# Pairwise Sequence Alignment 

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## Protein Evolution

"For many protein sequences, evolutionary history can be traced back 1-2 billion years"
-William Pearson

- When we align sequences, we assume that they share a common ancestor
- They are then homologous
- Protein fold is much more conserved than protein sequence
$\checkmark$ DNA sequences tend to be less informative than protein sequences


## Definition

- Homology: related by descent
- Homologous sequence positions

| ATTGCGC | $\rightarrow$ ATTGCGC |
| :--- | :--- |
| ATICCGC | $\rightarrow$ ATCCGC |

## Orthologous and paralogous

- Orthologous sequences differ because they are found in different species (a speciation event)
- Paralogous sequences differ due to a gene duplication event
- Sequences may be both orthologous and paralogous


## Pairwise Alignment

- The alignment of two sequences (DNA or protein) is a relatively straightforward computational problem.
- There are lots of possible alignments.
- Two sequenċes can always be aligned.
- Sequence alignments have to be scored.
- Often there is more than one solution with the same score.


## Methods of Alignment

- By hand - slide sequences on two lines of a word processor
- Dot plot
- with windows
$\bullet$ Rigorous mathematical approach
- Dynamic programming (slow, optimal)
$\bullet$ Heuristic methods (fast, approximate)
- BLAST and FASTA
- Word matching and hash tables0


## Align by Hand

## GATCGCCTA_TTACGTCCTGGAC <----> AGGCATACGTA_GCCCTTTCGC

You still need some kind of scoring system to find the best alignment

## Percent Sequence Identity

The extent to which two nucleotide or amino acid sequences are invariant


## Dotplot:

A dotplot gives an overview of all possible alignments


## Dotplot:

In a dotplot each diagonal corresponds to a possible (ungapped) alignment


One possible alignment:

## Insertions / Deletions in a Dotplot


$\begin{array}{llllllllll}I & I & I & I & I & & I & I & 1 & I \\ T & A & C & T & G & T & T & C & A & T\end{array}$


Hemoglobin $\alpha$-chain

## Word Size Algorithm

$$
\begin{aligned}
& T \begin{array}{|lll|lllll}
\hline \mathbf{A} & \mathbf{C} & \mathbf{G} & \mathrm{G} & \mathrm{~T} & \mathbf{A} & \mathrm{~T} & \mathrm{G} \\
\mathbf{1} & 1 & & 1 & 1 & 1 & 1 & \\
\mathbf{A} & \mathbf{C} & \mathbf{A} & \mathrm{G} & \mathrm{~T} & \mathbf{A} & \mathrm{~T} & \mathrm{C}
\end{array} \\
& \begin{array}{c|ccc|cccc}
\mathrm{T} & \mathrm{~A} & \mathbf{C} & \mathbf{G} & \mathbf{G} & \mathrm{~T} & \mathrm{~A} & \mathrm{~T} \\
\mathrm{I} & \mathrm{G} \\
\mathrm{I} & \mathrm{C} & & 1 & \mathrm{I} & 1 & 1 & \\
\text { A } & \mathbf{C} & \mathbf{A} & \mathbf{G} & \mathrm{T} & \mathrm{~A} & \mathrm{~T} & \mathrm{C}
\end{array} \\
& \begin{array}{lll|lll|lll}
\text { T } & \text { A } & C & \mathbf{G} & \mathbf{G} & \mathbf{T} & A & T & G \\
1 & 1 & & 1 & 1 & 1 & 1 & \\
& A & C & A & \mathbf{G} & \mathbf{T} & A & \mathrm{~A} & \mathrm{C}
\end{array} \\
& \begin{array}{llll|lll|ll}
\mathrm{T} & \mathrm{~A} & \mathrm{C} & \mathrm{G} & \mathrm{G} & \mathbf{T} & \mathbf{A} & \mathrm{~T} & \mathrm{G} \\
\mathrm{I} & 1 & & 1 & 1 & 1 & 1 & \\
\mathrm{~A} & \mathrm{C} & \mathrm{~A} & \mathrm{G} & \mathbf{T} & \mathbf{A} & \mathrm{~T} & \mathrm{C}
\end{array}
\end{aligned}
$$

Word Size $=3$


## Window / Stringency



Scoring Matrix Filtering

Matrix: PAM250<br>Window = 12<br>Stringency $=9$



Hemoglobin $\alpha$-chain

## Considerations

- The window/stringency method is more sensitive than the wordsize method (ambiguities are permitted).
- The smaller the window, the larger the weight of statistical (unspecific) matches.
- With large windows the sensitivity for short sequences is reduced.
- Insertions/deletions are not treated explicitly.


## Alignment methods

- Rigorous algorithms = Dynamic Programming
- Needleman-Wunsch (global)
- Smith-Waterman (local)
-Heuristic algorithms (faster but approximate)
- BLAST
- FASTA


# Basic principles of dynamic programming 

- Creation of an alignment path matrix
- Stepwise calculation of score values
- Backtracking (evaluation of the optimal path)


## Dynamic Programming

- Dynamic Programming is a very general programming technique.
- It is applicable when a large search space can be structured into a succession of stages, such that:
- the initial stage contains trivial solutions to subproblems
- each partial solution in a later stage can be calculated by recurring a fixed number of partial solutions in an earlier stage
- the final stage contains the overall solution


## Creation of an alignment path matrix

## Idea: <br> Build up an optimal alignment using previous solutions for optimal alignments of smaller subsequences

- Construct matrix F indexed by $i$ and $j$ (one index for each sequence)
- $F(i, j)$ is the score of the best alignment between the initial segment $x_{1 . . . i}$ of $x$ up to $x_{i}$ and the initial segment $y_{1 . . . j}$ of $y$ up to $y_{j}$
- Build $F(i, j)$ recursively beginning with $F(0,0)=0$
$s: A G C A C A C-A \quad A G-C A C A C A$ $t: A-C A C A C T A$ or $A C A C A C T-A$



## Creation of an alignment path matrix

- If $F(i-1, j-1), F(i-1, j)$ and $F(i, j-1)$ are known we can calculate $F(i, j)$
- Three possibilities:
- $x_{i}$ and $y_{j}$ are aligned, $F(i, j)=F(i-1, j-1)+s\left(x_{i}, y_{j}\right)$
- $x_{i}$ is aligned to a gap, $F(i, j)=F(i-1, j)-d$
- $y_{j}$ is aligned to a gap, $F(i, j)=F(i, j-1)-d$
- The best score up to $(i, j)$ will be the largest of the three options


## Backtracking



Optimal global alignment: HEAGAWGHE-E - - P-AW-HEAE

## Global vs. Local Alignments

- Global alignment algorithms start at the beginning of two sequences and add gaps to each until the end of one is reached.
- Local alignment algorithms finds the region (or regions) of highest similarity between two sequences and build the alignment outward from there.


## Global Alignment



Local Alignment




## Global Alignment

Two closely related sequences:

1 TGTCGATTAAGCGGTCGTAGCTGACCTGAGATTGCCCGATGGCGTAGTAGCTGACC 56

1 TGTCGATTATGCGGTCGTAG . . GACCTGAGTTTCCCCGATGGCGTAGTAGGTGACC 54
needle (Needleman \& Wunsch) creates an end-to-end alignment.

## Global Alignment

## Two sequences sharing several regions of local similarity:

1 AGGATTGGAㅗㅗTGCTCAGAㅗㅗGCAGCTAAㅗㅅㅗGCGTGTATGCAGGATTGGAㅗㅗTTAㅗㅗㅅㅗGAGGAGGTAGACCG ..... 67


# Global Alignment (Needleman -Wunsch) 

- The the Needleman-Wunsch algorithm creates a global alignment over the length of both sequences (needle)
- Global algorithms are often not effective for highly diverged sequences - do not reflect the biological reality that two sequences may only share limited regions of conserved sequence.
- Sometimes two sequences may be derived from ancient recombination events where only a single functional domain is shared.
- Global methods are useful when you want to force two sequences to align over their entire length


## Local Alignment (Smith-Waterman)

- Local alignment
- Identify the most similar sub-region shared between two sequences
- Smith-Waterman
- EMBOSS: water


## Parameters of Sequence Alignment

## Scoring Systems:

- Each symbol pairing is assigned a numerical value, based on a symbol comparison table.


## Gap Penalties:

- Opening: The cost to introduce a gap
- Extension: The cost to elongate a gap


## DNA Scoring Systems -very simple

## Sequence 1 Sequence 2

 actaccagttcatttgatacttctcaaataccattaccgtgttaactgaaaggacttaaagact

|  | $\mathbf{A}$ | $\mathbf{G}$ | $\mathbf{C}$ | $\mathbf{T}$ |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| A | 1 | 0 | 0 | 0 |  |
| G | 0 | 1 | 0 | 0 | Match: 1 |
| C | 0 | 0 | 1 | 0 | Mismatch: 0 |
| T | 0 | 0 | 0 | 1 | Score = 5 |
|  |  |  |  |  |  |

## Protein Scoring Systems



## Protein Scoring Systems

- Amino acids have different biochemical and physical properties that influence their relative replaceability in evolution.



## Protein Scoring Systems

- Scoring matrices reflect:
- \# of mutations to convert one to another
- chemical similarity
- observed mutation frequencies
- the probability of occurrence of each amino acid
- Widely used scoring matrices:
- PAM
- BLOSUM


## PAM matrices

> Family of matrices PAM 80, PAM 120, PAM 250
> The number with a PAM matrix represents the evolutionary distance between the sequences on which the matrix is based
> Greater numbers denote greater distances

## PAM (Percent Accepted Mutations) matrices

- The numbers of replacements were used to compute a so-called PAM-1 matrix.
- The PAM-1 matrix reflects an average change of $1 \%$ of all amino acid positions. PAM matrices for larger evolutionary distances can be extrapolated from the PAM-1 matrix.
- PAM250 $=250$ mutations per 100 residues.
- Greater numbers mean bigger evolutionary distance


## PAM (Percent Accepted Mutations) matrices

- Derived from global alignments of protein families . Family members share at least 85\% identity (Dayhoff et al., 1978).

- Construction of phylogenetic tree and ancestral sequences of each protein family
- Computation of number of replacements for each pair of amino acids


## PAM 250



## PAM - limitations

## - Based on only one original dataset

> Examines proteins with few differences ( $85 \%$ identity)

- Based mainly on small globular proteins so the matrix is biased


## BLOSUM matrices

> Different BLOSUM $\_$matrices are calculated independently from BLOCKS (ungapped local alignments)
> BLOSUMr is based on a cluster of BLOCKS of sequences that share at least $/$ percent identity
> BLOSUM6゙2 represents closer sequences than BLOSUM45

## BLOSUM (Blocks Substitution Matrix)

- Derived from alignments of domains of distantly related proteins (Henikoff \& Henikoff, 1992).

- Occurrences of each amino acid pair in each column of each block alignment is counted.
- The numbers derived from all blocks were used to compute the BLOSUM matrices.

$$
\begin{aligned}
& A-C=4 \\
& A-E=2 \\
& C-E=2 \\
& A-A=1 \\
& C-C=1
\end{aligned}
$$

## The Blosum50 Scoring Matrix

|  | A | R | N | D | C | Q | E | G | H | I | L | K | M | I |  | P | S | T | W | W | Y | V | B | Z | X |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | 5 | -2 | -1 | -2 | -1 | -1 | -1 | 0 | -2 | -1 | -2 | -1 | -1 | 1 -3 | 3 | -1 | 1 | 0 | 0 - | -3 | -2 | 0 | -2 | -1 | -1 | -5 |
| R | -2 | 7 | -1 | -2 | -4 | 1 | 0 | -3 | 0 | -4 | -3 | 3 | -2 | 2 | 3 | -3 | -1 | -1 | 1. | -3 | -1 | - -3 | -1 | 0 | -1 | -5 |
| N | -1 | -1 | 7 | 2 | -2 | 0 | 0 | 0 | 1 | -3 | -4 | 0 | -2 | 2 - | 4 | -2 | 1 | 0 | 0 - | -4 | -2 | -3 | 4 | 0 | -1 | -5 |
| D | -2 | -2 | 2 | 8 | -4 | 0 | 2 | -1 | -1 | -4 | -4 | -1 | -4 |  | 5 | -1 | 0 | -1 | 1 | -5 | -3 | -4 | 5 | 1 | -1 | -5 |
| C | -1 | -4 | -2 | -4 | 13 | -3 | -3 | -3 | -3 | -2 | -2 | -3 | -2 | -2 | 2 | -4 | -1 | -1 | - | - | -3 | -1 | -3 | -3 | 2 | -5 |
| Q | -1 | 1 | 0 | 0 | -3 | 7 | 2 | -2 | 1 | -3 | -2 | 2 | 0 | - | 4 | -1 | 0 | -1 | 1 - | -1 | -1 | -3 | 0 | 4 | -1 | -5 |
| E | -1 | 0 | 0 | 2 | -3 | 2 | 6 | -3 | 0 | -4 | -3 | 1 | -2 | -2 | 3 | -1 | -1 | -1 | - | -3 | -2 | -3 | 1 | 5 | 1 | -5 |
| G | 0 | -3 | 0 | -1 | -3 | -2 | -3 | 8 | -2 | -4 | -4 | -2 | -3 | - | 4 | -2 | 0 | -2 |  | -3 | -3 | -4 | 1 | -2 | 2 | 5 |
| H | -2 | 0 | 1 | -1 | -3 | 1 | 0 | -2 | 10 | -4 | -3 | 0 | -1 | - | 1 | -2 | -1 | -2 | 2 | -3 | 2 | -4 | 0 | 0 | -1 | -5 |
| I | -1 | -4 | -3 | -4 | -2 | -3 | -4 | -4 | -4 | 5 | 2 | -3 | 2 | 2 | 0 | -3 | -3 | -1 |  | -3 | -1 | 4 | -4 | -3 | -1 | -5 |
| L | -2 | -3 | -4 | -4 | -2 | -2 | -3 | -4 | -3 | 2 | 5 | -3 | 3 | 3 | 1 | -4 | -3 | -1 | 1 | -2 | -1 | 1 | -4 | -3 | -1 | 5 |
| K | -1 | 3 | 0 | -1 | -3 | 2 | 1 | -2 | 0 | -3 | -3 | 6 | -2 | 2 | 4 | -1 | 0 | -1 |  | -3 | -2 | -3 | 0 | 1 | -1 | 5 |
| M | -1 | -2 | -2 | -4 | -2 | 0 | -2 | -3 | -1 | 2 | 3 | -2 | 7 | 70 | 0 | -3 | -2 | -1 |  | -1 | 0 | 1 | -3 | -1 | -1 | -5 |
| F | -3 | -3 | -4 | -5 | -2 | -4 | -3 | -4 | -1 | 0 | 1 | -4 | 0 | 0 | 8 | -4 | -3 | -2 |  | 1 | 4 | -1 | -4 | -4 | -2 | -5 |
| P | -1 | -3 | -2 | -1 | -4 | -1 | -1 | -2 | -2 | -3 | -4 | -1 | -3 | -4 | 4 | 10 | -1 | -1 | - | - | -3 | -3 | -2 | -1 | -2 | -5 |
| S | 1 | -1 | 1 | 0 | -1 | 0 | -1 | 0 | -1 | -3 | -3 | 0 | -2 | 2 |  | -1 | 5 | 2 | 2 | -4 | -2 | -2 | 0 | 0 | -1 | -5 |
| T | 0 | -1 | 0 | -1 | -1 | -1 | -1 | -2 | -2 | -1 | -1 | -1 | -1 | -2 | 2 | -1 | 2 | 5 | - | -3 | -2 | 0 | 0 | -1 | 0 | -5 |
| W | -3 | -3 | -4 | -5 | -5 | -1 | -3 | -3 | -3 | -3 | -2 | -3 | -1 | 1 | 1 | -4 | -4 | -3 |  | 15 | 2 | -3 | -5 | -2 | -3 | -5 |
| Y | -2 | -1 | -2 | -3 | -3 | -1 | -2 | -3 | 2 | -1 | -1 | -2 | 0 | 0 | 4 | -3 | -2 | -2 | 2 | 2 | 8 | -1 | -3 | -2 | -1 | -5 |
| V | 0 | -3 | -3 | -4 | -1 | -3 | -3 | -4 | -4 | 4 | 1 | -3 |  | 1 - | 1 | -3 | -2 | 0 | 0 | -3 | -1 | 5 | -4 | -3 | -1 | -5 |
| B | -2 | -1 | 4 | 5 | -3 | 0 | 1 | -1 | 0 | -4 | -4 | 0 | -3 | 3 - | 4 | -2 | 0 | 0 | $0-$ | -5 | -3 | -4 | 5 | 2 | -1 | -5 |
| Z | -1 | 0 | 0 | 1 | -3 | 4 | 5 | -2 | 0 | -3 | -3 | 1 | -1 | 1 - | 4 | -1 | 0 | -1 | 1 | -2 | -2 | -3 | 2 | 5 | -1 | -5 |
| X | -1 | -1 | -1 | -1 | -2 | -1 | -1 | -2 | -1 | -1 | -1 | -1 | -1 | 1 -2 | 2 | -2 | -1 | 0 | 0 | -3 | -1 | -1 | -1 | -1 | -1 | -5 |
| * | -5 | -5 | -5 | -5 | -5 | -5 | -5 | -5 | -5 | -5 | -5 | -5 | -5 | 5 | 5 | -5 | -5 | -5 | 5 | -5 | -5 | -5 | -5 | -5 | -5 | 1 |

## BLOSUM (Blocks Substitution Matrix)

- Sequences within blocks are clustered according to their level of identity.
- Clusters are counted as a single sequence.
- Different BLOSUM matrices differ in the percentage of sequence identity used in clustering.
- The number in the matrix name (e.g. 62 in BLOSUM62) refers to the percentage of sequence identity used to build the matrix.
- Greater numbers mean smaller evolutionary distance.


# PAM Vs. BLOSUM PAM100 = BLOSUM90 PAM120 = BLOSUM80 PAM160 = BLOSUM60 PAM200 = BLOSUM52 PAM250 = BLOSUM45 

## More distant sequences

-PAM120 for general use -PAM60 for close relations PAM250 for distant relations
-BLOSUM62 for general use -BLOSUM80 for close relations -BLOSUM45 for distant relations

## TIPS on choosing a scoring matrix

- Generally, BLOSUM matrices perform better than PAM matrices for local similarity searches (Henikoff \& Henikoff, 1993).
- When comparing
proteins one should use matrices, for distantly related proteins higher PAM or lower BLOSUM matrices.
- For database searching the commonly used matrix is BLOSUM62.


## Scoring Insertions and Deletions



The creation of a gap is penalized with a negative score value.

## Why Gap Penalties?

Gaps not permitted
Score: 0
1 GTGATAAGACAACAAGAACCGGTGGCATTGTGG 29 ||| | | ||| | || || |
1 GTGTCGGG소소G소G소T소소CTCCGATGGTTG 29

Match = 5
Mismatch $=-4$

Gaps allowed but not penalized Score: 88

1 GTG. ATAG. ACACAGA . CCGGT . . GGCATTGTGG 29 ||| || | | ||| || | | || || |
1 GTGTAT.GGA. AGAAGATACC. .TCCG. . ATGGTTG 29

## Why Gap Penalties?

- The optimal alignment of two similar sequences is usually that which
- maximizes the number of matches and
- minimizes the number of gaps.
-There is a tradeoff between these two
- adding gaps reduces mismatches
- Permitting the insertion of arbitrarily many gaps can lead to high scoring alignments of non-homologous sequences.
- Penalizing gaps forces alignments to have relatively few gaps.


## Gap Penalties

$\cdot$ How to balance gaps with mismatches?

- Gaps must get a steep penalty, or else you' 11 end up with nonsense alignments.
-In real sequences, muti-base (or amino acid) gaps are quit common
-genetic insertion/deletion events
- "Affine" gap penalties give a big penalty for each new gap, but a much smaller "gap extension" penalty.


## Scoring Insertions and Deletions

$$
\begin{aligned}
& \text { match = } 1 \\
& \text { mismatch = } 0
\end{aligned}
$$

Total Score: 4


Total Score: $\quad 8-3.2=4.8$

Gap parameters:
$d=3 \quad$ (gap opening)
$e=0.1$ (gap extension)
$g=3 \quad$ (gap lenght)
$\gamma(g)=-3-(3-1) 0.1=-3.2$


## Modification of Gap Penalties

Score Matrix: BLOSUM62


