

Random Graphs

CS224W

Network models

- Why model?
 - simple representation of complex network
 - can derive properties mathematically
 - predict properties and outcomes

 - Also: to have a strawman
 - In what ways is your real-world network different from hypothesized model?
 - What insights can be gleaned from this?
-

Downloading NetLogo

- <https://ccl.northwestern.edu/netlogo/>
- Models specific to this class:
<http://web.stanford.edu/class/cs224w/NetLogo/>



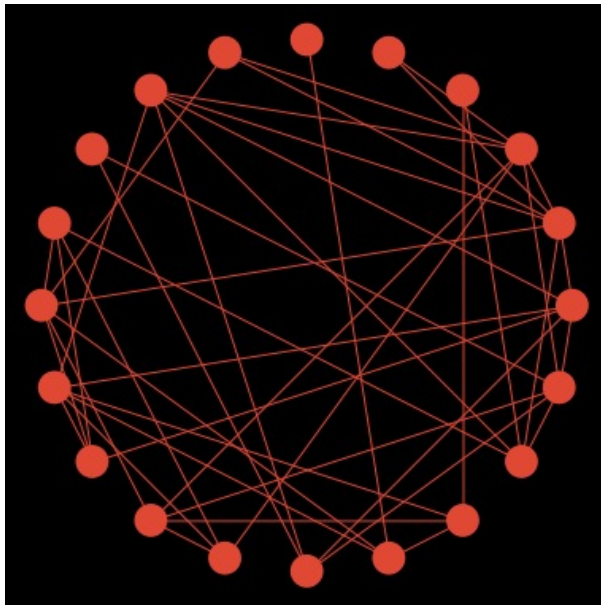
Erdős and Rényi



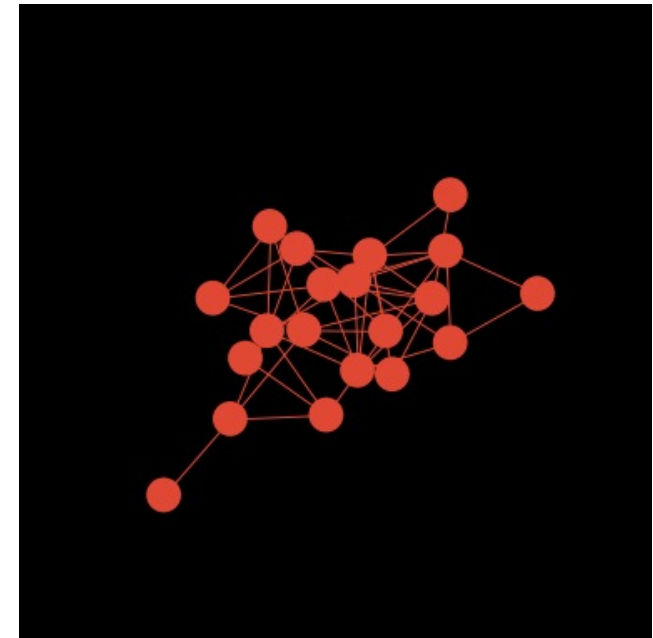
Erdős-Renyi: simplest network model

- ▣ Assumptions
 - ▣ nodes connect at random
 - ▣ network is undirected
- ▣ Key parameter (besides number of nodes N) : p or M
 - ▣ p = probability that any two nodes share an edge
 - ▣ M = total number of edges in the graph

what they look like



after spring
layout



Degree distribution

- (N,p)-model: For each potential edge we flip a biased coin
 - with probability p we add the edge
 - with probability $(1-p)$ we don't
- Alternate notation: G_{np}

Quiz Q:

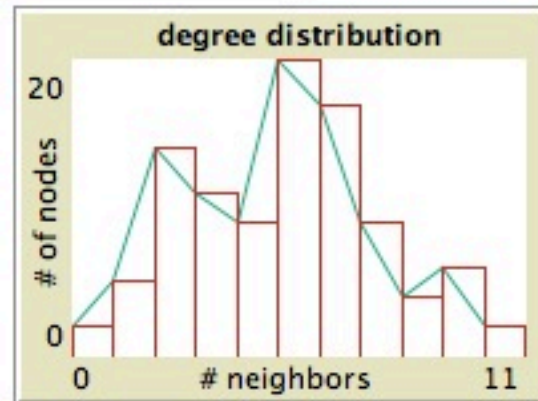
- As the size of the network increases, if you keep p , the probability of any two nodes being connected, the same, what happens to the average degree
 - a) stays the same
 - b) increases
 - c) decreases

<http://web.stanford.edu/class/cs224w/NetLogo/ErdosRenyiDegDist.nlogo>

num-nodes 100 GC size 98 av. deg 4.84

On prob-or-num? num-neighbors 2

Erdoes-Renyi prob-link 0.050



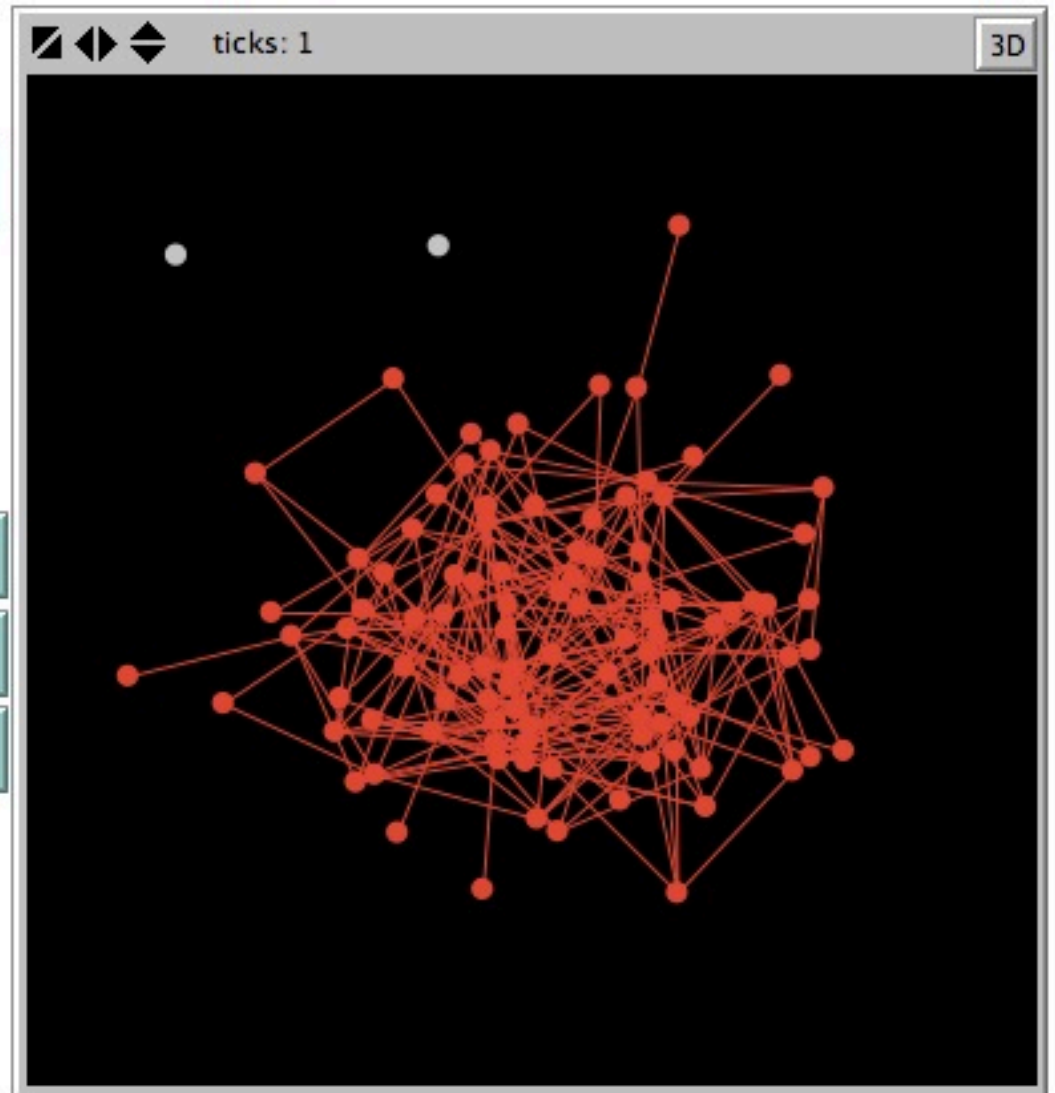
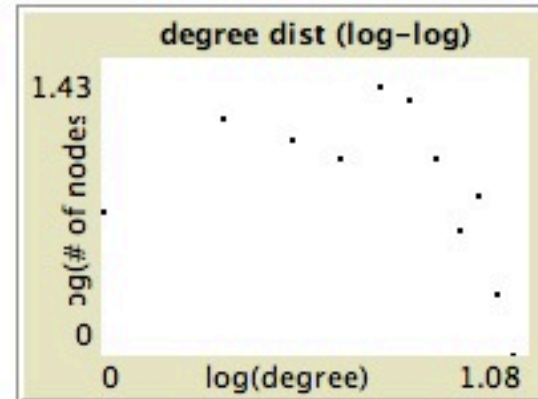
redo layout

layout options

spring-cons... 0.2

repulsion-st... 0.2

spring-length 15



<http://web.stanford.edu/class/cs224w/NetLogo/ErdoesRenyiDegDist.nlogo>

Degree distribution

- What is the probability that a node has 0,1,2,3... edges?
 - Probabilities sum to 1
-

How many edges per node?

- Each node has $(N - 1)$ tries to get edges
- Each try is a success with probability p
- The binomial distribution gives us the probability that a node has degree k :

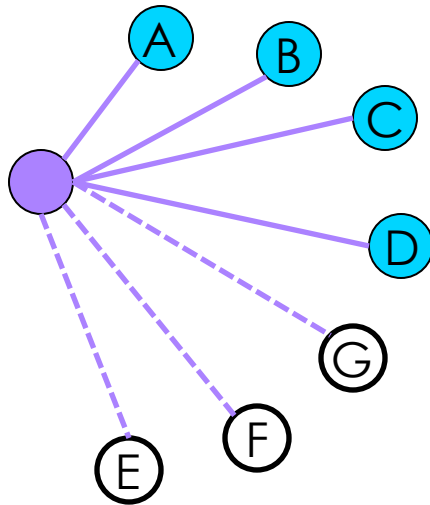
$$B(N - 1; k; p) = \binom{N - 1}{k} p^k (1 - p)^{N - 1 - k}$$

Quiz Q:

- The maximum degree of a node in a simple (no multiple edges between the same two nodes) N node graph is
 - a) N
 - b) $N - 1$
 - c) $N / 2$

Explaining the binomial distribution

- 8 node graph, probability p of any two nodes sharing an edge
- What is the probability that a given node has degree 4?



Binomial coefficient: choosing 4 out of 7

Suppose I have 7 blue and white nodes, each of them uniquely marked so that I can distinguish them. The blue nodes are ones I share an edge with, the white ones I don't.



How many different samples can I draw containing the same nodes but in a different order (the order could be e.g. the order in which the edges are added (or not)? e.g.



binomial coefficient explained

Ⓒ Ⓔ Ⓒ Ⓓ Ⓑ Ⓕ Ⓐ

If order matters, there are **7!** different orderings:

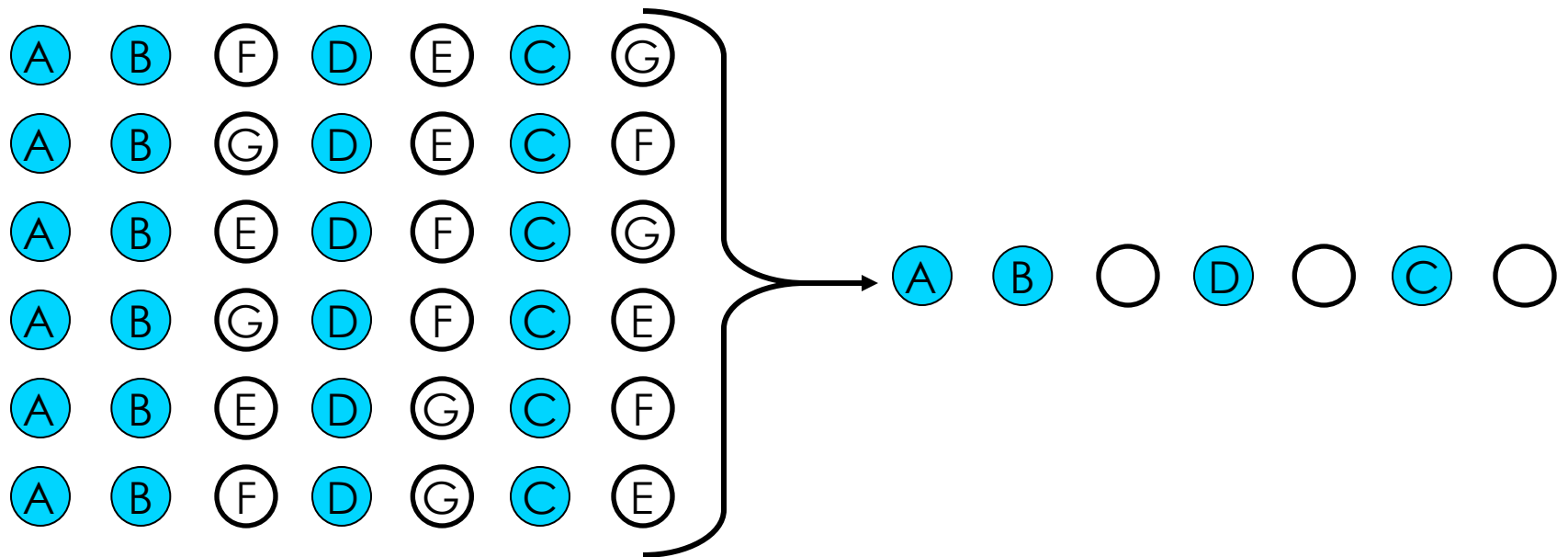
I have 7 choices for the first spot, 6 choices for the second (since I've picked 1 and now have only 6 to choose from),
5 choices for the third, etc.

$$7! = 7 * 6 * 5 * 4 * 3 * 2 * 1$$

binomial coefficient

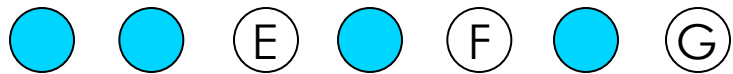
Suppose the order of the nodes I don't connect to (white) doesn't matter.

All possible arrangements ($3!$) of white nodes look the same to me.



Instead of $7!$ combinations, we have $7!/3!$ combinations

binomial coefficient explained



The same goes for the blue nodes, if we can't tell them apart, we lose a factor of 4!

binomial coefficient explained

number of ways of choosing k items out of $(n-1)$

$$\begin{aligned} & \text{number of ways of arranging } \mathbf{n-1} \text{ items} \\ = & \frac{\text{number of ways of arranging } \mathbf{n-1} \text{ items}}{(\# \text{ of ways to arrange } \mathbf{k} \text{ things}) * (\# \text{ ways to arrange } \mathbf{n-1-k} \text{ things})} \\ = & \frac{\mathbf{n-1!}}{\mathbf{k! (n-1-k)!}} \end{aligned}$$

Note that the binomial coefficient is symmetric – there are the same number of ways of choosing k or $n-1-k$ things out of $n-1$

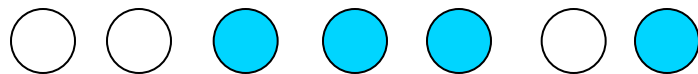
Quiz Q:

- What is the number of ways of choosing 2 items out of 5?
 - 10
 - 120
 - 6
 - 5

Now the distribution

- p = probability of having edge to node (blue)
- $(1-p)$ = probability of not having edge (white)
- The probability that you connect to 4 of the 7 nodes in some particular order (two white followed by 3 blues, followed by a white followed by a blue) is

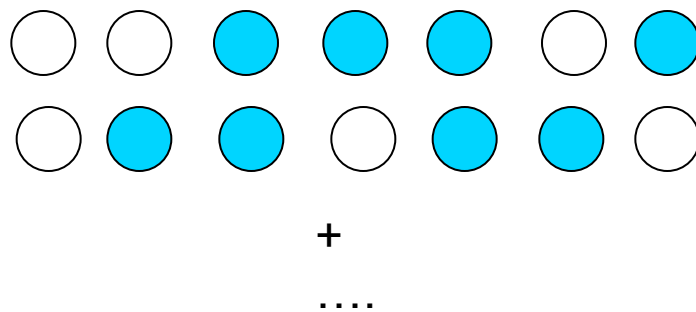
$$P(\text{white}) * P(\text{white}) * P(\text{blue}) * P(\text{blue}) * P(\text{blue}) * P(\text{white}) * P(\text{blue}) \\ = p^4 * (1-p)^3$$



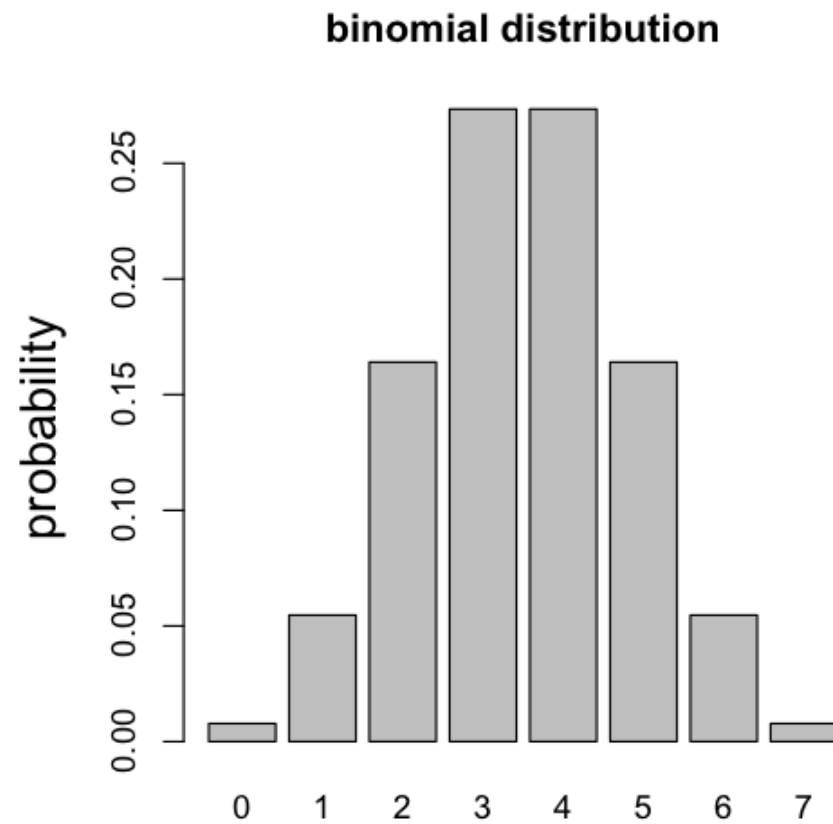
Binomial distribution

- If order doesn't matter, need to multiply probability of any given arrangement by number of such arrangements:

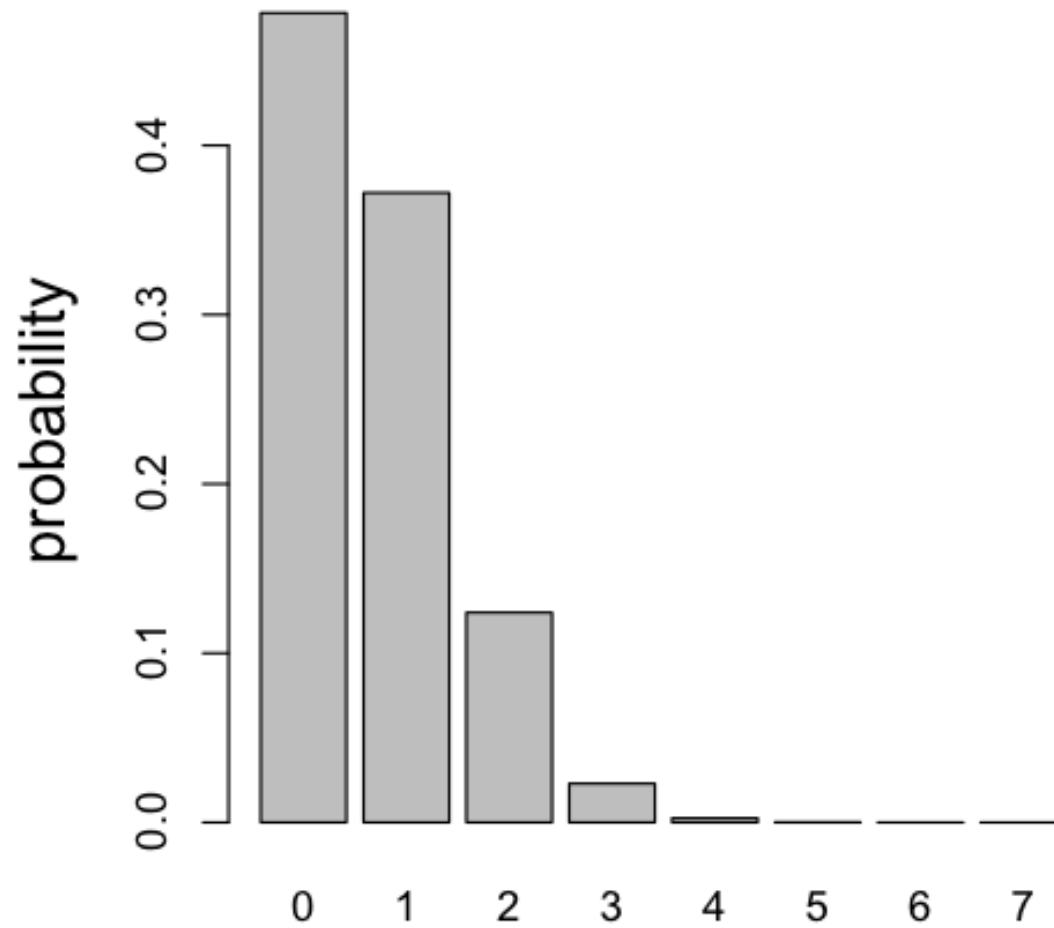
$$B(7; 4; p) = \binom{7}{4} p^4 (1-p)^3$$



if $p = 0.5$



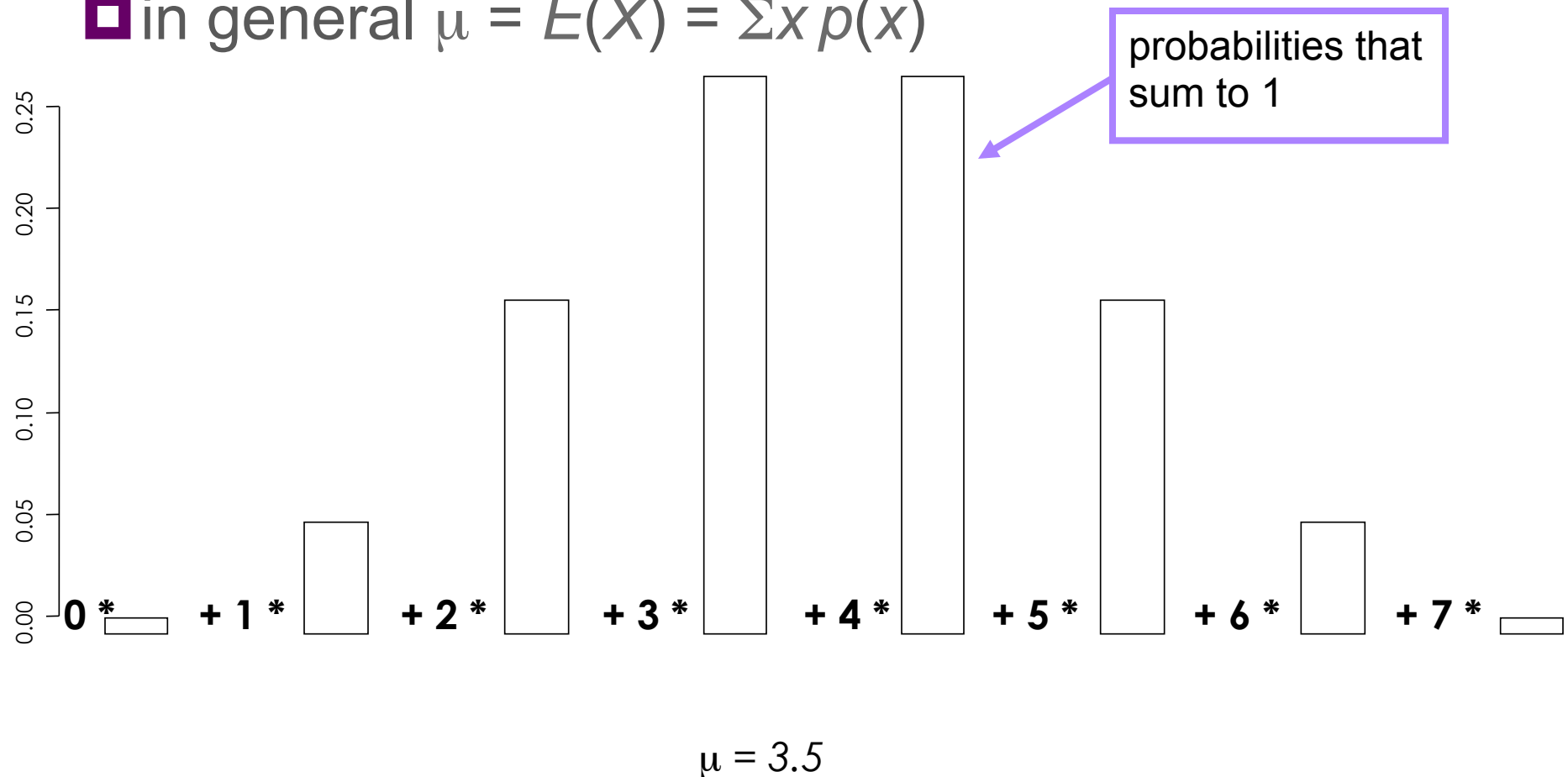
$$p = 0.1$$



What is the mean?

□ Average degree $\langle k \rangle = z = (n-1) * p$

□ in general $\mu = E(X) = \sum x p(x)$

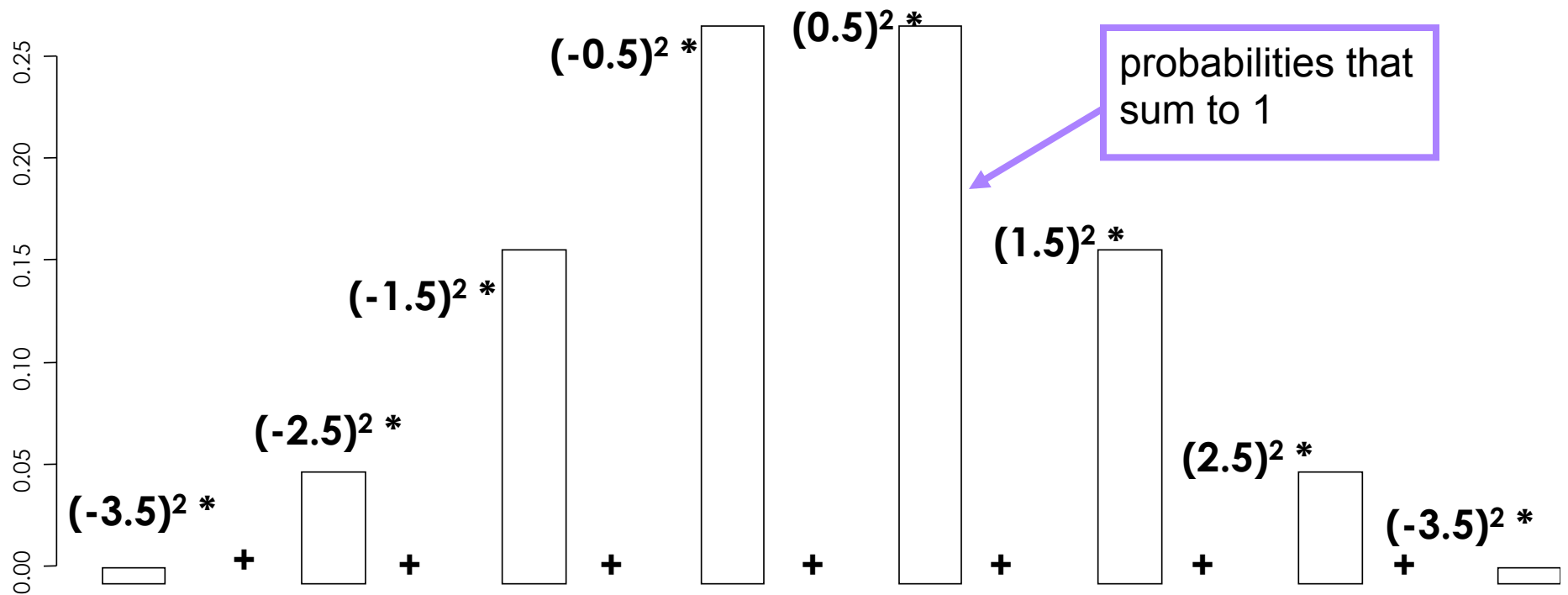


Quiz Q:

- What is the average degree of a graph with 10 nodes and probability $p = 1/3$ of an edge existing between any two nodes?
 - 1
 - 2
 - 3
 - 4

What is the variance?

- variance in degree
 $\sigma^2 = (n-1) * p * (1-p)$
- in general $\sigma^2 = E[(X-\mu)^2] = \sum (x-\mu)^2 p(x)$



Approximations

$$p_k = \binom{n-1}{k} p^k (1-p)^{n-1-k}$$

$$p_k = \frac{z^k e^{-z}}{k!}$$

$$p_k = \frac{1}{\sigma\sqrt{2\pi}} e^{-\frac{(k-z)^2}{2\sigma^2}}$$

Binomial



limit p small

Poisson

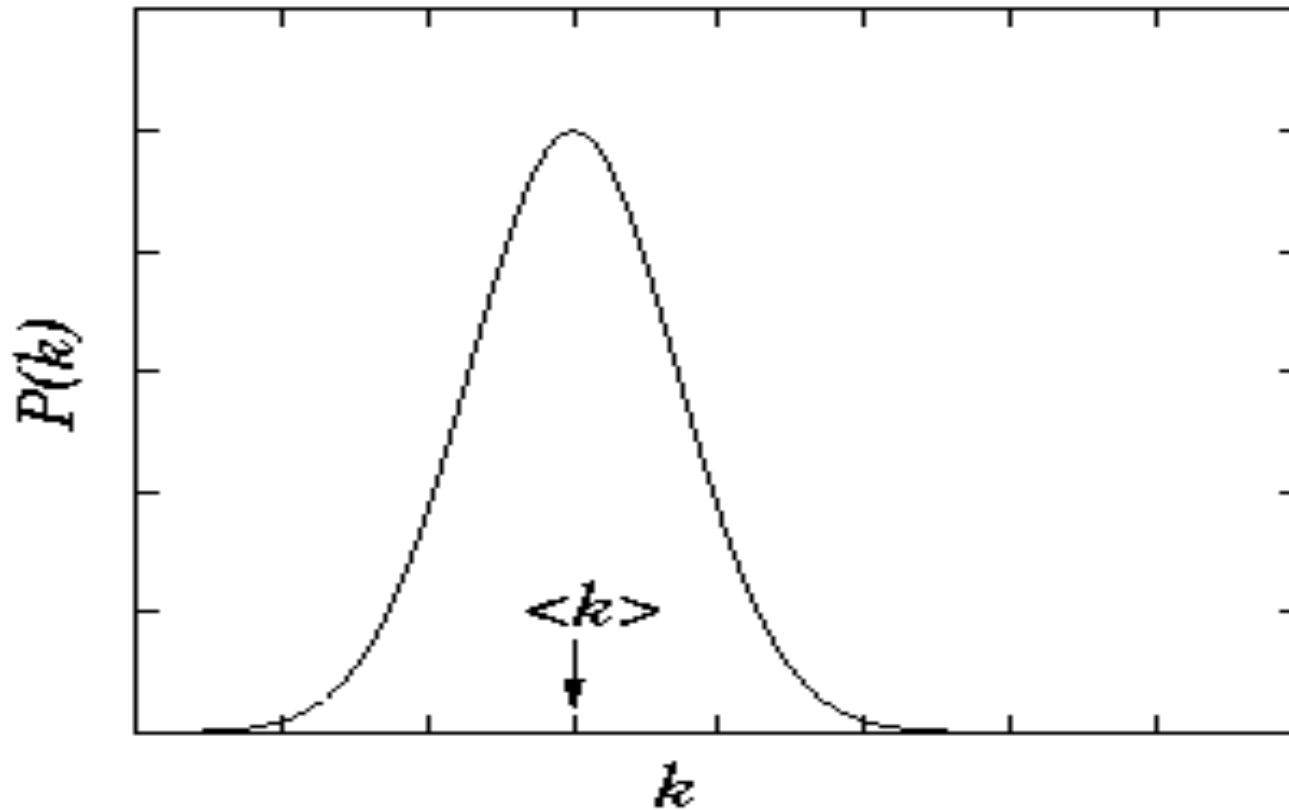


limit large n

Normal

Poisson distribution

Poisson distribution



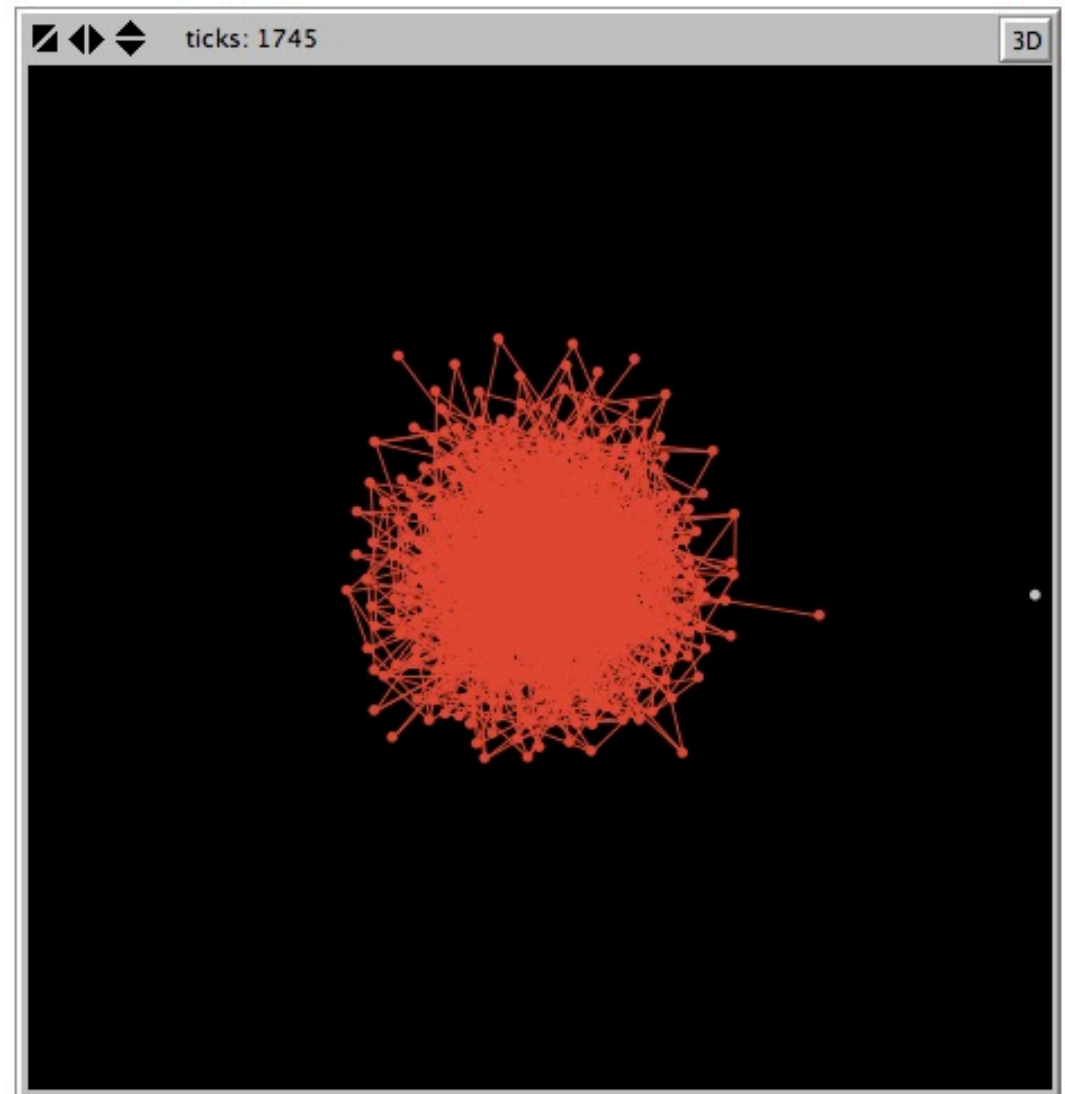
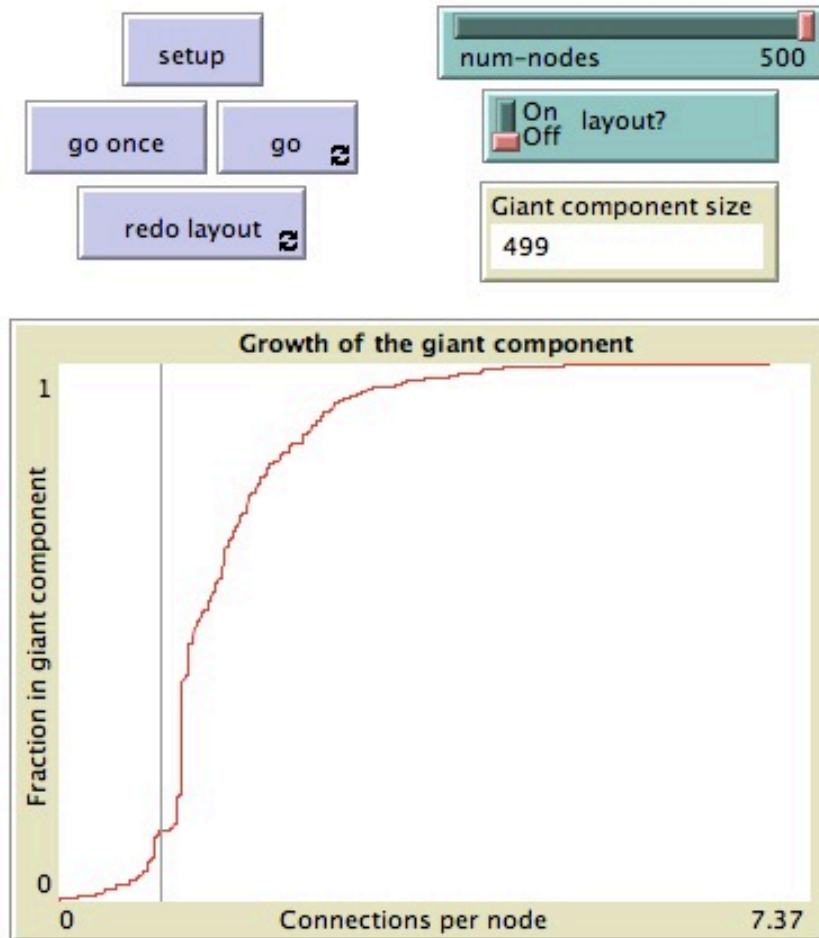
What insights does this yield? No hubs

- You don't expect large hubs in the network

Insights

- ❑ Previously: degree distribution / absence of hubs
 - ❑ Emergence of giant component
 - ❑ Average shortest path
-

Emergence of the giant component

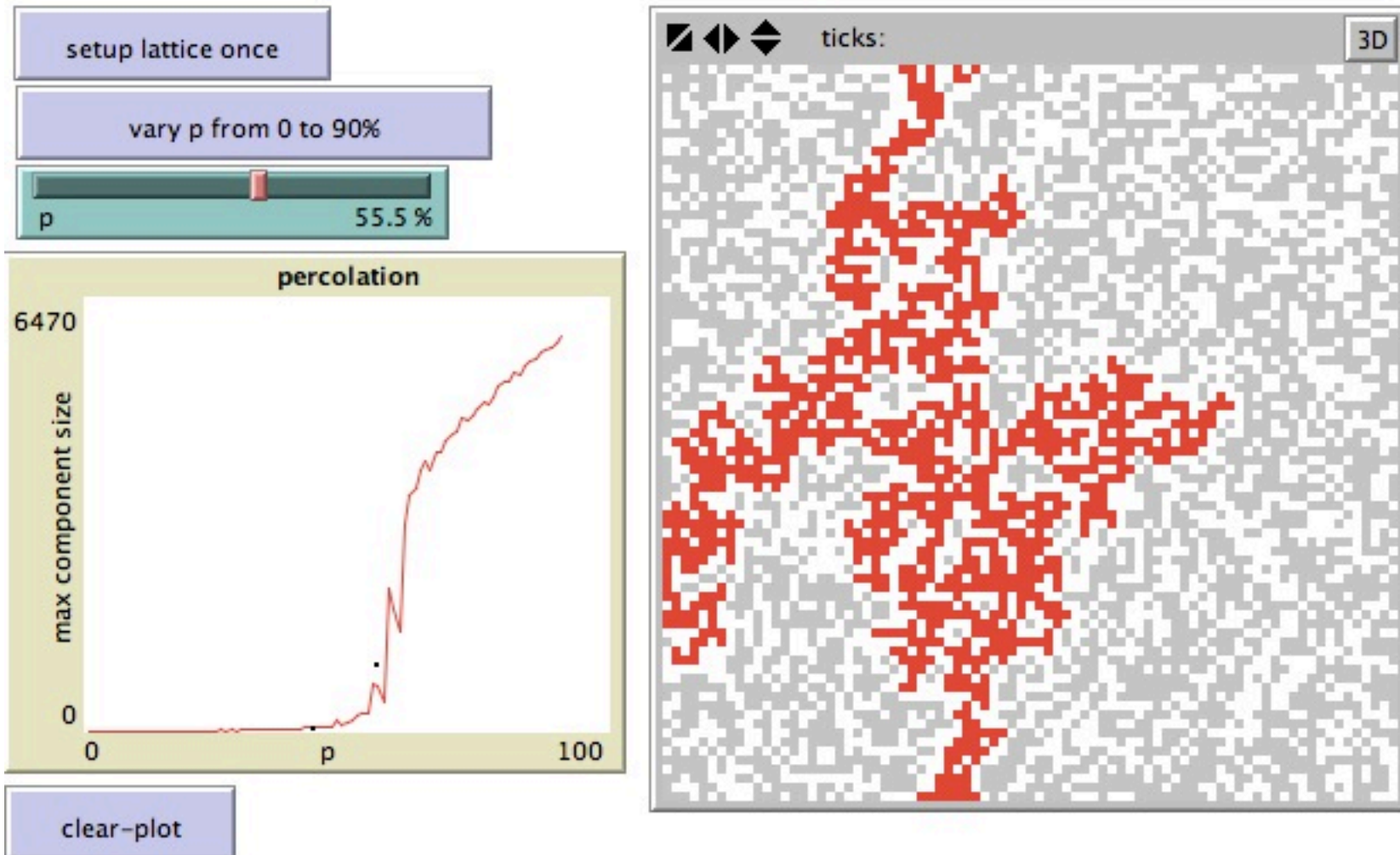


(standard model in NetLogo library) <http://ccl.northwestern.edu/netlogo/models/GiantComponent>

Quiz Q:

- What is the average degree z at which the giant component starts to emerge?
 - 0
 - 1
 - $3/2$
 - 3

Percolation on a 2D lattice

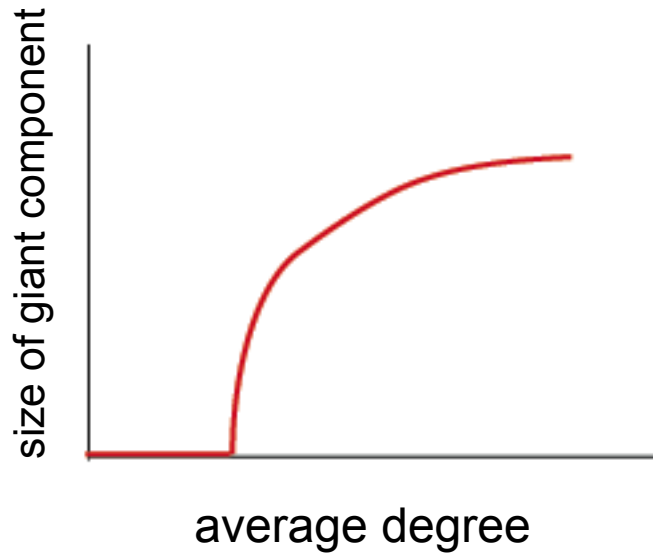


<http://web.stanford.edu/class/cs224w/NetLogo/LatticePercolation.nlogo>

Quiz Q:

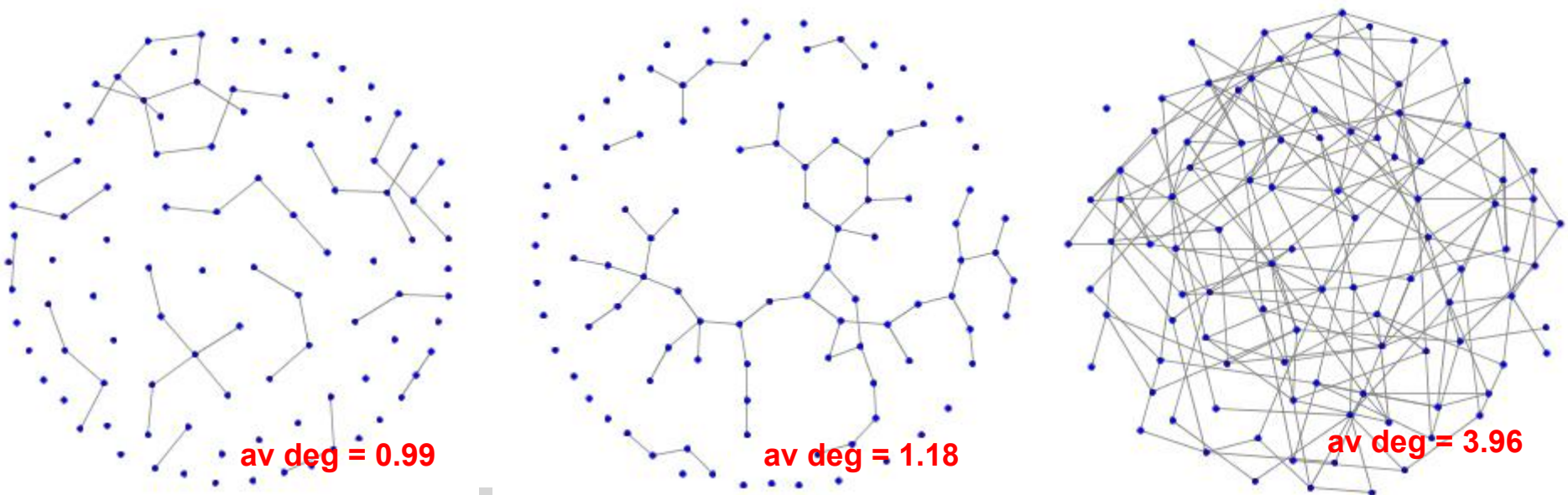
- What is the percolation threshold of a 2D lattice: fraction of sites that need to be occupied in order for a giant connected component to emerge?
 - 0
 - $\frac{1}{4}$
 - $\frac{1}{3}$
 - $\frac{1}{2}$

Percolation threshold



Percolation threshold: how many edges need to be added before the giant component appears?

As the average degree increases to $z = 1$, a giant component suddenly appears



“Evolution” of the G_{np}

What happens to G_{np} when we vary p ?

Back to Node Degrees of G_{np}

- Remember, expected degree $E[X_v] = (n-1)p$
- If want $E[X_v]$ be independent of n

let: $p=c/(n-1)$

Probability of a node being isolated

▣ **Observation:** If we build random graph G_{np} with $p=c/(n-1)$ we have many isolated nodes

▣ **Why?**

$$P[v \text{ has degree } 0] = (1 - p)^{n-1} = \left(1 - \frac{c}{n-1}\right)^{n-1} \xrightarrow{n \rightarrow \infty} e^{-c}$$

$$\lim_{n \rightarrow \infty} \left(1 - \frac{c}{n-1}\right)^{n-1} = \left(1 - \frac{1}{x}\right)^{-x \cdot c} = \left[\lim_{x \rightarrow \infty} \left(1 - \frac{1}{x}\right)^{-x} \right]^{-c} = e^{-c}$$

Use substitution $\frac{1}{x} = \frac{c}{n-1}$

e (by definition)

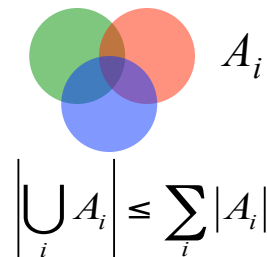
No Isolated Nodes

- How big do we have to make p before we are likely to have no isolated nodes?
- We know: $P[v \text{ has degree } 0] = e^{-c}$
- Event we are asking about is:
 - $I =$ some node is isolated
 - $I = \bigcup_{v \in N} I_v$ where I_v is the event that v is isolated

□ We have:

$$P(I) = P\left(\bigcup_{v \in N} I_v\right) \leq \sum_{v \in N} P(I_v) = ne^{-c}$$

Union bound



No Isolated Nodes

□ We just learned: $P(I) = n e^{-c}$

□ Let's try:

□ $c = \ln n$ then: $n e^{-c} = n e^{-\ln n} = n \cdot 1/n = 1$

□ $c = 2 \ln n$ then: $n e^{-2 \ln n} = n \cdot 1/n^2 = 1/n$

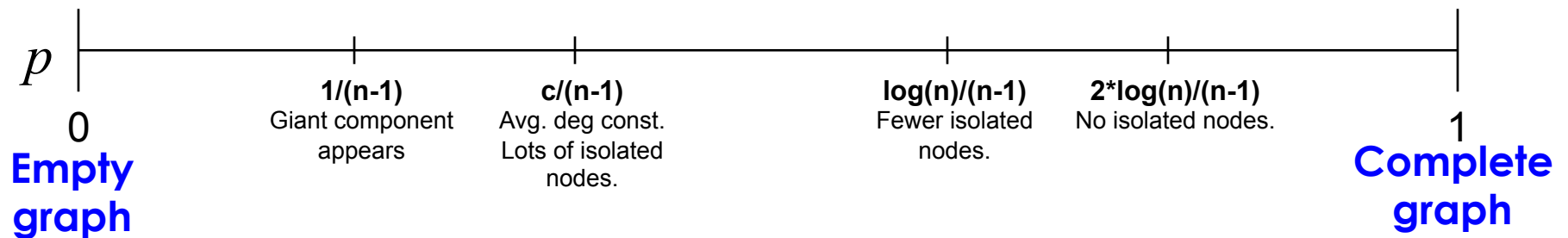
□ So if:

□ $p = \ln n$ then: $P(I) = 1$

□ $p = 2 \ln n$ then: $P(I) = 1/n \rightarrow 0$ as $n \rightarrow \infty$

“Evolution” of a Random Graph

Graph structure of G_{np} as p changes:



Emergence of a Giant Component:

avg. degree $k=2E/n$ or $p=k/(n-1)$

□ $k=1-\varepsilon$: all components are of size $\Omega(\log n)$

□ $k=1+\varepsilon$: 1 component of size $\Omega(n)$, others have size $\Omega(\log n)$

Giant component – another angle

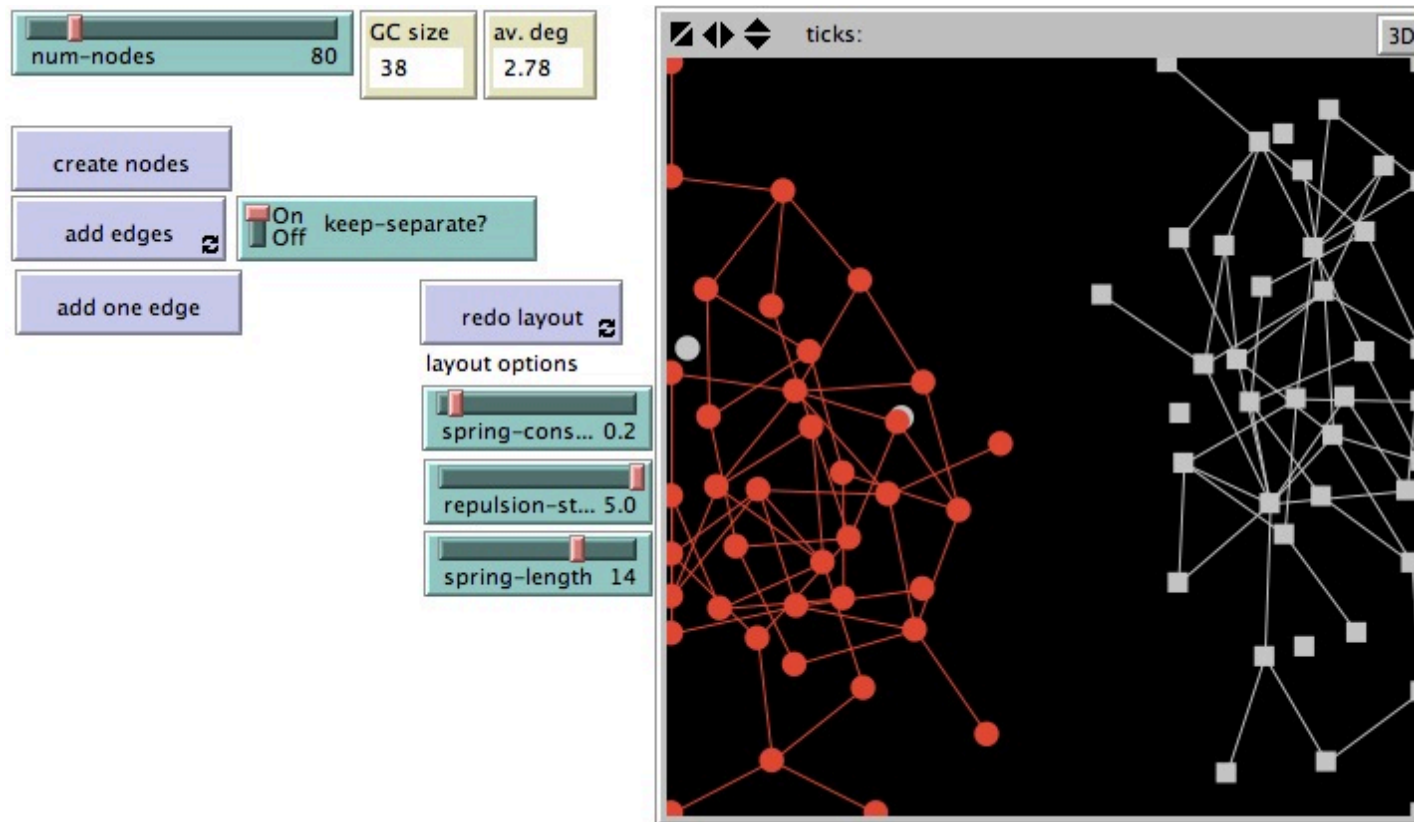
- How many other friends besides you does each of your friends have?
 - By property of degree distribution
 - the average degree of your friends, you excluded, is z
 - so at $z = 1$, each of your friends is expected to have another friend, who in turn have another friend, etc.
 - the giant component emerges
-

Giant component illustrated



Why just one giant component?

- What if you had 2, how long could they be sustained as the network densifies?



<http://web.stanford.edu/class/cs224w/NetLogo/ErdosRenyiTwoComponents.nlogo>

Quiz Q:

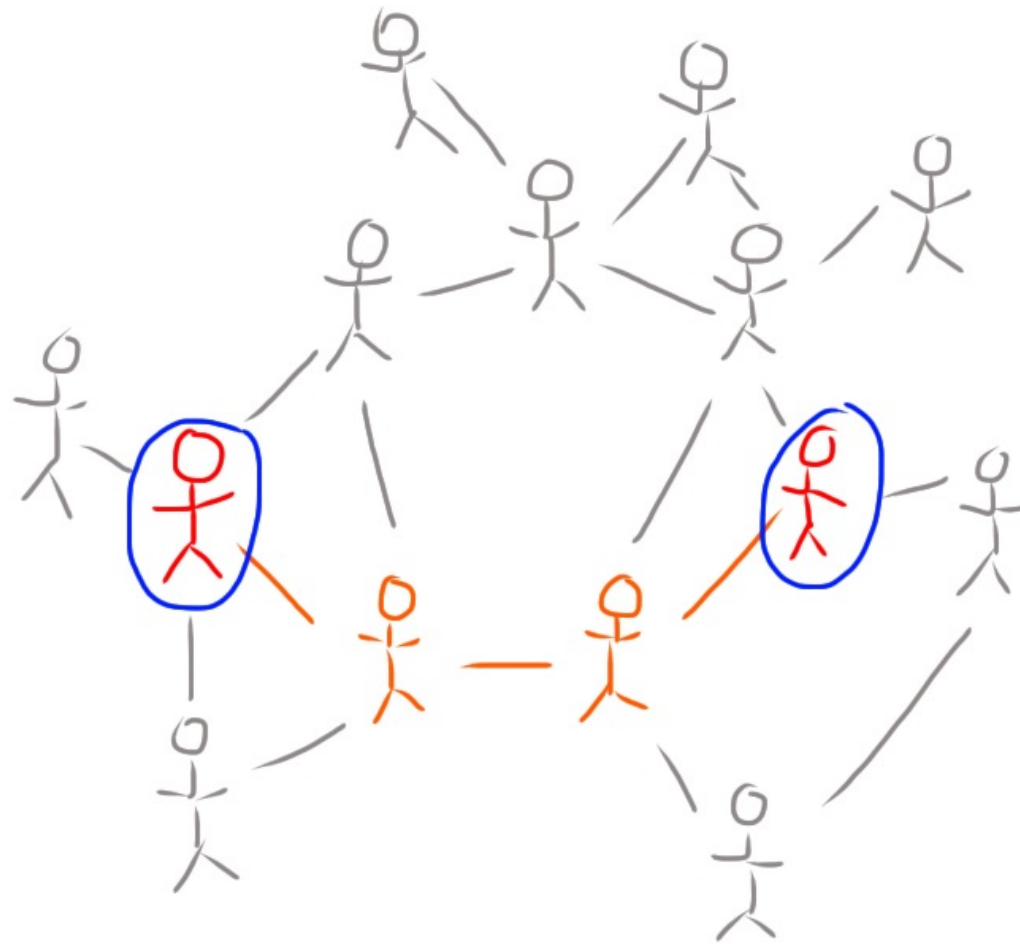
- If you have 2 large-components each occupying roughly $1/2$ of the graph, how long does it typically take for the addition of random edges to join them into one giant component
 - 1-4 edge additions
 - 5-20 edge additions
 - over 20 edge additions

Average shortest path

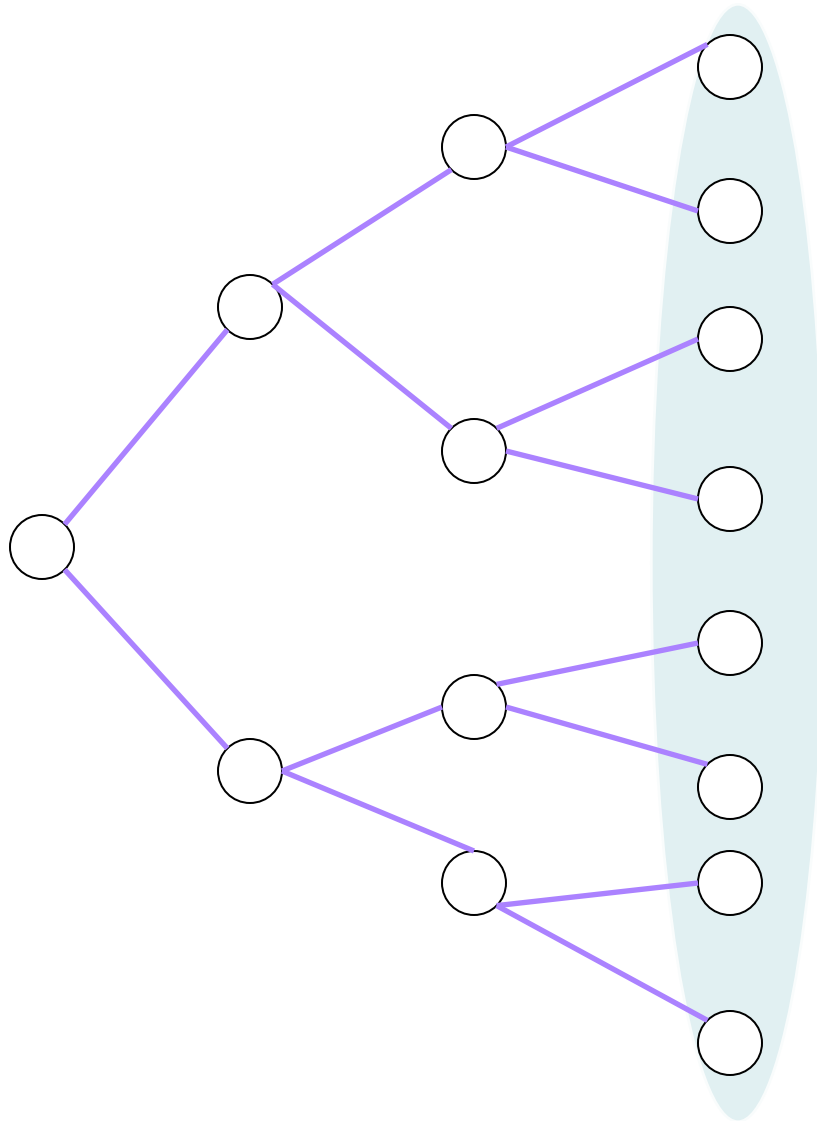
- How many hops on average between each pair of nodes?
- again, each of your friends has **$z = \text{avg. degree}$** friends besides you
- ignoring loops, the number of people you have at distance l is

z^l

Average shortest path



friends at distance l



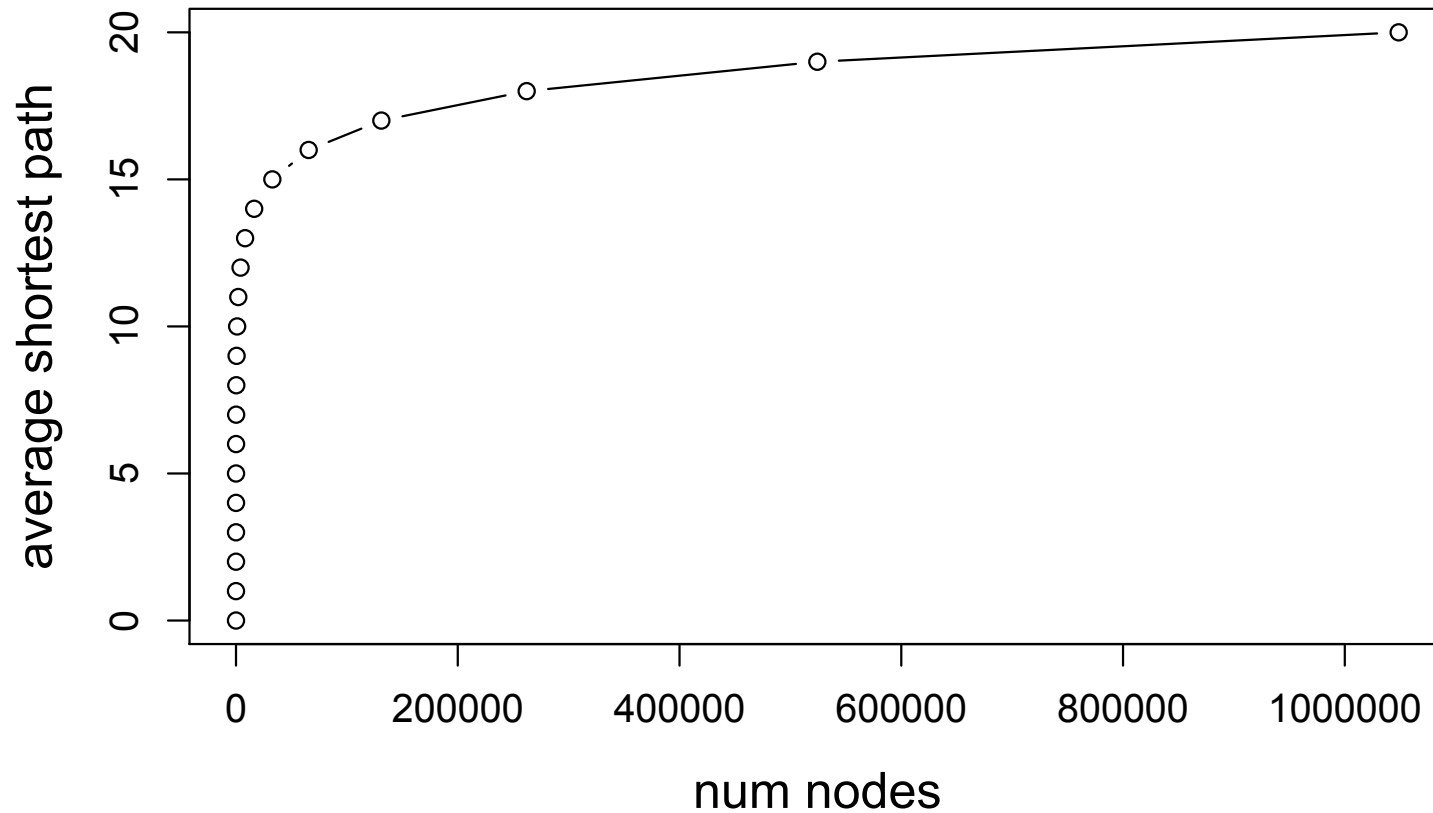
$$N_l = z^l$$

scaling:
average shortest path l_{av}

$$l_{av} \sim \frac{\log N}{\log z}$$

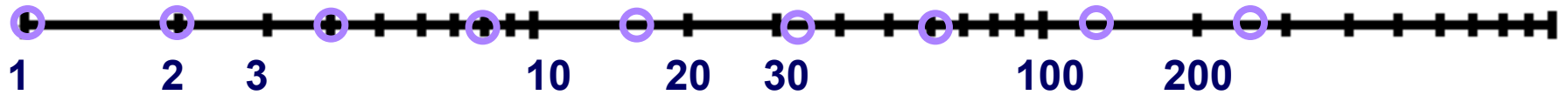
What this means in practice

- Erdős-Renyi networks can grow to be very large but nodes will be just a few hops apart



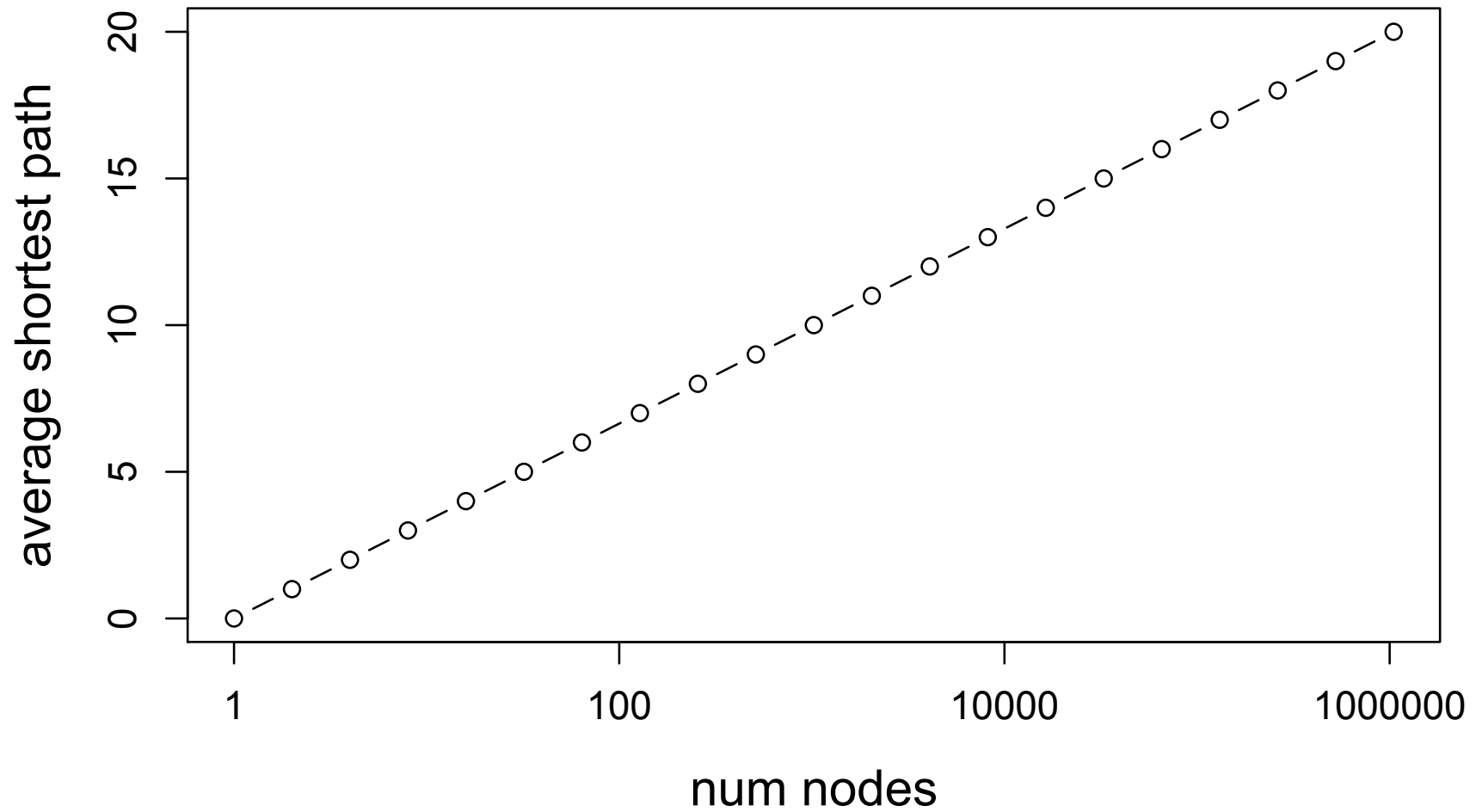
Logarithmic axes

- powers of a number will be uniformly spaced



- $2^0=1, 2^1=2, 2^2=4, 2^3=8, 2^4=16, 2^5=32, 2^6=64, \dots$

Erdős-Renyi avg. shortest path



Quiz Q:

- If the size of an Erdős-Renyi network increases 100 fold (e.g. from 100 to 10,000 nodes), how will the average shortest path change
 - it will be 100 times as long
 - it will be 10 times as long
 - it will be twice as long
 - it will be the same
 - it will be 1/2 as long

Realism

- ❑ Consider alternative mechanisms of constructing a network that are also fairly “random”.
 - ❑ How do they stack up against Erdős-Renyi?
 - ❑ <http://web.stanford.edu/class/cs224w/NetLogo/RandomGraphs.nlogo>
-

Introduction model

- Prob-link is the p (probability of any two nodes sharing an edge) that we are used to
 - But, with probability prob-intro the other node is selected among one of our friends' friends and not completely at random
-

Introduction model

num-nodes 50 GC size 50 av. deg 5.72

On prob-or-num?
 Off

Erdos-Renyi prob-link 0.20

introduction prob-intro 0.83

static-geo

growth num-neighbors 1

rand-encounter

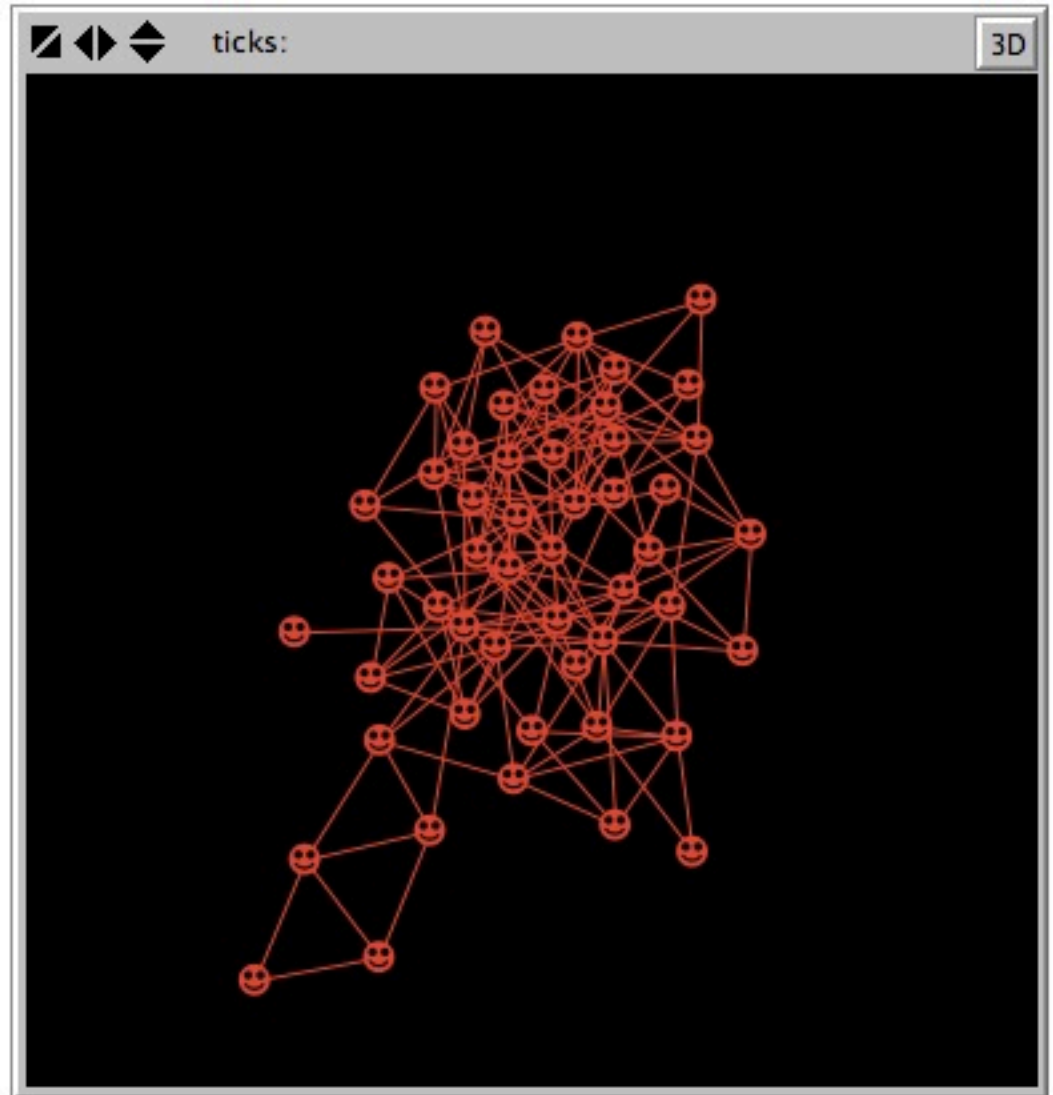
redo layout

layout options

spring-length 11

spring-constant 0.1

repulsion-strength 2.8



Quiz Q:

- Relative to ER, the introduction model has:
 - more edges
 - more closed triads
 - longer average shortest path
 - more uneven degree
 - smaller giant component at low p
-

Static Geographical model

- Each node connects to num-neighbors of its closest neighbors
 - use the num-neighbors slider, and for comparison, switch PROB-OR-NUM to 'off' to have the ER model aim for num-neighbors as well
 - turn off the layout algorithm while this is running, you can apply it at the end
-

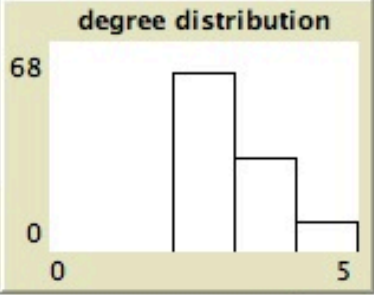
static geo

num-nodes 99

On prob-or-num?
 Off

GC size	av. deg	ASP
47	2.55	9.3

degree distribution



ticks: 1 3D

Erdos-Renyi

prob-link 0.02

introduction

prob-intro 0.83

static-geo

growth

num-neighbors 2

rand-encounter

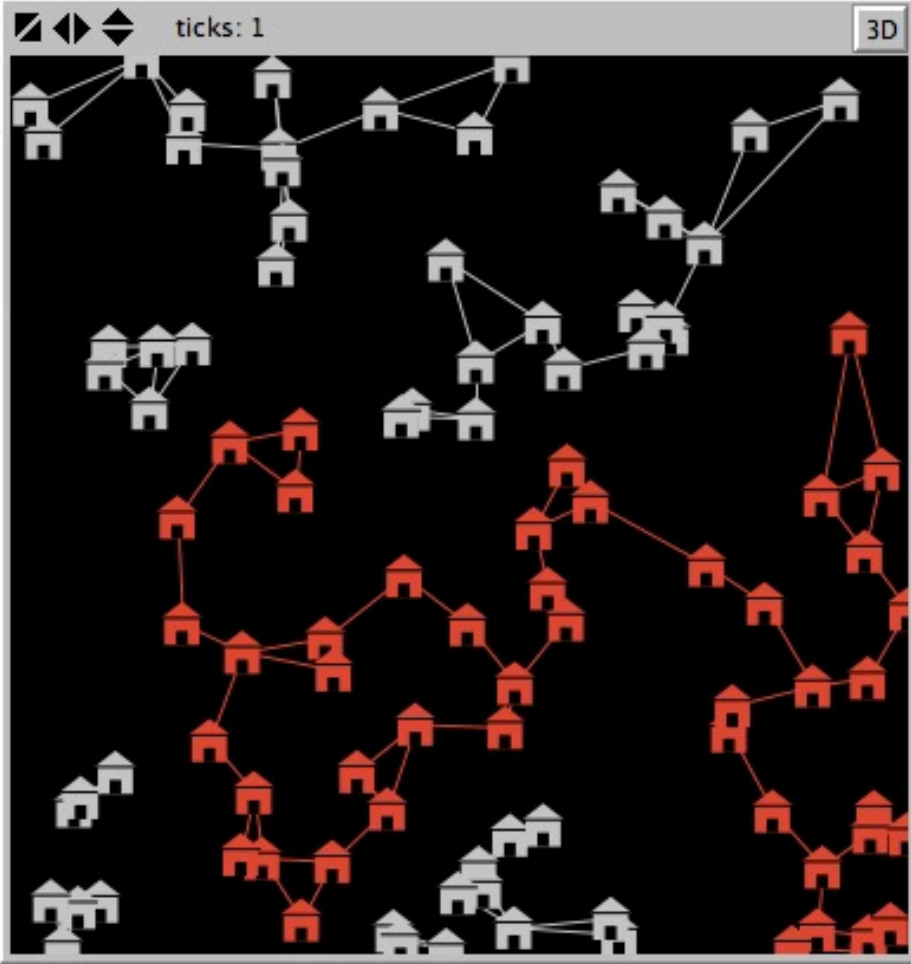
redo layout

layout options

spring-length 6

spring-constant 0.1

repulsion-strength 0.4



Quiz Q:

- Relative to ER, the static geographical model has :
 - longer average shortest path
 - shorter average shortest path
 - narrower degree distribution
 - broader degree distribution
 - smaller giant component at a low number of neighbors
 - larger giant component at a low number of neighbors

Random encounter

- ❑ People move around randomly and connect to people they bump into
 - ❑ use the num-neighbors slider, and for comparison, switch PROB-OR-NUM to 'off' to have the ER model aim for num-neighbors as well
 - ❑ turn off the layout algorithm while this is running (you can apply it at the end)
-

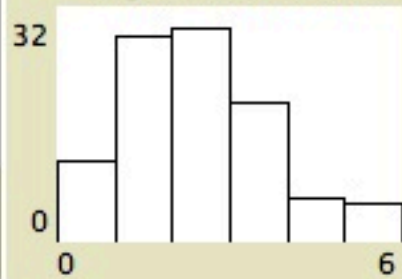
random encounters

num-nodes 99

On Off prob-or-num?

GC size	av. deg	ASP
25	2	4.4

degree distribution



Erdos-Renyi

prob-link 0.02

introduction

prob-intro 0.83

static-geo

rand-encounter

num-neighbors 2

growth

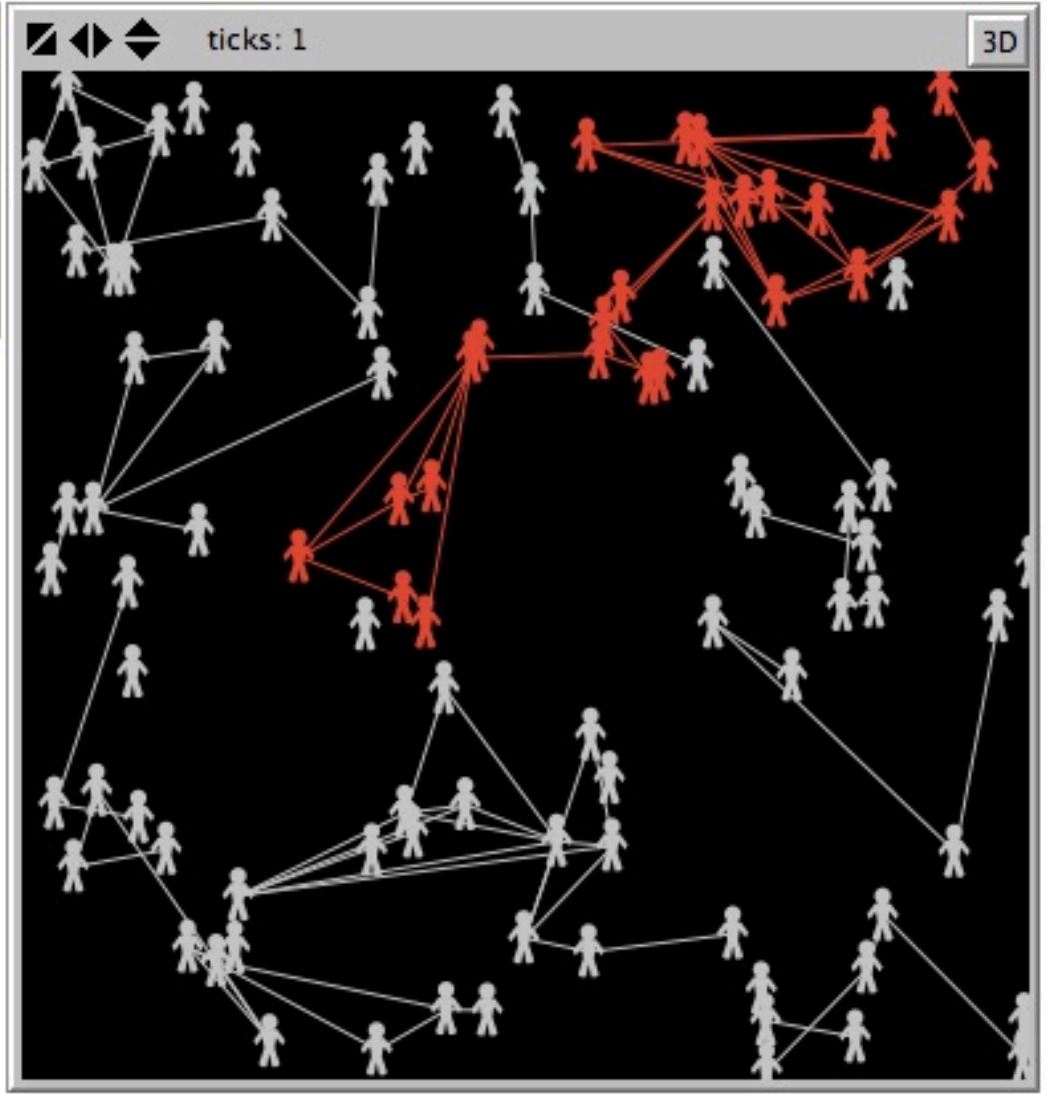
redo layout

layout options

spring-length 6

spring-constant 0.1

repulsion-strength 2.4



Quiz Q:

- Relative to ER, the random encounters model has :
 - more closed triads
 - fewer closed triads
 - smaller giant component at a low number of neighbors
 - larger giant component at a low number of neighbors
-

Growth model

- ▣ Instead of starting out with a fixed number of nodes, nodes are added over time
 - ▣ use the num-neighbors slider, and for comparison, switch PROB-OR-NUM to 'off' to have the ER model aim for num-neighbors as well
-

growth model

num-nodes 99

On
 Off prob-or-num?

GC size	av. deg	ASP
65	1.88	3.9

degree distribution

ticks: 1 3D

Erdoes-Renyi

introduction

static-geo

rand-encounter

growth

redo layout

prob-link 0.02

prob-intro 0.83

num-neighbors 2

layout options

spring-length 6

spring-constant 0.1

repulsion-strength 2.4

Quiz Q:

- Relative to ER, the growth model has :
 - more hubs
 - fewer hubs
 - smaller giant component at a low number of neighbors
 - larger giant component at a low number of neighbors

other models

- in some instances the ER model is plausible
- if dynamics are different, ER model may be a poor fit