 - n_s is number of jobs in service
 - *r* is response time (service time plus queueing delay)
 - w is waiting time (queueing delay only)
- System must be "stable" to have an interesting steady state solution
 - Number of jobs in the system is finite
 - Requires the relation $\lambda < m\mu$ hold unless
 - the population is finite (queue length is bounded)
 - the buffer capacity is finite (arrivals are lost when queue is full)
 - (in these cases, system is always stable)

Notation and Basic "Facts"

• Number of jobs in the system

 $-n = n_q + n_s$ (jobs are either waiting or in service)

 $E[n] = E[n_q] + E[n_s]$ (or $\overline{n} = \overline{n}_q + \overline{n}_s$)

- and, if the service rates are independent of queue length
 - $Cov(n_q, n_s) = 0$
 - $Var[n] = Var[n_q] + Var[n_s]$
- Number and time
 - -r = w + s (response time is the sum of queueing delay and service)
 - but, *r*, *w*, and *s* are random variables, so $\overline{r} = \overline{w} + \overline{s}$
 - and, if the service rates are independent of queue length
 - Cov(w,s) = 0
 - Var[r] = Var[w] + Var[s]

Little's Law

- Very important result -- Part of the queueing folk literature for the past century
- Formal proof due to J. D. C. Little (1961)
- Relates mean queue length to arrival rate and mean response time
- Mathematically (in seady state),

 $\overline{n} = \lambda \overline{r}$

- Applies to any "black box" queue under the following assumptions
 - System is work conserving
 - Number of jobs entering is same as number leaving (system is stable)
- Also applies to any transient interval, without requirement that system be stable.
- Note that these are very general conditions, and can apply for any system ("black box") in which customers leave and enter subject to the above constraints.
- An intuitive proof...

Little's Law (Continued)

- Sketch of proof (of steady-state case):
 - During a long interval, arrivals \approx departures (else no stability)
 - Area under the curve is total job time units (jobs x time)
 - Mean queue length \overline{n} is average curve height (area/time)
 - Mean time in system \bar{r} is area/arrivals
 - Mean arrival rate is arrivals/time
- A very general result:
 - No assumptions about arrival or service processes
 - Holds for any queueing discipline (simply charge the area differently)



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jobs x time

iobs

Avg time in system

arrival

rate

jobs

time

jobs x time

time

Avg number in system

Analysis of Single Queues

- Plan:
 - Start with one of the simplest queues, an M/M/1
 - Model as a "birth-death" process
 - Generalize result to other types of queues



- A *birth-death process* is a Markov process in which states are numbered a integers, and transitions are only permitted between "neighboring" states.
- Steady state solution of a birth death process (Kleinrock, Queueing Systems, vol. 1):
 - (*Theorem*) steady state probability p_n of being in state *n* is

$$p_n = \frac{\lambda_0 \lambda_1 \dots \lambda_{n-1}}{\mu_1 \mu_2 \dots \mu_n} p_0 \qquad n = 1, 2, \dots, \infty$$

- where p_0 is the probability of being in state 0
- Now for a proof ...

Birth-Death (Steady-State) State Occupancy Proof

• If stable, in the steady state (by Markov process solution described earlier) $0 = \lambda_{j-1} p_{j-1} - (\mu_j + \lambda_j) p_j + \mu_{j+1} p_{j+1}$ Flow balance at state j or

$$p_{j+1} = \left(\frac{\mu_j + \lambda_j}{\mu_{j+1}}\right) p_j - \frac{\lambda_{j-1}}{\mu_{j+1}} p_{j-1} \qquad j = 1, 2, 3, \dots$$

and

$$p_1 = \frac{\lambda_0}{\mu_1} p_0$$

• And the solution is...

$$p_{n} = \frac{\lambda_{0}\lambda_{1}...\lambda_{n-1}}{\mu_{1}\mu_{2}...\mu_{n}} p_{0}$$

= $p_{0} \prod_{j=0}^{n-1} \frac{\lambda_{j}}{\mu_{j+1}}$ $n = 1, 2, ..., \infty$

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Birth-Death (Steady-State) State Occupancy Proof, cont.

• Finally, because

$$\sum_{j=0}^{\infty} p_j = 1$$

we have

$$p_0 = \frac{1}{1 + \sum_{n=1}^{\infty} \prod_{j=0}^{n-1} \frac{\lambda_j}{\mu_{j+1}}}$$

M/M/1 Queue Analysis

- M / M / 1 is a special case of a birth death process
 - $\lambda_i = \lambda_j$ for all *i* and *j*
 - $\mu_i = \mu_j$ for all *i* and *j*



• By simplification

$$p_n = \left(\frac{\lambda}{\mu}\right)^n p_0 \qquad n = 1, 2, \dots, \infty$$

• By tradition, the ratio

$$\rho = \frac{\lambda}{\mu}$$

is called the "traffic intensity" and

$$p_n = \rho^n p_0$$

and

$$p_0 = \frac{1}{1 + \rho + \rho^2 + \dots + \rho^\infty} = 1 - \rho$$

• By substitution

$$p_n = (1 - \rho)\rho^n \qquad n = 0, 1, 2, \dots \infty$$

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M/M/1 Queue Analysis (Continued)

- Utilization *U* is simply $1 p_0 = \rho$
- Mean queue length E[n] (or \overline{n})

$$\overline{n} = \sum_{n=1}^{\infty} np_n$$
$$= \sum_{n=1}^{\infty} n(1-\rho)\rho^n$$
$$= \frac{\rho}{1-\rho}$$

"almost" mean of a geometric random variable---factor out a rho first

$$Var[n] = E[n^{2}] - (E[n])^{2}$$
$$= \left(\sum_{n=1}^{\infty} n^{2} (1-\rho) \rho^{n}\right) - (E[n])^{2}$$
$$= \frac{\rho}{(1-\rho)^{2}}$$

• Probability of *n* or more jobs in the system

$$\sum_{j=n}^{\infty} p_j = \sum_{j=n}^{\infty} (1-\rho)\rho^j = \rho^n$$

M/M/1 Queue Analysis (Continued)

• Mean response time \overline{r} (or R) via Little's Law

$$\overline{n} = \lambda \overline{r}$$

yields

$$\overline{r} = \frac{\overline{n}}{\lambda} = \frac{\rho}{(1-\rho)\lambda} = \frac{1}{\mu-\lambda}$$

where the response time approaches ∞ as $\lambda \rightarrow \mu$

• CDF of response time is

$$F(r) = 1 - e^{-r\mu(1-\rho)}$$

• Mean number of jobs in the queue $E[n_q]$ (or \overline{n}_q)

$$\overline{n}_q = \sum_{n=1}^{\infty} (n-1)p_n = \frac{\rho^2}{1-\rho}$$

M/M/1 Queue Example

- Consider the following queue
 - $\lambda = 0.3$
 - $-\mu = 0.5$
- We can calculate the following statistics
 - utilization \boldsymbol{U}

$$U = \rho = \frac{\lambda}{\mu} = \frac{0.3}{0.5} = 0.6$$

- mean number of jobs in the system \overline{n}

$$\overline{n} = \frac{\rho}{1 - \rho} = \frac{0.6}{0.4} = 1.5$$

- mean response time \overline{r}

$$\overline{r} = \frac{1}{\mu - \lambda} = \frac{1}{0.2} = 5.0$$

M/M/1 Queue Example (Continued)

- Consider changing λ
 - hold μ fixed at 0.5
 - examine changes in performance metrics



M/M/m Queues

- M/M/m queues
 - *m* servers rather than one server
 - Reasonable model of
 - a bank queue with multiple tellers
 - a shared memory multiprocessor
- Assumptions
 - *m* servers
 - All servers have the same service rate μ
 - Single queue for access to the servers
 - Arrival rate λ
 - Formally

$$\begin{split} \lambda_n &= \lambda & n = 0, 1, \dots, \infty \\ \mu_n &= \begin{cases} n\mu & n = 0, 1, \dots, m-1 \\ m\mu & n = m, m+1, \dots \infty \end{cases} \end{split}$$

• What are the state occupancy probabilities?



- State occupancy probabilities
 - Just another birth death process
 - Recall general form of the probability occupancies earlier

$$p_n = \frac{\lambda_0 \lambda_1 \dots \lambda_{n-1}}{\mu_1 \mu_2 \dots \mu_n} p_0 \qquad n = 1, 2, \dots, \infty$$

• By simple substitution of the λ_j and μ_j , we have

$$p_{n} = \begin{cases} \frac{\lambda^{n}}{n!\mu^{n}}p_{0} & n = 1, 2, ..., m-1 \\ \frac{\lambda^{n}}{m!m^{n-m}\mu^{n}}p_{0} & n = m, m+1, ..., \infty \end{cases}$$

or equivalently (with $\rho = \lambda/(m\mu)$)

$$p_{n} = \begin{cases} \frac{(m\rho)^{n}}{n!} p_{0} & n = 1, 2, ..., m - 1 \\ \frac{\rho^{n} m^{m}}{m!} p_{0} & n = m, m + 1, ..., \infty \end{cases}$$

• And, because

$$\sum_{n=0}^{\infty} p_n = 1$$

we have

$$p_0 = \left[1 + \frac{(m\rho)^m}{m!(1-\rho)} + \sum_{n=1}^{m-1} \frac{(m\rho)^n}{n!}\right]^{-1}$$



- In a similar manner to that for the M/M/1 queue,
 - We can derive the "standard" measures

(queue length, utilization, response time, etc.)

- You should do these derivations yourself
- Mean number of jobs in the system $\overline{n} = \overline{n}_q + \overline{n}_s$

$$\overline{n} = m\rho + \frac{\rho\varsigma}{1-\rho}$$

where

$$\rho = \frac{\lambda}{m\mu}$$

$$\varsigma = P(\ge m \text{ jobs}) = \sum_{n=m}^{\infty} p_n = \frac{(m\rho)^m}{m!(1-\rho)} p_0$$

observe that ς is

- the probability an arriving job must queue

- also known as Erlang's C formula

• Expected number of jobs in service \overline{n}_s

$$\overline{n}_s = \sum_{n=1}^{m-1} np_n + \sum_{n=m}^{\infty} mp_n = m\rho$$

• Utilization of each server

- *m* servers
- $m\rho$ mean jobs in service
- individual server utilization must be ho
- Mean response time $\overline{r} = \overline{w} + \overline{s}$ (just apply Little's law)

$$\overline{r} = \frac{\overline{n}}{\lambda}$$
$$= \frac{1}{\mu} \left(1 + \frac{\varsigma}{m(1 - \rho)} \right)$$

• Mean waiting time \overline{w} (Little's law again)

$$\overline{w} = \frac{\overline{n}_q}{\lambda}$$

$$= \frac{\overline{n} - \overline{n}_s}{\lambda}$$

$$= \frac{\rho}{m\mu(1 - \rho)}$$
• r_q (q percentile of waiting time)
$$r_q = \max\left(0, \frac{\overline{w}}{\varsigma} \ln \frac{100\varsigma}{100 - q}\right)$$

M/M/m Queue Example

- Consider changing *m*
 - hold λ and μ fixed
 - examine changes in performance metrics
- Observations
 - M/M/m queue has asymptote at $\frac{\lambda}{m\mu}$
 - substantial performance gains with even two servers



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M/M/1 and M/M/m Queue Comparison

- Which is better?
 - *m* queues each with an arrival rate λ/m
 - one queue with *m* servers and an arrival rate of λ
- Suppose we use mean response time as our metric...
 - m M/M/1 queues

$$\bar{r} = \frac{1}{\mu - \lambda/m}$$

- one M/M/m queue

$$\bar{r} = \frac{1}{\mu} \left(1 + \frac{\varsigma}{m(1-\rho)} \right)$$

where

$$\varsigma = \frac{\left(\lambda / \mu\right)^m}{m! \left(1 - \lambda / (m\mu)\right)} p_0$$

and

$$p_{0} = \left[1 + \frac{(\lambda / \mu)^{m}}{m!(1 - \lambda / (m\mu))} + \sum_{n=1}^{m-1} \frac{(\lambda / \mu)^{n}}{n!}\right]^{-1}$$

Queueing Comparison

- Consider the following
 - service rate μ fixed at 4, divided evenly among *m* servers
 - fix $\lambda = 2$
 - m M/M/1 queues (arrival rate to each is λ/m)
 - One M/M/m queue (total arrival rate is λ)
 - Increase *m*
- What happens to response time in both queues? Why?

Mean Response Time as function of m



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Queueing Comparison

- Consider the following
 - service rate μ fixed at 2.1, divided evenly among *m* servers
 - varying λ (subject to stability constraint)
 - m M/M/1 queues (arrival rate to each is λ/m)
 - One M/M/m queue (total arrival rate is λ)
- What happens as λ approaches 2.1? Why?

Mean Response Time as a Function of Arrival Rate



Queueing system with Maximum Arrival Rate 2, total service rate 2.1, 4 servers

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Extrapolation Scenarios

- Given queueing formulae, standard questions include
 - Performance measures for different parameters
 - Parameters values needed to satisfy a particular performance constraint
- Examples:
 - What is the mean response time if arrival rate doubles?
 - What is the mean queue length if service rate decreases by one third?
 - What is the number of servers for mean response time less than five minutes?
- Approach:
 - Plug and crank
 - Repeated solution with different parameter values

Extrapolation Scenarios (Continued)

- Concrete example
 - multiprocessor system (two processors)
 - mean job service time is 15 seconds
 - mean job interarrival time is 12 seconds
- By inspection
 - mean service rate is 4.0 jobs/minute (per processor)
 - mean arrival rate is 5.0 jobs/minute

and by plug and crank, we have mean response time \overline{r}

$$\overline{r} = \frac{1}{\mu} \left[1 + \frac{\varsigma}{m(1-\rho)} \right] = 0.41 \text{ minutes (24.6 seconds)}$$

• How many processors do we need to have $\bar{r} < 0.3$ minutes?

- solve for m = 3, 4, ...
- find smallest value of *m* such that $\overline{r} < 0.3$
- here, m = 3 satisfies this constraint

M/M/m/B Queues

- Finite buffers
 - no more than *B* jobs in total can be
 - queued
 - *and* in service

(i.e., total number of jobs in the system must be less than *B*)

- jobs arriving when *B* jobs are present are discarded
- More formally, this implies

$$\lambda_n = \lambda$$
 $n = 1, 2, ..., B-1$

and

$$\mu_n = \begin{cases} n\mu & n = 1, 2, \dots, m-1 \\ m\mu & n = m, m+1, \dots, B \end{cases}$$

- Observations
 - $-B \ge m$ or servers are wasted
 - birth-death process
 - finite number of states



• Applying the state occupancy formula

$$p_{n} = \begin{cases} \frac{\lambda^{n}}{n! \mu^{n}} p_{0} & n = 1, 2, ..., m-1 \\ \frac{\lambda^{n}}{m! m^{n-m} \mu^{n}} p_{0} & n = m, m+1, ..., B \end{cases}$$

• And, because $\rho = \frac{\lambda}{m\mu}$

$$p_{n} = \begin{cases} \frac{(m\rho)^{n}}{n!} p_{0} & n = 1, 2, ..., m-1 \\ \frac{\rho^{n} m^{m}}{m!} p_{0} & n = m, m+1, ..., B \end{cases}$$

• Finally, the probability of zero jobs in the system is

$$p_{0} = \sum_{n=0}^{B} p_{n} = \left[1 + \frac{\left(1 - \rho^{B-m+1}\right)(m\rho)^{m}}{m!(1-\rho)} + \sum_{n=1}^{m-1} \frac{(m\rho)^{n}}{n!}\right]^{-1}$$

- Now, we can use the state occupancy probabilities to compute
 - mean response time
 - mean queue lengths
 - effective arrival rates

• Mean queue length \overline{n} (queue plus service)

$$\overline{n} = \sum_{n=1}^{B} n p_n$$

and mean number in the queue

$$\overline{n}_q = \sum_{n=m+1}^{B} (n-m) p_n$$

• Arrivals are constrained by waiting space

- effective arrival rate $\widetilde{\lambda}$ is less than λ

- jobs enter the system only when buffers are available

$$\widetilde{\lambda} = \sum_{n=0}^{B-1} \lambda p_n = \lambda (1 - p_B)$$

and the difference $\lambda - \tilde{\lambda}$ is the loss rate

- Because jobs are not lost after entry, the mean response time is

$$\overline{r} = \frac{\overline{n}}{\widetilde{\lambda}} = \frac{\overline{n}}{\lambda(1 - p_B)}$$

by Little's law

- Finally, the utilization U of each server is

$$U = \frac{\widetilde{\lambda}}{m\mu} = \rho(1 - p_B)$$

Other Queues

- Other queues can be solved to varying degrees...
- Exact solutions are possible for
 - $M/E_r/1$ (Erlangian service)
 - M/D/1 (special case of M/G/1)
 - M/M/1 with bulk arrivals (restricted cases)
- Analysis is more difficulty for:
 - *G/M*/1
 - M/G/1
 - *G*/*G*/1

M/G/1 Queues

- *M/G/*1
 - General service time distribution
 - Otherwise, similar to M/M/1 queues
 - The most complex, readily solvable single queue
- Solution approach
 - First, some additional mathematical machinery
 - Then, comparisons with M/M/1 queues
- Service time distribution is general
 - Service history matters
 - Denote service time already received by $X_0(t)$
- Arrival distribution is negative exponential
 - Arrival history does not matter
 - But we do need to know the number of customers N(t) present
 - N(t) is non-Markovian because it depends on service time
- State-space description
 - States are $[N(t), X_0(t)]$
 - Mixed discrete/continuous, two-dimensional description
 - Analysis via this method (supplementary variables) is ugly
 - Use the method of embedded Markov chains...

- What has changed from M/M/1?
 - Two-dimensional state space
 - State space is now continuous (due to $X_0(t)$)
- Ideally
 - Convert $[N(t), X_0(t)]$ to one-dimensional N(t)
 - Implicitly specify remaining service duration $X_0(t)$
- How do we do this?
 - Look only at selected points in time
 - Compute new metrics only at those points
 - Choose those points to implicitly carry $X_0(t)$
 - departures instants make great choices
 - Remaining (residual) service $X_0(t)$ is zero!
 - At that instant, we can treat the behavior like a Markov chain
 - N(t) is the number of customers left behind
 - This is an *embedded Markov chain*; for details (see Kleinrock, vol. 1) but we haven't specified the distribution of departure instants

- A informal derivation follows (see Kleinrock vol. 1 for details)...
- Notation
 - Arrival rate λ (Poisson process)
 - General service time distribution
 - mean \bar{x}
 - variance
- What is the expected time until a customer that arrives completes service?
 - Mean time needed to service customers already waiting
 - Mean time is $\bar{n}_q \bar{x}$
 - Note that this is independent of the distribution of *x*
 - *plus* the residual time for customer in service ...
- Residual life requires yet another aside...

Residual Life

- What is a "renewal"?
 - Informally, a point where random variables which describe a model are memoryless given current state, with respect to past state.
- Renewal example
 - Consider a queue with general service distribution, and Poisson arrival process
 - Most time points are not renewal points, since remaining service time depends on service time completed.
 - However, times at which service completes are renewal points, since arrival process is Poisson.
- Need to determine the residual lifetime of a customer in service:
 - Denote this random variable as *R*
 - Distribution of *R* depends on
 - Distribution of original variable *A* (the service time distribution) at its renewal point and some time expended after the renewal point

Residual Life (Continued)

- Suppose
 - a(t) is the pdf of A (original variable)
 - the original lifetime has expended time t_e
 - then r(t), the pdf of R (the residual lifetime) is

$$r(\tau - t_e | t_e) = a(\tau | \tau > t_e)$$
$$= \frac{a(\tau)}{P(A > \tau_e)}$$
$$= \frac{a(\tau)}{1 - \int_0^{t_e} a(s) ds}$$

- Intuition
 - in general, knowing about the expended time helps
 - in short, knowledge changes the pdf
 - we saw that this was *not* true for the exponential distribution
 - the geometric distribution is the only case in the discrete domain
- Average residual lifetime \overline{r} (claim without proof)
 - depends only on the first two moments of the original pdf f(x)
 - mean \bar{f}
 - second moment \bar{f}^2 (not the variance!)
 - mean residual lifetime is

$$\bar{r} = \frac{\bar{f}^2}{2\,\bar{f}}$$

Residual Life (Continued)

- Example (computer part)
 - suppose the pdf b(t) of the failure time is uniform
 - and suppose the mean value is 10

$$b(t) = \begin{cases} \frac{1}{20} & 0 < \tau \le 20\\ 0 & \text{otherwise} \end{cases}$$

- if the part has been in use for 5 time units, then

$$b(t+t_e) = \begin{cases} \frac{1}{20} & t_e < t+t_e \le 20\\ 0 & \text{otherwise} \end{cases}$$

and

$$1 - \int_0^5 b(s) ds = 1 - \frac{1}{20} \cdot 5 = 0.75$$

and finally

$$r(\tau - t_e | 5) = \begin{cases} \frac{1}{15} & 0 < \tau + t_e \le 15\\ 0 & \text{otherwise} \end{cases}$$

- notice that

$$\bar{r} = \frac{\bar{f}^2}{2\bar{f}} = \frac{133.33}{2 \cdot 10} = 6.67$$

Observe

- pdf of residual time is not the same as the original pdf
- Knowledge of past behavior changes the pdf
- There are only two exceptions
 - negative exponential distribution (continuous)
 - geometric distribution (discrete)

- How long does a new arrival have to wait for service?
 - mean time needed to service customers already waiting
 - * let \overline{x} denote the mean service time
 - * mean time is $\overline{n}_{q}\overline{x}$
 - * note that this is independent of the distribution of x
 - *plus* the residual time for customer in service
 - * recall that this is

$$\bar{t} = \frac{\bar{x}^2}{2\bar{x}}$$

assuming a customer is in service

* the probability of a customer in service is ρ

• Combining items, the waiting time for a new arrival is

$$\bar{r}_q = \bar{n}_q \bar{x} + \rho \frac{x^2}{2\bar{x}}$$

Little's Law again!

• Expected number of arrivals during this interval is $\lambda \bar{r}$, so

$$\bar{r}_q = \bar{r}_q \lambda \bar{x} + \rho \frac{\bar{x}^2}{2\bar{x}}$$

and by rearranging terms

$$\bar{r}_q = \frac{\lambda \bar{x}^2}{2(1-\rho)}$$

• As we just saw, the mean time to receive service is is

$$\bar{r}_q = \frac{\lambda \bar{x}^2}{2(1-\rho)}$$

• Adding the mean service time yields the mean response time

$$\overline{r} = \overline{x} + \frac{\lambda \overline{x}^2}{2(1-\rho)}$$

• Normally, both \bar{r} and \bar{r}_q are expressed as (verify the math)

$$\bar{r} = \frac{1}{\mu} + \frac{\lambda(1 + C_s^2)}{2\mu^2(1 - \rho)}$$

and

$$\bar{r}_q = \frac{\lambda(1+C_s^2)}{2\mu^2(1-\rho)}$$

where C_s^2 is the coefficient of variation

$$C_s^2 = \frac{\sigma^2}{\bar{x}^2}$$

and

 $-\sigma^2$ is the variance of the mean service time $-\frac{1}{\mu}$ is the mean service time

- and by simplification (yields original formulation): $1 + C_s^2 = 1 + \frac{\sigma^2}{\overline{r}^2} = 1 + \frac{\overline{x}^2 - \overline{x}^2}{\overline{r}^2} = \frac{\overline{x}^2}{\overline{r}^2} = \mu^2 \overline{x}^2$

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• Via Little's law, the mean number in the system is

$$\overline{n} = \lambda \overline{r}$$

$$= \frac{\lambda}{\mu} + \frac{\lambda^2 \left(1 + C_s^2\right)}{2\mu^2 (1 - \rho)}$$

$$= \rho + \frac{\rho^2 \left(1 + C_s^2\right)}{2(1 - \rho)}$$

Observations

- this is the famous Pollaczek Khinchin (PK) formula
- learn it, remember it, treasure it!
- C_s is *one* for the negative exponential distribution, so

$$\overline{n} = \rho + \frac{\rho^2 (1+1)}{2(1-\rho)} = \rho + \frac{\rho^2}{1-\rho} = \frac{\rho}{1-\rho}$$

as we knew before

- C_s is zero for the deterministic distribution (M / D / 1 queue)

$$\overline{n} = \rho + \frac{\rho^2 (1+0)}{2(1-\rho)} = \frac{\rho}{(1-\rho)} \left(\frac{2-\rho}{2}\right)$$

- The value of C_s has profound implications
 - larger C_s increases mean queue length and response time
 - values grow linearly with C_s

Queueing Comparison

- Consider the following
 - M/D/1 queue ($C_s = 0$)
 - M/M/1 queue ($C_s = 1$)
 - M/G/1 queue ($C_s > 1$)



Queueing Example

- Consider the following
 - arrival rate $\lambda = 0.6$
 - service rate $\mu = 1.0$
 - M/D/1, M/M/1, and M/G/1 queues

and compare mean response times

• *M*/*M*/1

$$\bar{r} = \frac{1}{\mu - \lambda} = \frac{1}{1.0 - 0.6} = 2.5$$

• *M*/*D*/1

$$\bar{r} = \frac{1}{\mu} + \frac{\lambda \left(1 + C_s^2\right)}{2\mu^2 \left(1 - \rho\right)} = \frac{1}{1.0} + \frac{0.6(1 + 0)}{2(1.0)(1 - 0.6/1.0)} = 1.75$$

•
$$M/G/1 \ (C_s = 2.0)$$

$$\bar{r} = \frac{1}{\mu} + \frac{\lambda \left(1 + C_s^2\right)}{2\mu^2 (1 - \rho)} = \frac{1}{1.0} + \frac{0.6(1 + 1)}{2(1.0)(1 - 0.6/1.0)} = 3.25$$

Queueing Example (Continued)

- Consider M/M/1 and M/G/1 queues
 - assume same arrival rates for both
 - desire same mean response times
 - must solve for ratio of service rates
- *M*/*M*/1



• *M/G/*1

$$\bar{r} = \frac{1}{\mu_g} + \frac{\lambda \left(1 + C_s^2\right)}{2\mu_g^2 \left(1 - \lambda / \mu_g\right)}$$

• Equating, we have

$$\frac{1}{\mu_m - \lambda} = \frac{1}{\mu_g} + \frac{\lambda \left(1 + C_s^2\right)}{2\mu_g^2 \left(1 - \lambda / \mu_g\right)}$$

• Let's look at some numerical solutions...

M/G/1 via Embedded DTMC

- M/G/1 can be analyzed from the point of view of an *embedded DTMC*
- Note : state can be defined as (n, r) where n is the number in system, and r is remaining time of the job in service.
- Future behavior depends only on n at instants when r = 0—job departures
- The state of the embedded DTMC is the number of jobs in system at the time the last job left service

M/G/1 via Embedded DTMC

- So what is P_{ij} ? The probability that exactly j-i-1 jobs arrived while the last job received service.
- If the arrival process is Poisson, and the service time is x, then the number of arrivals during service is Poisson distributed with mean λx.
- Let f(x) be the pdf for the service time distribution, then for i > 0 and i − 1 ≤ j

$$P_{ij} = \int_0^\infty f(x) \frac{(\lambda x)^{(j-i-1)}}{(j-i-1)!} \exp\{-\lambda x\}$$

and for i = 0, $P_{0,1} = 1$ and $P_{0,n} = 0$ for n > 1.

M/G/1 via Embedded DTMC

- Let $\{\pi_i^*\}$ be the equilibrium state probabilities for the embedded DTMC.
- The mean time (from the M/G/1 queue point of view) the DTMC is in a state i is
 - for i > 1, $1/\mu$, the mean of the general service time distribution.
 - for i == 0, $1/\lambda$ the mean time between arrivals
- From $\{\pi_i^*\}$ and mean occupancy times we can get the stationary distribution of the M/G/1 queue.
 - Define $G = (1/\lambda)\pi_0^* + (1 \pi_0^*)/\mu$ —weighted sum of state probabilities



Queueing Example (Continued)

- Comparison Example (Continued)
 - arrival rate $\lambda = 0.6$
 - M/M/1 queue (service rate $\mu_m = 1.0$)
 - M/G/1 queue (service rate μ_g)

