Combinatory Categorial Grammar

Constraining surface realisation in OpenCCG

OpenCCG surface realisation



Recommended Reading

- Michael White. 2006.
 Efficient Realization of Coordinate Structures in
 Combinatory Categorial Grammar. Research on
 Language and Computation, 4(1):39–75.
- Mark Steedman and Jason Baldridge. Combinatory Categorial Grammar. To appear in Robert Borsley and Kersti Borjars (eds.) Constraint-based approaches to grammar: alternatives to transformational syntax. Oxford: Blackwell. PDF (Will appear in February 2011.)

Categorial Grammar

Categorial grammars are lexicalised grammars

- a grammar is just a "dictionary"
- there are no language-specific grammar rules
- a grammar is a **mapping** from words to structures



Lexicalised grammars

Many kinds of lexicalised grammar

- Categorial grammars (including CCGs)
- Lexicalised Tree Adjoining Grammars (LTAGs)
- CFGs in Greibach Normal Form

Lexicalised grammars are more efficient than arbitrary CFGs for NLG

• search space is simpler (Koller & Striegnitz, 2002)

Categories

Two kinds of category

- "atomic" categories
- "complex" categories

Categorial grammars (CGs)

A CG is a mapping from words to categories

· i.e. a set of word-category pairs

What do categories look like?



Atomic categories

Each CG is built around a finite set of atomic categories

- simple, non-composite, atomic symbols
- · similar to the symbols of a CFG

Examples:

- S sentence/clause
- NP noun phrase
- N noun
- PP preposition phrase

Atomic categories in XML

Use ${\tt atomcat}$ elements with a ${\tt type}$ attribute

```
<atomcat type="S"/>
```

```
<atomcat type="NP"/>
```

Arithmetic notation

Arithmetic notation gives us a finite set of digits

- 0, 1, 2, . . ., 9
- And a small set of operators for describing an infinite set of numbers: e.g.,
- concatenation: 23, 456, 92789
- addition: 2+7, 7+23, 456+65
- subtraction: 45 6, (2+6) (67- 34)

Complex categories

- Complex categories are built up from atomic category symbols
- From any *finite* set of atomic categories, can construct an *infinite* set of complex categories using two operators
 - directional slash operators: / and $\$

Traditional arithmetic notation is a useful analogy

Recursive definition

Categories are defined recursively

Atomic categories constitute the "base"

· every atomic category is also a category

The recursion involves the slash operators

- if X and Y are both categories, then so is (X/Y)
- if X and Y are both categories, then so is (X Y)

Simple examples

category	meaning
(S\NP)	verb phrase, intransitive verb
(NP/N)	determiner
(N\N)	noun post-modifier, relative clause
(PP/NP)	preposition
(PP\NP)	postposition

Embedded examples

category	meaning
((S\NP)/NP)	transitive verb
((S\NP)/NP)/NP	ditransitive verb
((N\N)/NP)	post-nominal preposition
((S\NP)\(S\NP))	adverb
((S\NP)\((S\NP)/NP))	reflexive pronoun
((N\N)/(S\NP))	relative pronoun

Notational conveniences

Drop outermost parentheses

- (S\NP) \Rightarrow S\NP
- $((N\backslash N)/(S\backslash NP)) \implies (N\backslash N)/(S\backslash NP)$

Assume left associativity of / and $\$

- ((S\NP)/NP)/NP \Rightarrow S\NP/NP/NP
- $(N\N)/(S\NP) \Rightarrow N\N/(S\NP)$

Complex categories in XML

How to represent S\NP:

```
<complexcat>
<atomcat type="S"/>
<slash dir="\"/>
<atomcat type="NP"/>
</complexcat>
```

S\NP/NP in XML



NN/(SNP) in XML







Our lexicon

Giovanni's :- NP pasta :- NP

serves :- S\NP/NP

rocks :- S\NP

restaurant :- N

great :- N/N

a : NP/N

that :- N\N/(S\NP)

Remember HLDS?

- The input to the OpenCCG realiser is a hybrid logic dependency structure
- So our categorial lexicon needs to include HLDS in some way
- We need to be able to relate the grammatical sentences with their HLDS *(interpretation)*
- And also to relate HLDSs to the grammatical sentences that can realise them (generation)

What does our grammar do?

- It tells us which strings of words are grammatical and which are not.
- It assigns derivational structure to the grammatical strings.
- · But what about semantics?

Adding HLDS to our lexicon

Two steps:

- 1. Add a nominal to each atomic category symbol
- 2. Add a set of elementary predications of hybrid logic to each lexical category
- Then relax and let forward and backward application (i.e. unification) take care of the rest!

Our lexicon again

Giovanni's :- NP pasta :- NP serves :- S\NP/NP rocks :- S\NP restaurant :- N great :- N/N a : NP/N that :- N\N/(S\NP)

Adding nominals in XML

1. Adding nominals to categories

Giovanni's :- NP_x pasta :- NP_x serves :- S_e\NP_x/NP_y rocks :- S_e\NP_x restaurant :- N_x great :- N_x/N_x a : NP_x/N_x that :- N_x\N_x/(S_e\NP_x)

 $N_{x} \setminus N_{x}$

Subscripts to atomic category symbols

- Referential indices: unique labels for object or event evoked by the word
- By convention, use x, y, z for *objects*, and e, f, g for *events*

• Coindexed nominals indicate the referent of the argument is the same as referent of result, e.g., "great"

Nominal coindexation in XML





So where are we?

- · We' ve seen how to define a lexicon in CG
- We' ve learned about two important operators in CG, i.e., forward and backward application
- · We've seen how to combine words both
 - Syntactically (derivations, unification), and
 - Semantically (set union of elementary predications)
- But, Combinatory Categorial Grammar gives us much more

From CG to CCG

CCG is an "extension" of CG.

CCG has more rules:

- forward and backward type raising
- forward and backward composition

Everything else remains the same -

• in particular the HLDS representations.

Forward type raising



Type Raising

- CCG includes type-raising rules, which turn arguments into functions over functions over such arguments
- Forward type raising

$$\xrightarrow{X} T$$

• Example:



 The rules are order preserving. Here we turn an NP into a rightward looking function over leftward functions, preserving the linear order of constituents



Further Reading

- Jason Baldridge and Geert-Jan Kruijff. 2003. "Multi-Modal Combinatory Categorial Grammar". In *Proceedings of EACL 2003*.
- Mike White and Jason Baldridge. 2003. "Adapting Chart Realization to CCG". In *Proceedings of ENLG 2003*.
- Jason Baldridge and Geert-Jan Kruijff. 2002. "Coupling CCG with Hybrid Logic Dependency Semantics". In *Proceedings of ACL 2002*.