## Quantum Computing with Haskell <br> and FPGA simulation <br> shuchang.zhou@gmail.com Jan. 18, 2018

## Why quantum computing?

- Can crack elliptic curve cryptography ...
- And threaten your Bitcoin
- 0

○ 1
Classical Bit Qubit



## Why study quantum computing <br> ... even when you don't have a quantum computer

- Many fast classic algorithms can be traced to simulations of quantum algorithms.

| Discrete Fourier <br> Transform | Simulated Annealing | Probabilistic <br> checking | BPP |
| :--- | :--- | :--- | :--- |
| Quantum Fourier <br> Transform | Quantum Annealing | Deutsch's algorithm | BQP |



Why study quantum computing
... even when you don't have a quantum computer
"Quantum computing may be the key, to understanding Deep Learning."


## Pure states (eigenvectors)

## Quantum Mechanics

- Schrödinger's equation
$i \hbar \frac{\partial}{\partial t} \Psi(\mathbf{r}, t)=\left[\frac{-\hbar^{2}}{2 \mu} \nabla^{2}+V(\mathbf{r}, t)\right] \Psi(\mathbf{r}, t)$
- von Neumann's equation
- Lax-pair, isospectral

$$
i \hbar \frac{\partial \rho}{\partial t}=[H, \rho]
$$

 $n=1$

## Bit to Qubit ("Q-bit")

- Bit: +1 / -1
- Qubit: 2DOF
- Angles: theta / phi
- Two complex numbers + norm constraint

■ $|\psi\rangle=\alpha|0\rangle+\beta|1\rangle$, (superposition)

- $|\alpha|^{2}+|\beta|^{2}=1$.
- $|\alpha|^{2}$ and $|\beta|^{2}$ are probabilities
- Bra-ket notation
- <x| for Bra, or row-vector, or transposed vector.
- |x> for vector


## Classical Bit



Qubit

## Exponential number of bits for simulating Qubits

- N qubits need $2^{\mathrm{N}}$ classic bits
- $\quad(a|0>+b| 1>)(c|0>+d| 1>)$
- $=a c|00>+a d| 01>+b c|10>+b d| 11>$
- Entanglement: when (a, b; c, d) not rank 1

$$
\text { - } \quad a|00>+b| 01>+c|10>+d| 11>
$$

## Superposition and Measurement



## Quantum operations

- Reversible
- Apply a unitary transform: all kinds of gates
- Unconditional
- Conditional
- Irreversible (Quantum decoherence)
- "Create" a qubit
- Measurement

$$
H=\frac{1}{\sqrt{2}}\left[\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right] \text { Hadamard }
$$

$$
X=\left[\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right] \quad \text { Pauli-X }
$$

$$
Y=\left[\begin{array}{cc}
0 & -i \\
i & 0
\end{array}\right] \quad \text { Pauli-Y }
$$

$$
R_{\phi}=\left[\begin{array}{cc}
1 & 0 \\
0 & e^{i \phi}
\end{array}\right] \begin{aligned}
& \text { Phase Shift, pi / } 4 \text { for } \\
& \text { "pi-over-eight" gate }
\end{aligned}
$$

$$
\cdots \quad\left[\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right]
$$

## Modeling quantum operations with Haskell

- Reversible
- Apply a unitary transform
- Unconditional
- Conditional
- Irreversible
- "Create" a qubit
- Measurement
-- | The underlying data type of a U unitary operation
data U = UReturn -- A List like construct
| Rot Qbit Rotation U
| Swap Qbit Qbit U
| Cond Qbit (Bool -> U) U
| Ulet Bool (Qbit -> U) U

Has side-effects, how to model?

## Monad and the "multi-world"

- Monad is a representation for computation graph
- Construct first, run later
- Haskell put everything with side-effects in IO monad
- putStrLn :: String -> IO ()
- "write : World -> Filename -> String -> World"
- type IO a = World -> (a, World)

World
10 a

Run

New World
a

## Example: IO Monad

- (>>=) :: IO a -> (a -> IO b) $->$ IO b (action1 >>= action2) world0 = let ( $a$, world1) $=$ action1 world0
(b, world2) = action2 a world1 in (b, world2)
- "Bind" operation
- return :: a -> IO a
return a world0 $=(\mathrm{a}$, world0)

Input 1
A Monad contains type a


## Example: State Monad

- newtype State s a = State $\{$ runState $::$ s -> $(\mathrm{a}, \mathrm{s})\}$
- return $\mathrm{a}=$ State $\$ \backslash \mathrm{~s}->(\mathrm{a}, \mathrm{s})$
- (>>=) :: State s a ->(a -> State s b) ->State s b
- $m \gg=k=$ State $\$$ ls $->$ let $\left(a, s^{\prime}\right)=$ runState $m$ s in runState (k a) s'



## Example: "RNN monad"

- newtype Rnn s i a = Rnn \{runRnn :: (i, s) -> (a, s) \}
- return a = Rnn \$ <br>(i, s) -> (a, s)
- (>>=) :: Rnn sia -> (a -> Rnn sjb) -> Rnn s (i, j) b
- $m \gg=k=R n n \$ \backslash(i, j), s)$-> let $\left(a, s^{\prime}\right)=\operatorname{runRnn} m(i, s)$ in runRnn $(k a)\left(j, s^{\prime}\right)$



## QIO Haskell package

- QIO models the "irreversible" part: decoherence of the qubits
- Forming a monad
instance Monad QIO
mkQbit :: Bool $\rightarrow$ QIO Qbit
applyU :: U $\rightarrow$ QIO ()
measQbit :: Qbit $\rightarrow$ QIO Bool


## QIO Monad can be simulated or sampled

- "run" for sampling
- "sim" for distributional representation

Prob :: * $\rightarrow *$<br>instance Monad Prob<br>run :: QIO a $\rightarrow$ IO a<br>sim :: QIO a $\rightarrow$ Prob a<br>runC : : QIO a $\rightarrow$ a

## Creating qubits

```
-- | Initialise a qubit in the |0> state
q0 :: QIO Qbit
q0 = mkQ False
-- | Initialise a qubit in the |1> state
q1 :: QIO Qbit
q1 = mkQ True
```

-- | Initialise a qubit in the |+> state. This is done by applying a Hadamard gate to the |0> state.
qPlus :: QIO Qbit
qPlus = do qa <- q0
applyU (uhad qa)
return qa

-- | Initialise a qubit in the |-> state. This is done by applying a Hadamard gate to the |1> state.
qMinus :: QIO Qbit
qMinus = do qa <-q1
applyU (uhad qa)
return qa


## Measuring and "sharing"

-- | Create a random Boolean value, by measuring the state |+> randBit :: QIO Bool
randBit $=$ do qa <-qPlus
$x<-$ measQbit qa
return $x$
-- | This function can be used to "share" the state of one qubit, with another
-- newly initialised qubit. This is not the same as "cloning", as the two qubits
-- will be in an entangled state. "sharing" is achieved by simply initialising
-- a new qubit in state $|0\rangle$, and then applying a controlled-not to that qubit,
-- depending on the state of the given qubit.
share :: Qbit -> QIO Qbit
share qa $=$ do qb <-q0

$$
\begin{aligned}
& \text { applyU (cond qa (la -> if a then (unot qb) } \\
& \text { else (mempty)) ) } \\
& \text { return qb }
\end{aligned}
$$



## Deutsch-Jozsa's algorithm

- Given a balanced/constant boolean function (Bool^k -> Bool)
- Do a 2-classification

| const True | const False | $\backslash x->x$ | $\backslash x->$ not $x$ |
| :--- | :--- | :--- | :--- |
| 1 | 1 | 0 | 0 |

- Exact solution on a Quantum computer requires 1 evaluation
- Exact solution on a classic computer requires exponential many evaluations
- ... But if allowing bounded errors, require k answers to obtain $\epsilon \leq 1 / 2^{k-1}$


## Manual work out

$U_{f}$ maps $|x\rangle|y\rangle$ to $|x\rangle|y \oplus f(x)\rangle$

$$
\begin{aligned}
& \frac{1}{2}(|0\rangle(|f(0) \oplus 0\rangle-|f(0) \oplus 1\rangle)+|1\rangle(|f(1) \oplus 0\rangle-|f(1) \oplus 1\rangle)) \\
& =\frac{1}{2}\left((-1)^{f(0)}|0\rangle(|0\rangle-|1\rangle)+(-1)^{f(1)}|1\rangle(|0\rangle-|1\rangle)\right)
\end{aligned}
$$

$$
\frac{1}{2}(|0\rangle+|1\rangle)(|0\rangle-|1\rangle) . \quad=(-1)^{f(0)} \frac{1}{2}\left(|0\rangle+(-1)^{f(0) \oplus f(1)}|1\rangle\right)(|0\rangle-|1\rangle)
$$



## Clash: Haskell for FPGA

- C入aSH http://www.clash-lang.org/
- A Haskell spin-off
- Models wires as infinite stream, and sequential logic as State machines
- counter :: Signal (Unsigned 2)
- counter = register 0 (liftA (+1) topEntity)
- > sampleN $8 \$$ topEntity
- $[0,1,2,3,0,1,2,3]$
- Dependent type for bit width (partial support)
- Type checking for bit width checking
- (++) :: Vec n a -> Vec mat> $\operatorname{Vec}(\mathrm{n}+\mathrm{m}) \mathrm{a}$


## Clash: Haskell for FPGA

```
mealy :: (s -> i -> (s, 0)) -> s -> Signal i -> Signal o
mac :: Int -- Current state
    -> (Int,Int) -- Input
    -> (Int,Int) -- (Updated state, output)
mac s (x,y) = (s',s)
    where s' = x * y + s
```

topEntity :: Signal (Int, Int) -> Signal Int
topEntity $=$ mealy mac 0


## Clash/FPGA: implement Complex Number

```
type CC = Vec 2 RR
c0 = 0 :> 0 :> Nil
c1 = 1 :> 0 :> Nil
sqr_norm :: CC -> RR
sqr_norm (a :> b :> Nil) = a * a + b * b
cadd :: CC -> CC -> CC
cadd = zipWith (+)
cmul :: CC -> CC -> CC
cmul (a :> b :> Nil) (c : > d :> Nil) = (a * c - b * d) :> (a * d + b * c) :>
Nil
dotProduct xs ys = foldr cadd c0 (zipWith cmul xs ys)
matrixVector m v = map (`dotProduct` v) m
```


## Clash/FPGA: Qubit

```
type QBit = Vec 2 CC
q0 :: Signal QBit
q0 = register (c1 :> c0 :> Nil) q0
q1 :: Signal QBit
q1 = register (c0 :> c1 :> Nil) q1
qPlus = hadamardG q0
qMinus = hadamardG q1
    ( hr ch
hadamard :: QBit -> QBit
hadamard = matrixVector ((h
```



```
hadamardG :: Signal QBit -> Signal QBit 
measure :: Signal QBit -> Signal RR
measure = register 0 . liftA (\ x -> sqr_norm (x !! 1))
```


## Multi-Qubit interaction

From $(a|0>+b| 1>)(c|0>d| 1>)$ to $a c|00>+a d| 01>+b c|10>+b d| 11>$

```
explode :: Signal QBit -> Signal QBit -> Signal (Vec 4 CC)
```

explode qx $q y=$ register (repeat c0) \$ liftA2 outer qx qy
where

```
        outer :: QBit -> QBit -> Vec 4 CC
        outer (x0 :> x1 :> Nil) y = (map (cmul x0) y) ++ (map (cmul x1) y)
```

measure0 : : Signal (Vec 4 CC) -> Signal RR

Measures $\left|\left|00>\left.\right|^{2}+\left||01>|^{2}\right.\right.\right.$

## Deutsch-Jozsa's algorithm


deutsch $u$ : : Vec 2 RR $->$ Vec 4 CC $->$ Vec 4 CC deutsch $u(f 0:>$ fl $:>$ Nil) $=$ matrixVector (make complex (

$$
\left(\begin{array}{cccc}
1-f 0 & f 1 & 0 & 0 \\
f 0 & 1-f 1 & 0 & 0 \\
0 & 0 & 1-f 0 & f 1 \\
0 & 0 & f 0 & 1-f 1
\end{array}\right)
$$

hadamard I : Vec 4 CC -> Vec 4 CC
hadamard I =

```
matrixVector (make_complex (
                    (h :> 0 :> h :> 0 :> Nil) :>
(0 :> h :> 0 :> h :> Nil) :>
(h :> 0 :> - h :> 0 :> Nil) :>
(0 :> h :> 0 :> - h :> Nil) :> Nil))
```

$H \otimes I=\left(\begin{array}{cccc}h & 0 & h & 0 \\ 0 & h & 0 & h \\ h & 0 & -h & 0 \\ 0 & h & 0 & -h\end{array}\right)$
where $h=\$ \$(f L i t(1 /$ sqrt 2$)):: R R$

$$
\begin{aligned}
& (\mathrm{fO}:>(1-\mathrm{fl}):>0:>0:>\mathrm{Nil}):>\quad 1\left(\begin{array}{lllll}
0 & 0 & f 0 & 1-f 1
\end{array}\right) \\
& \text { (0 : > } 0:>(1-\mathrm{f} 0):>\mathrm{fl}:>\mathrm{Nil}):> \\
& \text { ( } 0:>0 \quad:>\mathrm{fO}:>(1-\mathrm{fl}) \quad:>\mathrm{Nil}):>\mathrm{Nil}) \text { ) }
\end{aligned}
$$

## Deutsch-Jozsa's algorithm



```
deutsch :: Vec 2 RR -> Signal RR
deutsch f0f1 =
    let xy = explode qPlus qMinus in
    let xy2 = register (repeat c0) $ liftA (deutsch_u f0f1) xy in
    let xy3 = register (repeat c0) $ liftA hadamard_I xy2 in
    measure0 xy3
```

topEntity : : Signal (Vec 4 RR)
topEntity $=$ bundle (map deutsch (f0 :> f1 :> f2 :> f3 :> Nil))
where $f 0=0$ :> $0 \quad:>$ Nil
f1 = 1 :> 1 :> Nil
f2 = 0 :> 1 :> Nil
f3 = 1 :> 0 :> Nil
sampleN 8 \$ topentity
$[<0.0,0.0,0.0,0.0>,<0.0,0.0,0.0,0.0>,<0.0,0.0,0.0,0.0>,<0.0,0.0,0.0,0.0>,<0.0,0.0,0.0,0.0\rangle$ $,<0.999847412109375,0.999847412109375,0.0,0.0>,<0.999847412109375,0.999847412109375,0.0,0$.
$0>,<0.999847412109375,0.999847412109375,0.0,0.0>$ ]

## Synthesizing on FPGA

yosys> show
Deutsch_explode

Problem: no usage of ALU, very resource intensive.

## Future work

- Try more Quantum Computing algorithms
- Do the matrix multiplications in multiple cycles


## Congratulation for becoming one of the rarest

 species!

Ross Freeman

## Backup after this slide

## Quantum Computing

- Qubit
- Inherently reversible
- Quantum coherence exploits entanglement
- Quantum Decoherence
- Introduction to Quantum Information


## Schmidt decomposition (yet another SVD)

- vector w in tensor product space H_1 lotimes H_2
- separable state
- entangled state
- Schmidt rank

$$
w=\sum_{i=1}^{m} \alpha_{i} u_{i} \otimes v_{i} .
$$

- Schmidt decomposition
- Partial trace
- von Neumann entropy

$$
-\sum_{i}\left|\alpha_{i}\right|^{2} \log \left|\alpha_{i}\right|^{2}
$$

- matrix w with first dimension being H_1 and second being H_2
- rank 1 matrix
- $\operatorname{rank}(w)>1$
- rank
- SVD: w = U S V^T
- $S^{\wedge} 2$
- entropy of square of singular values


## A Finite Input Response Filter in Clash

```
dotp :: SaturatingNum a
    => Vec (n + 1) a
    -> Vec (n + 1) a
    -> a
```


dotp as bs = fold boundedPlus (zipWith boundedMult as bs)
fir
:: (Default a, KnownNat n, SaturatingNum a, HasClockReset domain gated synchronous)
=> Vec $(\mathrm{n}+1) \mathrm{a}$-> Signal domain a -> Signal domain a
fir coeffs $x_{-} t=y \_t$
where

```
y_t = dotp coeffs <$> bundle xs
xs = window x_t
```

