# Optimization in ComPASS-4 

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## The Plan

1. Optimization Formulations and Taxonomy

- Stochastic Optimization
- Multiobjective Optimization
- Simulation-Based Optimization
- Derivative-Free Optimization
- Global Optimization

2. An Example LPA Optimization to Highlight Challenges
3. POPAS
4. Why not Blackbox Optimization
5. APOSMM

## Mathematical/Numerical Nonlinear Optimization

## Optimization is the "science of better"

Find parameters (controls) $x=\left(x_{1}, \ldots, x_{n}\right)$ in domain $\Omega$ to improve objective $f$

$$
\min \left\{f(x): x \in \Omega \subseteq \mathbb{R}^{n}\right\}
$$

$\diamond$ (Unless $\Omega$ is very special) Need to evaluate $f$ at many $x$ to find a good $\hat{x}_{*}$
$\diamond$ Focus on local solutions: $f\left(\hat{x}_{*}\right) \leq f(x) \forall x \in \mathcal{N}\left(\hat{x}_{*}\right) \cap \Omega$
$\diamond$ constraints defined the feasibility region $\Omega$



## Stochastic Optimization

Addresses situations where you obtain a nondeterministic quantity $F(x, \xi)$

$$
\min \{f(x)=\mathrm{E}\{F(x, \xi)\}: x \in \Omega\}
$$

$\diamond x \in \mathbb{R}^{n}$ decision variables
$\diamond \xi$ vector of random variables

- independent of $x$
- $P(\xi)$ distribution function for $\xi$
- $\xi$ has support $\Xi$
$\diamond F(x, \cdot)$ functional form of uncertainty for decision $x$
$\diamond \Omega \subseteq \mathbb{R}^{n}$ set defined by deterministic constraints
- Also: stochastic/probabilistic constraints
$\diamond$ Nonstationarity: does $\operatorname{Var}\{F(x, \xi)\}$ depend on $x$ ?


## Multiobjective Optimization

Simultaneously minimize $n_{f}>1$ objectives

$$
\min _{x \in \Omega} f_{1}(x), \cdots, f_{n_{f}}(x)
$$

" $x^{1}$ dominates $x^{2}$ " if:
$\diamond f_{i}\left(x^{1}\right) \leq f_{i}\left(x^{2}\right)$ for all $i$, and
$\diamond f_{i}\left(x^{1}\right)<f_{i}\left(x^{2}\right)$ for at least one $i$
" $x^{1}$ is nondominated in $\mathcal{X}$ " if there is no $x^{2} \in \mathcal{X}$ that dominates $x^{1}$


Pareto optimal solutions: A set $\mathcal{P}$ of points are nondominated in $\Omega$
$\diamond$ Especially useful when missing a currency exchange between objectives
$\diamond$ Significantly more expensive than single-objective optimization


## Simulation-Based Optimization

$$
\min _{x \in \mathbb{R}^{n}}\{f(x)=F[\mathbf{S}(\mathbf{x})]: c(\mathbf{S}(\mathbf{x})) \leq 0, x \in \mathcal{B}\}
$$

$\diamond S$ (numerical) simulation output, (here deterministic)
$\diamond$ Derivatives $\nabla_{x} S$ often unavailable or prohibitively expensive to obtain/approximate directly
$\diamond$ Some AD hurdle (e.g., proprietary/legacy/coupled/mixed-language codes)
$\diamond$ Single evaluation of $S$ could take seconds/minutes/hours/days
Evaluation is a bottleneck for optimization
$\mathcal{B}$ compact, known region (e.g., finite bound constraints)

Computing advances have driven this research area.


Argonne's AVIDAC (1953 vacuum tubes)


Argonne's BlueGene/Q (2012 0.79M cores)


Argonne's Theta (2017 0.23M cores)


Sunway TaihuLight (2016 11M cores)

## Derivative-Free/Zero-Order Optimization

"Some derivatives are unavailable for optimization purposes"

## Derivative-Free/Zero-Order Optimization

"Some derivatives are unavailable for optimization purposes"

## The Challenge: Optimization is tightly coupled with derivatives

Typical optimality (no noise, smooth functions)

$$
\nabla_{x} f\left(x^{*}\right)+\lambda^{T} \nabla_{x} c_{E}\left(x^{*}\right)=0, c_{E}\left(x^{*}\right)=0
$$


(sub)gradients $\nabla_{x} f, \nabla_{x} c$ enable:
$\diamond$ Faster feasibility
$\diamond$ Faster convergence

- Guaranteed descent
- Approximation of nonlinearities
$\diamond$ Better termination
- Measure of criticality $\left\|\nabla_{x} f\right\|$ or $\left\|\mathcal{P}_{\Omega}\left(\nabla_{x} f\right)\right\|$
$\diamond$ Sensitivity analysis
- Correlations, standard errors, UQ, ...


## Ways to Get Derivatives

## Handcoding (HC)

"Army of students/programmers"
? Prone to errors/conditioning
? Intractable as number of ops increases

## Algorithmic/Automatic Differentiation (AD)

"Exact* derivatives!"
? No black boxes allowed
? Not always automatic/cheap/well-conditioned

## Finite Differences (FD)

"Nonintrusive"
? Expense grows with $n$
? Sensitive to stepsize choice/noise


Caution
High noise levels
$\rightarrow$ [Moré \& W.; SISC 2011], [Moré \& W.; TOMS 2012]
... then apply derivative-based method (that handles inexact derivatives)

## Algorithmic Differentiation

$\rightarrow$ [Coleman \& Xu; SIAM 2016], [Griewank \& Walther; SIAM 2008]

## Computational Graph

$\diamond y=\sin (a * b) * c$
$\diamond$ Forward and reverse modes
$\diamond$ AD tool provides code for your derivatives

Write codes and formulate problems with AD in mind!


Many tools (see www. autodiff.org):

F OpenAD<br>F/C Tapenade, Rapsodia<br>$\mathrm{C} / \mathrm{C}++\mathrm{ADOL}-\mathrm{C}, \mathrm{ADIC}$

Matlab ADiMat, INTLAB
Python/R ADOL-C

Also done in AMPL, GAMS, JULIA!

## The Price of Algorithm Choice: Solvers in PETSc/TAO



Toolkit for Advanced Optimization
[Munson et al.; mcs.anl.gov/tao]

## Increasing level of user input:

nm Assumes $\nabla_{x} f$ unavailable, black box
pounders Assumes $\nabla_{x} f$ unavailable, exploits problem structure

Imvm Uses available $\nabla_{x} f$

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THIS TALK
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DFO methods should be designed to beat finite-difference-based methods

Observe: Constrained by budget on \#evals, method limits solution accuracy/problem size

## Why Algorithms Matter: The Accelerator Case

Varying skew quadrupoles to meet beam size targets (in PELEGANT)

$\diamond$ Heuristics often "embarrassingly/naturally parallel"; PSO = particle swarm method

- Typically through stochastic sampling/evolution
- 1024 function evaluations per iteration
$\diamond$ Simplex is Nelder-Mead; POUNDERS is model-based trust-region algorithm
* one function evaluation per iteration


## Global Optimization, $\min _{x \in \Omega} f(x)$

Careful:
$\diamond$ Global convergence: Convergence (to a local solution/stationary point) from anywhere in $\Omega$
$\diamond$ Convergence to a global minimizer: Obtain $x^{*}$ with $f\left(x^{*}\right) \leq f(x) \forall x \in \Omega$

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## Anyone selling you global solutions when derivatives are unavailable:

either assumes more about your problem (e.g., convex $f$ )
or expects you to wait forever
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Instead:
$\diamond$ Rapidly find good local solutions and/or be robust to poor solutions
$\diamond$ Consider multistart approaches and/or structure of multimodality

## Why Multistart?

Best minimizer(s) approximate global minimizer $x^{*}, f\left(x^{*}\right) \leq f(x) \forall x \in \mathcal{D}$

## Multiple local minima are often of interest in practice

Design Multiple objectives/constraints might later be of interest
Distinctness $j$ best minimizers have physical meaning
Simulation Errors Spurious local minima from simulator anomalies

Uncertainty Some minima more sensitive to perturbations


## Increased opportunity for parallelism

Trilevel simulation/function $\rightarrow$ local solver $\rightarrow$ global solver

## Efficient local solvers

$\diamond$ (Local) surrogate-based, exploit problem structure

- least-squares objectives, (un)relaxable constraints, known nonsmoothness,...


## Motivating Example: Staging a Laser Plasma Accelerator


$\diamond$ Electron bunch is injected in a laser-induced plasma wave

- Typically when laser intensity reaches its first maximum
$\diamond$ Nonlinear effects $\Rightarrow$ plasma wave shrinks and electron bunch is lost
- Typically because bunch ends up in a defocusing region when laser intensity reaches its (first) minimum

Goal: Shape initial section of capillary to raise the minimum intensity and/or lower the maximum intensity.
$\rightarrow$ For a given $x$, we compute $v(t ; x)$, the (smooth) laser intensity at time $t$
Under ComPASS-3 with Carlo Benedetti \& Jean-Luc Vay (LBNL)

Motivating Example: $\min \left\{f(x): x \in \mathcal{D} \subset \mathbb{R}^{n}\right\}$

Simulation provides intensity at a discrete set of times
$t_{1}<\cdots<t_{p=|I|}:$

$$
B_{i}(x)=v\left(t_{i} ; x\right), \quad i \in I
$$

$f(x)=\max _{i \in \Theta_{1}(x)} v\left(t_{i} ; x\right)-\min _{i \in I} v\left(t_{i} ; x\right)$

$\Theta_{1}(x)=\left\{i \in I: i \leq \max \underset{j \in I}{\operatorname{argmin}} v\left(t_{j} ; x\right)\right\}$

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## Slice Through LPA Subproblem



This is a nonsmooth (piecewisesmooth) function of the parameters $x$

## LPA Feasible Region

| Variable | Range |
| :---: | :---: |
| Length | $2 \leq L \leq 6$ |
| Plasma channel radius | $1 \leq X_{\max } \leq 1.5$ |
| Minimum channel radius | $0.7 \leq X_{\min } \leq 1$ |
| Longitudinal location | $0 \leq Z_{\min } \leq 1$ |
| Laser focus position | $-1.2 \leq Z_{f} \leq 0$ |
| $c_{1}(x)=-X_{\max } Z_{\min }^{4}$ |  |
| $\quad-\left(X_{\max }-X_{\min }\right)\left(2 Z_{\min }-3 Z_{\min }^{2}\right)$ |  |
| $\leq 0$ |  |
| $c_{2}(x)=$ | $X_{\max }\left(Z_{\min }^{4}-4 Z_{\min }^{3}+3 Z_{\min }^{2}\right)$ |
|  | $+\left(X_{\max }-X_{\min }\right)\left(3 Z_{\min }^{2}-4 Z_{\min }+1\right)$ |
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$c(x) \leq 0$ are UNRELAXABLE: Simulator (often) fails in $\mathcal{D}^{c}$
QUAK
SBO constraint taxonomy $\rightarrow$ [Le Digabel \& W.; ANL/MCS-P5350-0515]

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## Numerical Experiments on LPA Problem

## Test multimodality:

$\diamond 51$ starting points $x^{0}$ generated uniformly from within $\mathcal{D}$
$\diamond$ Significant variation in $f\left(x^{0}\right)$
$\diamond$ Includes pathological $t_{1}=$
$\arg \max _{i \in \Theta_{1}\left(x^{0}\right)} v\left(t_{i}\right.$
$\diamond$ Maximum of $20 n v$ evaluations
(7.5 minutes each)
$\diamond 51$ CPU days


Solutions Found for I D^ Drahinm

## 51 Solutions:

$\diamond$ Converge to two solutions (A, B)
$\diamond \approx 10 \%$ to B
$\diamond$ Behavior after
$t_{\max \left\{i: i \in \Omega_{1}\right\}}$ unconstrained
$\diamond c\left(x^{A}\right), c\left(x^{B}\right)<0$
PS solutions remarkably consistent


Structured POUNDER code

## Solutions Found for LPA Problem

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Constrained Nelder-Mead code

## POPAS Activity Proposed for ComPASS-4

## Platform for Optimization of Particle Accelerators at Scale

$\diamond$ integrated platform for coordinating the evaluation and numerical optimization of accelerator simulations on leadership-class DOE computers
$\diamond$ orchestrate concurrent evaluations of OSIRIS, QuickPIC, Synergia, and MARS (or combinations thereof) with distinct inputs/parameter values
$\diamond$ account for resource requirements of the above
$\diamond$ API will allow the user to describe the mapping from simulation outputs and the derived quantities of interest used to define objective and constraint quantities

TH: Provide enough information so that optimization is efficient
"Simplest" (=Most Naive) Formulation: Blackbox $f$

Inputs

## Optimizer gives $x$, physicist provides $f(x)$

$\diamond f$ can be a blackbox (executable only or proprietary/legacy codes)
$\diamond$ Only give a single output

- no derivatives with respect to $x: \nabla_{x} S(x), \nabla_{x, x}^{2} S(x)$
- no problem structure


## Good solutions guaranteed in the limit, but:

$\diamond$ Computational budget limits number of evaluations
-SIMULATOR-
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## Two main styles of local algorithms

$\diamond$ Direct search methods (pattern search, Nelder-Mead, ...)
$\diamond$ Model- ("surrogate-" )based methods (quadratics, radial basis functions, ...)

## Black-Box Algorithms: Direct Search Methods

Pattern Search + Variants


Easy to parallelize $f$ evaluations

## Nelder-Mead + Variants

Popularized by Numerical Recipes
$\diamond$ Rely on indicator functions: $\left[f\left(x_{k}+s\right)<? f\left(x_{k}\right)\right] f\left(x_{k}\right)$, short memory
$\diamond$ Work with black-box $f(x)$, do not exploit structure $F[x, S(x)]$
$\diamond$ Convergence results for variety of settings

Survey $\rightarrow$ [Kolda, Lewis, Torczon; SIREV 2003]
Newer NM $\rightarrow$ [Lagarias, Poonen, Wright; SIOPT 2012] Tools $\rightarrow$ DFL [Liuzzi et al.], NOMAD [Audet et al.], . . .

## Making the Most of Little Information About Smooth $f$

$\diamond$ Overhead of the optimization routine is minimal (negligible?) relative to cost of evaluating simulation


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## Bank of data, $\left\{x_{i}, f\left(x_{i}\right)\right\}_{i=1}^{k}$

$=$ Points (\& function values) evaluated so far
$=$ Everything known about $f$
Goal:
$\diamond$ Make use of growing Bank as optimization progresses
$\diamond$ Limit unnecessary evaluations
(geometry/approximation)

## Derivative-Free, Model-Based Trust-Region Algorithms

$f$ expensive, no $\nabla f$
Substitute $\min \left\{m_{k}(x): x \in \mathcal{B}_{k}\right\}($ TRSP $)$ for $\min f(x)$ $m_{k}$ cheap, analytic derivatives

## Trust region:

$$
\mathcal{B}_{k}=\left\{x \in \Omega:\left\|x-x^{k}\right\| \leq \Delta_{k}\right\}
$$

## Basic algorithm

$\diamond$ Build model $m_{k}\left(\approx f\right.$ in $\left.\mathcal{B}_{k}\right)$
$\diamond x^{+} \approx \arg \min \left\{m_{k}(x): x \in \mathcal{B}_{k}\right\}$
$\diamond \rho_{k}=\frac{f\left(x^{k}\right)-f\left(x^{+}\right)}{m_{k}\left(x^{k}\right)-m_{k}\left(x^{+}\right)}$
$\diamond$ If $\rho_{k} \geq \eta_{1}>0$, accept $x^{k+1}=x^{+}$; Elseif $m_{k}$ is valid in $\mathcal{B}_{k}$, shrink $\Delta_{k}$ Else, improve $m_{k}$ in $\mathcal{B}_{k}$

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ORBIT: [W., Regis, Shoemaker, SISC 2008]

## Radial Basis Function Interpolation Models

Given
$\diamond$ base point $x_{k}$
$\diamond$ interpolation points
$\mathcal{Y}=\left\{y_{j}\right\}_{j=1}^{|\mathcal{Y}|} \subset \mathbb{R}^{n}$
$\diamond$ values $f\left(x_{k}+y_{j}\right)$ for $j=1, \ldots,|\mathcal{Y}|$
$\diamond$ radial kernel $\phi: \mathbb{R}_{+} \rightarrow \mathbb{R}$


Unique coefficients $\lambda$ and polynomial $p$ define interpolating RBF model

$$
m_{k}^{f}\left(x_{k}+s\right)=\sum_{j=1}^{|\mathcal{Y}|} \lambda_{j} \phi\left(\left\|s-y_{j}\right\|\right)+p(s)
$$

## Structure in Simulation-Based Optimization, $\min f(x)=F[x, S(x)]$

$f$ is often not a black box $S$
NLS Nonlinear least squares

$$
f(x)=\sum_{i}\left(S_{i}(x)-d_{i}\right)^{2}
$$

CNO Composite (nonsmooth) optimization

$$
f(x)=h(S(x))
$$

SKP Not all variables enter simulation

$$
f(x)=g\left(x_{I}, x_{J}\right)+h\left(S\left(x_{J}\right)\right)
$$

BLO Bilevel optimization

$$
\min \left\{S_{1}\left(x_{I}, x_{J}\right): x_{I} \in \arg \max _{y} S_{2}\left(y, x_{J}\right)\right\}
$$

SCO Only some constraints depend on simulation

$$
\min \left\{f(x): c_{1}(x)=0, c_{S}(x)=0\right\}
$$

Model-based methods offer one way to exploit such structure

## Nonlinear Least Squares $f(x)=\frac{1}{2} \sum_{i} R_{i}(x)^{2}$

## Obtain a vector of output $R_{1}(x), \ldots, R_{p}(x)$

$\diamond$ Model each $R_{i}$

$$
R_{i}(x) \approx m_{k}^{R_{i}}(x)=R_{i}\left(x_{k}\right)+\left(x-x_{k}\right)^{\top} g_{k}^{(i)}+\frac{1}{2}\left(x-x_{k}\right)^{\top} H_{k}^{(i)}\left(x-x_{k}\right)
$$

$\diamond$ Approximate:

$$
\begin{aligned}
\nabla f(x)= & \sum_{i} \nabla \mathbf{R}_{\mathbf{i}}(\mathbf{x}) R_{i}(x) \quad \longrightarrow \sum_{i} \nabla m_{k}^{R_{i}}(x) R_{i}(x) \\
\nabla^{2} f(x)= & \sum_{i} \nabla \mathbf{R}_{\mathbf{i}}(\mathbf{x}) \nabla \mathbf{R}_{\mathbf{i}}(\mathbf{x})^{\top}+\sum_{i} R_{i}(x) \nabla^{\mathbf{2}} \mathbf{R}_{\mathbf{i}}(\mathbf{x}) \\
& \longrightarrow \sum_{i} \nabla m_{k}^{R_{i}}(x) \nabla m_{k}^{R_{i}}(x)^{\top}+\sum_{i} R_{i}(x) \nabla^{2} m_{k}^{R_{i}}(x)
\end{aligned}
$$

$\diamond$ Model $f$ via Gauss-Newton or similar

$$
\text { regularized Hessians } \rightarrow \text { DFLS [Zhang, Conn, Scheinberg] }
$$ full Newton $\rightarrow$ POUNDERS [W., Moré]

## POUNDERS for $\chi^{2}$ (=Nonlinear Least Squares Calibration)

POUNDERS (in PETSc/TAO) well tested for calibration problems:

$$
f(x) \propto \sum_{i, j} W_{i, j}\left(S\left(x ; \theta^{i}\right)-d_{i}\right)\left(S\left(x ; \theta^{j}\right)-d_{j}\right)
$$



## Constraints in Simulation-Based Optimization


[le Digabel, W.; 2017]; [Regis, W.; OMS, 2017]

## Why Expressing Constraint Functions Matters

Augmented Lagrangian methods, $L_{A}(x, \lambda ; \mu)=f(x)-\lambda^{T} c(x)+\frac{1}{\mu}\|c(x)\|^{2}$

## $\min _{x}\{f(x): c(x)=0\}$

Four choices:

1. Penalize constraints
2. Treat $c$ and $f$ both as (separate) black boxes
3. Work with $f$ and $\nabla_{x} c$
4. Have both $\nabla_{x} f$ and $\nabla_{x} c$
$\rightarrow$ With Slava Kungurtsev


$$
n=15,11 \text { constraints }
$$

## What is APOSMM?

## Asynchronous Parallel Optimization Solver for Multiple Minima

$\diamond$ Better account for dynamic number of local runs
$\diamond$ Decouple local run from fixed resource
$\diamond$ Anticipate nontrivial $\operatorname{Var}[\operatorname{time}(f(x))]$
[Larson \& W. Asynchronously Parallel Optimization Solver for Finding Multiple Minima, Math. Program. Comput., 2018.]


## The (A)POSMM Algorithm

Repeat:
$\diamond$ Receive from worker(s) $w_{\ell} \in W$ that has evaluated its point
$\diamond$ If point was a sample point, update $r_{k}=\frac{1}{\sqrt{\pi}} \sqrt[n]{\operatorname{vol}(\mathcal{D}) \frac{5 \Gamma\left(1+\frac{n}{2}\right) \log \left(\left|\mathcal{S}_{k}\right|\right)}{\left|\mathcal{S}_{k}\right|}}$
$\diamond$ If point was a local optimization point, add subsequent point in the run (not in $\mathcal{H}_{k}$ ) to $Q_{L}$ if not terminated
$\diamond$ Start run(s) at all point(s) now satisfying conditions, adding subsequent point from each run to $Q_{L}$
$\diamond$ Merge/collapse runs within $Q_{L}$
$\diamond$ Send point(s) from $Q_{L}$ and/or $\mathcal{R}$ to worker(s)
$W$ Set of workers
$\mathcal{R}$ Stream of sample points (from $\mathcal{D}$ )
$\mathcal{S}_{k}$ Sample points after iteration $k$
$Q_{L}$ Queue of local optimization points (needed by $\mathcal{A}$ )
$\mathcal{H}_{k}$ History after $k$ evaluations

## Basic Idea: Multi Level Single Linkage (MLSL) Clustering

Where to start $\mathcal{A}$ in $k$ th iteration [Rinnooy Kan \& Timmer (MathProg, 1987)]


## Start $\mathcal{A}$ at each sample point $x^{i} \in \mathcal{S}_{k}$ provided:

$\diamond \mathcal{A}$ has not been started from $x^{i}$, and
$\diamond$ no other sample point $x^{j} \in \mathcal{S}_{k}$ with $f\left(x^{j}\right)<f\left(x^{i}\right)$ is within a distance

$$
r_{k}=\frac{1}{\sqrt{\pi}} \sqrt[n]{\operatorname{vol}(\mathcal{D}) \frac{5 \Gamma\left(1+\frac{n}{2}\right) \log (k N)}{k N}}
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Ex.: It. 1 Exploration

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Ex.: It. 1 Exploration

Thm [RK-T]- With probability 1, MLSL will start finitely many local runs.

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Ex.: It. 1 Refinement

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Ex.: It. 2 Exploration

Thm [RK-T]- With probability 1, MLSL will start finitely many local runs.

## (A)POSMM Framework



## Data Profiles: Ability to Find Approximate Global Minimizer

## 600 GKLS problems

(A) POSMM
$\diamond$ Makes rapid progress to $f_{G}$
$\diamond$ Outperforms other algorithms (even while demanding 14-fold concurrency) evaluations


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## Argonne/Optimization Milestones in ComPASS-4

| Activity | Institution(s) | Sec | Year |
| :--- | :--- | :--- | :--- |
| Develop API for POPAS prototype | ANL, FNAL, <br> UCLA | $\S 2.4$ | 1 |
| Identify optimizable elements in the MARS <br> and Synergia PIP-II models; connect with <br> POPAS prototype | FNAL, ANL | $\S 2.1 .1$ | 2 |
| Use MARS-Synergia-POPAS prototype for <br> preliminary optimization | FNAL, ANL | $\S 2.1 .1$ | 3 |
| Include prototype of structure-exploiting <br> optimization algorithm for standard <br> PIC/QuickPIC simulations; enable basic <br> execution of all ComPASS-4 codes in POPAS | ANL, FNAL, <br> UCLA | $\S 2.4$ | 3 |
| Link numerical optimization algorithm to <br> POPAS; Remove file I/O layer from POPAS | ANL, FNAL, <br> UCLA | $\S 2.4$ | 3 |
| Connect IOTA Synergia model with POPAS | FNAL, ANL | $\S 2.1 .1$ | 3 |
| Release POPAS; apply POPAS to standard <br> PIC/QuickPIC and Synergia | ANL, FNAL, <br> UCLA | $\S 2.4$ | 4 |
| Refine MARS-Synergia-POPAS | FNAL, ANL | $\S 2.1 .1$ | 4 |
| Apply IOTA Synergia-POPAS | FNAL, ANL | $\S 2.1 .1$ | 4 |
| Carry out parameter optimization on <br> PWFA-LC relevant parameters using <br> QuickPIC | UCLA, FNAL, <br> ANL | $\S 2.5 .2$ | 5 |


[^0]:    ORBIT: [W., Regis, Shoemaker, SISC 2008]

[^1]:    ORBIT: [W., Regis, Shoemaker, SISC 2008]

[^2]:    ORBIT: [W., Regis, Shoemaker, SISC 2008]

