



Loosely-Coupled/Tightly-Constrained (LCTC) Workflows for Inter-organizational, Compliance-Governed, Collaborative Computing for Multi-Tenant Cloud

Marc Edwards (SCI), Rhett Davis (NCSU), Paul Rad (UTSA)

Taxonomy of Loosely Couple Tightly Constrained Workflows

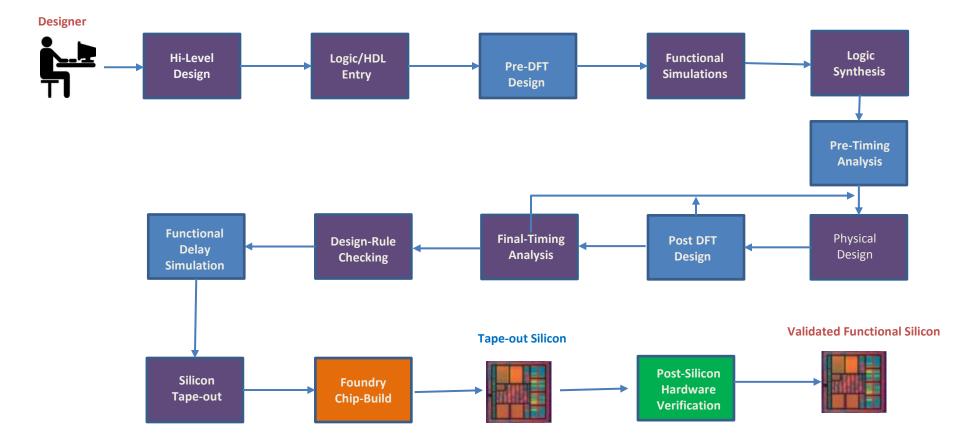
Workflow

https://en.wikipedia.org/wiki/Workflow

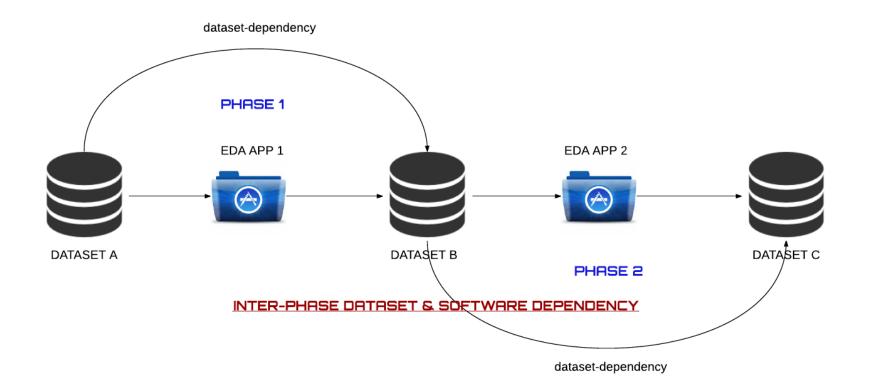
From Wikipedia, the free encyclopedia

A **workflow** consists of an orchestrated and repeatable pattern of business activity enabled by the systematic organization of resources into processes that transform materials, provide services, or process information.^[1] It can be depicted as a sequence of operations, declared as work of a person or group,^[2] an organization of staff, or one or more simple or complex mechanisms.

- We seek a computational framework for (1) IP DRM, (2) COMPUTATIONAL COMPLIANCE, (3) SECURE COLLABORATION (IP SHARING)
- Engineering Workflows (EW) dominate engineering design practices, e.g. semiconductor, FEA, engine analysis, tire design are well defined EWs.
- □ We distinguish *Engineering Workflows* from *Scientific Workflows (SW)*
 - SWs are "tightly-coupled" steps, EWs are "loosely-coupled" "phases"
- Engineering Workflows are classified as (1) Design To Release Manufacturing (DTRM) workflows, i.e. the computing activities that lead UP TO the point of releasing the product in contrast to (2) Manufacturing To Deployment (MTD) workflows.
- □ *"sequence of operations"* → EW *"sequences"* can be described as *PHASES* of a workflow.
- □ EACH **PHASE** of an workflow represents a *computational transformation* for an INPUT and OUTPUT dataset(s), i.e. hyper-complex function.
 - LINKED COLLECTION of PHASES constitute the EW



Workflow Inter-phase Dependencies



Silicon Cloud Overview

EW *Intellectual Property* (IP) is represented as INPUT & OUTPUT DATASETS

- We seek to track the *"transformational provenance"* of an IP's progression through the phases of a workflow, i.e. the *«execution provenance»* of the IP
- *"Execution Provenance"* of IP exemplifies the *DRM* component of the IP.

□ A key and distinguishing characteristic of DTRM EWs from SWs...

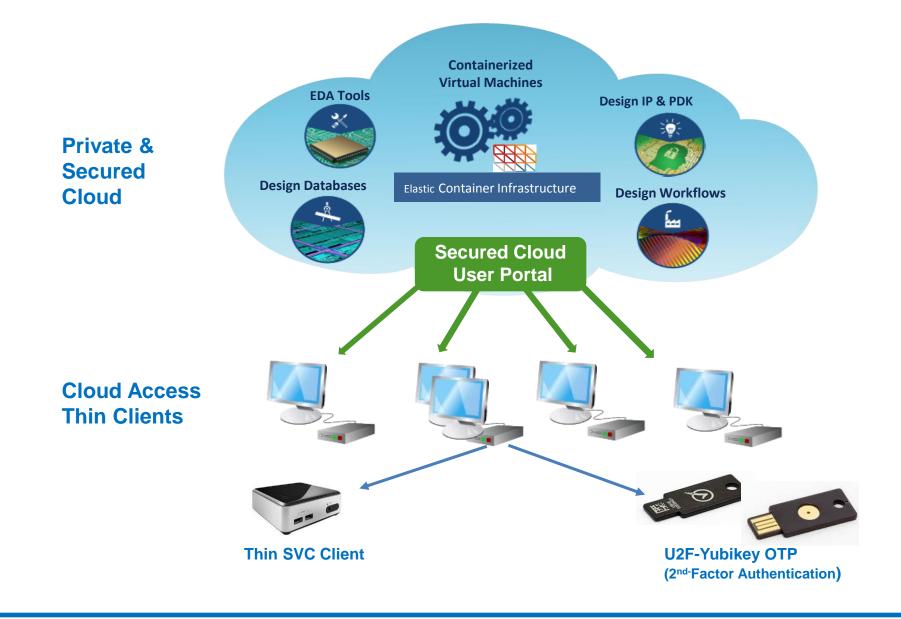
- The INTER-dependencies ("linked") of the *"transformational provenance"* of IP BETWEEN PHASES that are LOOSELY COUPLED, but...
 - "Transformational Provenance" TIGHTLY CONSTRAINS the IP within the DTRM workflow

 "TIGHTLY CONSTRAINED" component DEFINES the COMPLIANCE aspect of an Loosely-Coupled-Tightly-Constrained (LCTC) workflow

□ The multi-phase aspect of the workflow enables a *IP PHASE EXTRACTION/INSERTION POINT*

- Exposes a SECURITY HOLE within the workflow
- IP datasets can be "ejected" from the workflow
- CON → Destroys the IP security within the workflow
- PRO → Provide a mechanism for DRM through provenance record
- EW execution often requires multi-person/roles
- Demands a secure collaboration schema as EW phases are assigned through roles

Silicon Cloud I2 Infrastructure Overview



- Workflow-engaged inter-organizational collaboration through Role-Based-Access-Control (RBAC)
 - *Digital Rights Management* (DRM) scheme is enforced through a "Vaulted Cloud" graph model wherein digital IP is "execution provenanced" in real time in a graph database.
 - Graph database of multi-propertied nodes and edges
 - Resulting graph is mined for pattern matches that workflow and DRM rules of behavior
 - Enforcement, security, monitoring, compliance
 - VMs are "execution provenanced"
 - VMs do not exist as isolated computation engines, but have active database management during execution
 - Projects own VMs as *"computational objects"* that have lifecycled (cradle-2-grave) through a directed, acyclic graphoriented workflow.
 - The *Silicon Cloud* encapsulates a *WorkFlow-as-a-Service* (WFaaS) and tightly-coupled *Digital-Rights-Management-as-a-Service* (DRMaaS) architecture.
- Ephemeral inter-organizational collaboration through SDN
 - Projects are data isolated at VLAN, TCP/IP, and VPN layers
 - Inter-organizational collaboration is supported through a provenance VM with data-isolated, computational object (VM)





Graph Data Model For Workflow Compliance & IP Secure Collaboration (1)

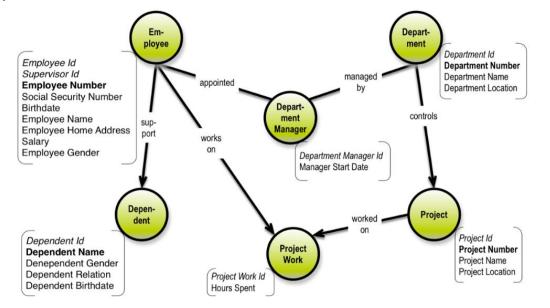
In computing, a graph database is a database that uses graph structures for semantic queries with nodes, edges and properties to represent and store data. A key concept of the system is the graph (or edge or relationship), which directly relates data items in the store. The relationships allow data in the store to be linked together directly, and in many cases retrieved with one operation.

Graph databases, by design, allow simple and fast retrieval^[citation needed] of complex hierarchical structures that are difficult to model^[according to whom?] in relational systems. Graph databases are similar to 1970s network model databases in that both represent general graphs, but network-model databases operate at a lower level of abstraction^[1] and lack easy traversal over a chain of edges.^[2]

Workflows resemble Directed Acyclic Graphs (DAGs), which are by name, GRAPHS, with associated NODES and EDGES

Assigning *«property-attributes»* to NODES & EDGES results in a *"property graph"*

- LCTC workflows are easily modeled as *"property graphs"*
 - NODES → Workflow PHASE with *«node-properties»* as app used in that phase
 - EDGES → IP transformation relationships between phases with «edge-properties» as dependency metadata



Graph Data Model For Workflow Compliance & IP Secure Collaboration (2)

DEFINE the workflow-reference-graph (WRG) as the BASELINE for...
 Workflow Compliance Patterns (WCP)

«execution provenance» is transactionally recorded as a node-edge relationships in a graph database (*Titan* and next *DSE-Graph*)

□ PROVENANCE TRAIL is created that TRACKS IP transformations across multi-phase workflow

Repeated cycling of PHASES of the LCTC workflow can be defined as the "annealing" process of the workflow, i.e. the design is maturing...
 (1) WITHIN a phase and (2) ACROSS ALL phases of the workflow

Key element visibly apparent within the reference workflow DAG is the opportunity for IP PHASE EXTRACTION & INSERTION

□ Encapsulating the entire workflow execution is a fundamental premise for cloud-based...

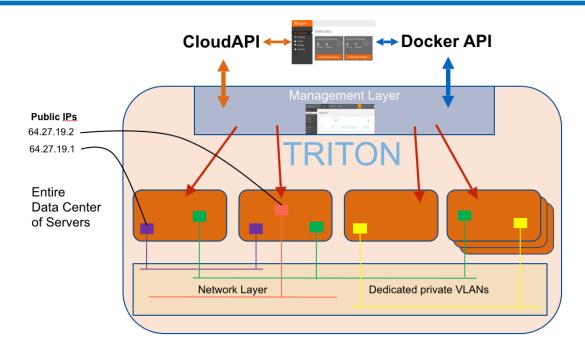
(1) workflow compliance, (2) IP collaboration, & (3) IP security

Use Tinkerpop's *Gremlin Query Language* to search for relevant WCPs

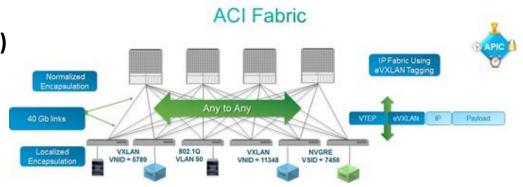
Defining Characteristics

- □ Engineering workflow applications are...
 - (1) Performance-driven, (2) large memory footprint, (3) I/O intensive
- □ Infrastructure is embodied through bare-metal server farms
- Compute workloads are (1) dominantly batch-oriented & (2) interactive graphics (2D & 3D)
- Primarily SINGLE-threaded apps...MULTI-threaded apps are increasingly important
 - Multi-threaded apps are more difficult to optimize performance across multi-core processors and INTER-server in generically networked server farms
- Linux desktops, e.g. Gnome, KDE, running an application's X-Windows GUI
- □ NFS file shares for inter-PHASE data sharing...Object Store?? ← What is that?
- Data security is the FIREWALL appliance...INTERNAL data security (RBAC)? What is that?
 - CHALLENGES FOR TRADITIONAL WORKFLOWS IN GENERALIZED PUBLIC CLOUDS
- □ Transition of BATCH compute loads
 - FIXED bare-metal server farm with LSF → DYNAMICALLY provisioned & DYNAMICALLYpowered (vCPUs) Linux containers
- □ Move away from VNC desktop to JavaScript HTML-rendered <canvas> bitmap

Leveraging Triton Cloud Middleware & Cisco ACI Fabric Infrastructure For Security



- Container virtualization and orchestration through open-source Triton/SmartOS Cloud Middleware
- Compute nodes → Cisco C240 M4 (20C) Servers
- L2/L3/L4 Switching Fabric →
- Cisco ACI (Nexus 9K)
 - 40G fabric backplane



TOWARDS HIGH PERFORMANCE VIRTUALIZATION

- Applications are virtualized through "application" & "multi-process" Linux Containers (Docker API)
- □ Containers as MICRO-serviced, MICRO-networked, MICRO-clustered "compute-shire"
 - Provide a deterministic, flexible, socket-based micro-cluster for MULTI-threaded apps
 - 20x compute density improvement over fully virtualized OS VMs
- □ From bare metal server to NON-hypervisor, OS-Virtualization with Linux Containers
- □ FIXED bare-metal server farm with LSF → DYNAMICALLY provisioned & DYNAMICALLYpowered (vCPUs) Linux containers
 - 7 sec container provisioning
- □ Move from process-isolation to OS-user space isolation
- □ No virtio, containers have direct access to kernel-mode device drivers & adapters, e.g. GPUs
- □ Applications are virtualized through Linux Containers (Docker API)
- □ Containers as MICRO-serviced, MICRO-networked, MICRO-clustered *compute-shire*
 - Provide a deterministic, flexible, socket-based micro-cluster for MULTI-threaded apps

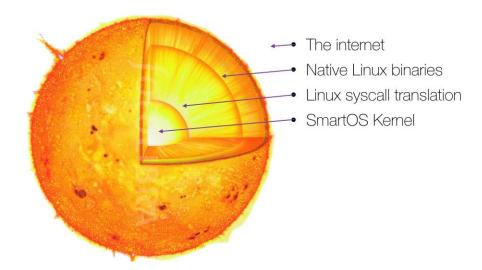
TOWARDS SOFTWARE-DEFINED-DATA-CENTER (SDDC)

□ From fixed, DHCP-based networking → Software-Defined-Network (SDN) (Cisco ACI)

- VM-2-VM micro-networked, contract-oriented TCP flows
- General Software Defined Data Center (SDDC) VXLAN, project-network isolation
- □ Integrate the container-orchestration with Cisco's leaf-spine ACI network
- ACI provides a RESTful, single point of control that is integrated with the container provisioning,
- □ "Project SDDCs" exist within a "Tenant-of-Tenants" (ToT) ACI architecture, i.e. project-tenant
- □ "Project Access Containers" (PACs) or are assigned as "computing identity" for EACH user
- PACs can only "project connect" to ONE project-tenant at a time and PACs are "morphed" with a project-owned PAC-filesystem that prevents removal of project-owned IP and crossproject IP contamination.
- Project IP sharing is facilitated through object-storage with ACI "contracts" between the object-store application and a project-network/VLAN.
- Node-Edge transactions are generated and recorded in graph database and Process or IP Compliance Violation Policies (CVP) are monitored against all IP Object Stores

TOWARDS SOFTWARE-DEFINED-DATA-CENTER (SDDC)

- □ FIXED bare-metal server farm with LSF → DYNAMICALLY provisioned & DYNAMICALLYpowered (vCPUs) Linux containers
- □ Applications are virtualized through Linux Containers (Docker API)
- □ Containers as MICRO-serviced, MICRO-networked, MICRO-clustered compute-shire
 - Provide a deterministic, flexible, socket-based micro-cluster for MULTI-threaded apps

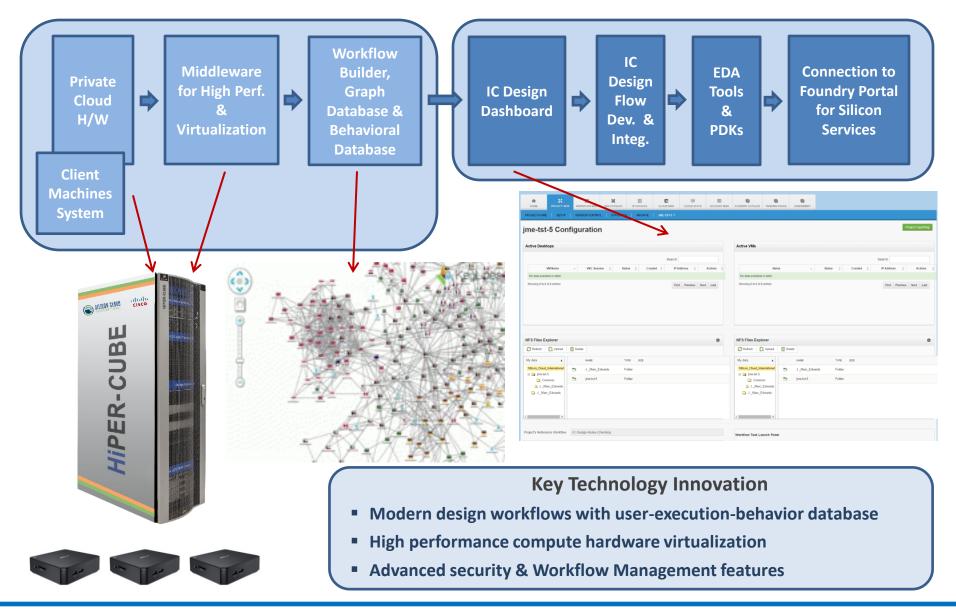


OBJECT STORAGE WITH COMPUTE

□ CHALLENGES FOR TRADITIONAL WORKFLOWS IN GENERALIZED PUBLIC CLOUDS

- □ Transition of BATCH compute loads
 - FIXED bare-metal server farm with LSF → DYNAMICALLY provisioned & DYNAMICALLYpowered (vCPUs) Linux containers
- □ Move from NAS to object-store-with-compute
- **CLI** phase/pipe processing for mapping & reducing *Data At Rest*

Summary of the SCI Cloud Technology Innovations





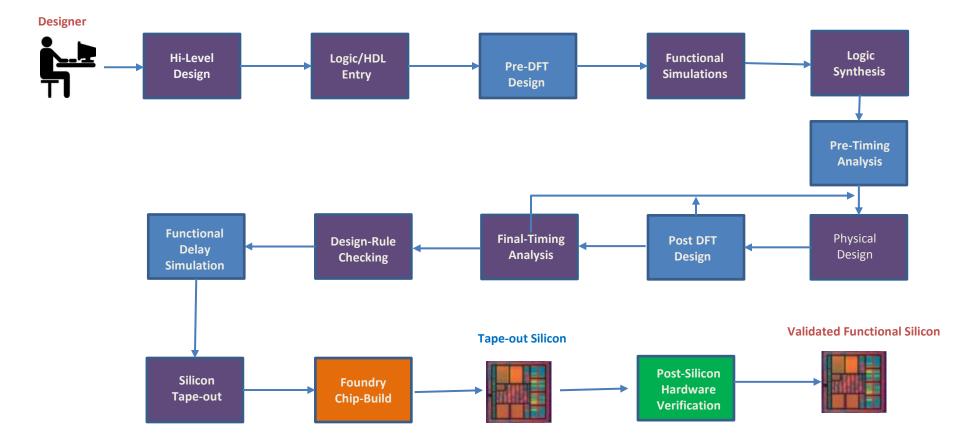


Workflow Compliance Example-1

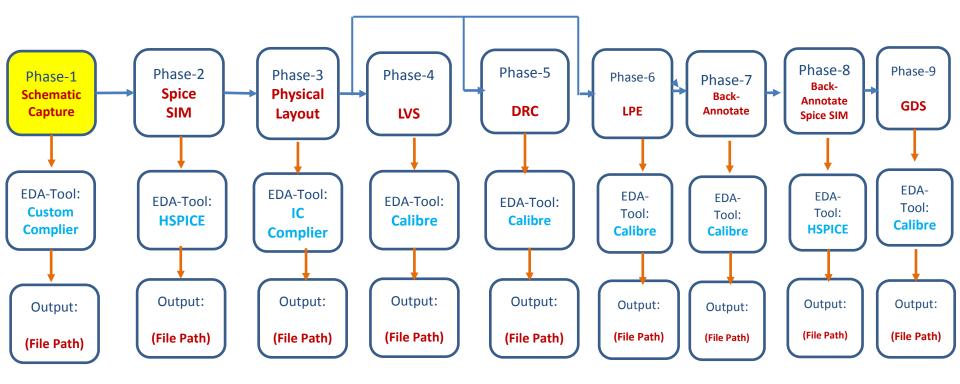
NCSU Semiconductor design Workflow Application

Professor W. Rhett Davis (NCSU)

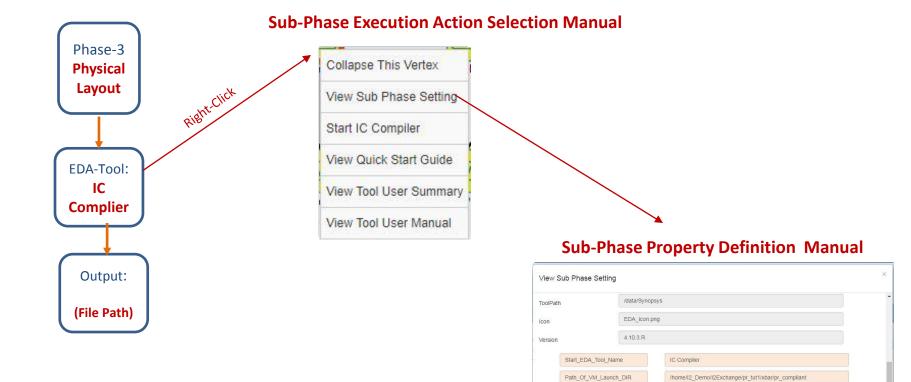
W. Rhett Davis consults for Silicon Cloud International (SCI) on the development of Semiconductor Workflows. He also conducts research in areas of interest similar to the business interest of SCI. The terms of this arrangement have been reviewed and approved by NC State University in accordance with its policy on objectivity in research.



Overall X-Bar Design Workflow Graph



Sub-Phase Action & Property Setting Manuals



View_Quick_Start_Guide

Eda_Tool_User_Summary

Eda_Tool_User_Manual

Compliance Policy_Action

Compliance_Policy (yes or no)

wf/docs/jazz-ca18hd_quickstart.pdf

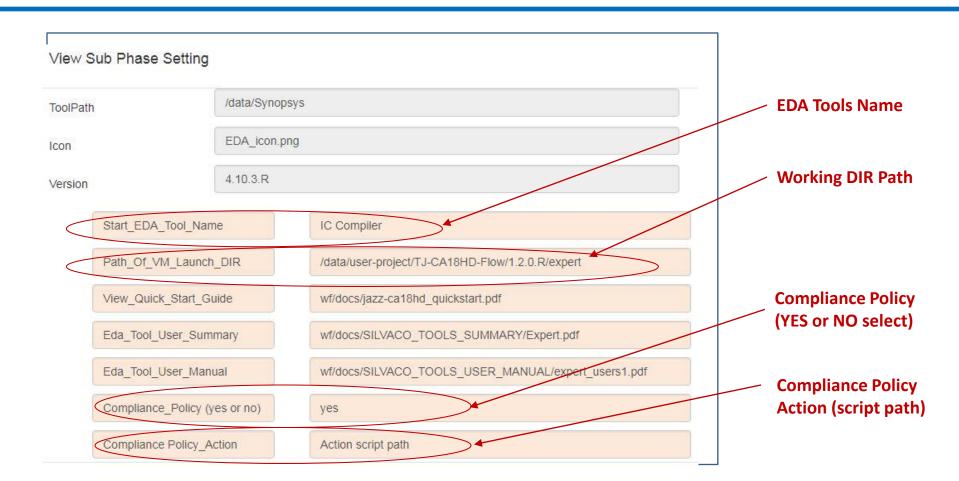
yes

wf/docs/SILVACO_TOOLS_SUMMARY/Expert.pdf

wf/docs/SILVACO_TOOLS_USER_MANUAL/expert_users1.pdf

/home/l2_Demo/l2Exchange/pr_tut1/xbar/pr_compliant/compliance_cf

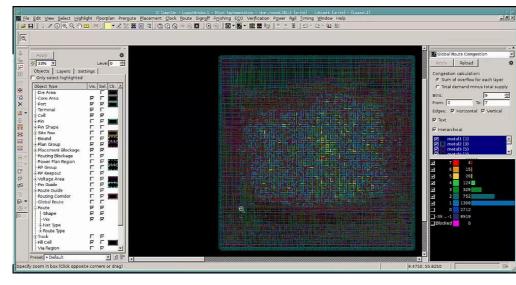
Sub-Phase Property Setting Manual (details)



Design Process Compliance Check Example:

□ To show a <u>non-compliant</u> design process case

- Design cell layout placement and routing in conflict result
- □ To show a <u>compliant</u> design process case
 - Design cell layout placement and routing in non-conflict result



Non-Compliant Layout Case

- 8 > File Edit View Select Highlight Floorplan Preroute Placement Clock Boute Signoff Finishing ECO Venfication Power Timing Window Help |4 ≠ 0 Q Q A 🛎 🗶 🔽 🗶 🗟 3 | Q Q Q A O 🗟 | Q O] 📓 ▼ 📓 🗟 | 2 ± 3 | 2 + Q + Q S # H 0 ÷ 100% ▼ Level Objects | Layers | Settings | Only select highlighted Object Type 표 -Die Area e × -Core Area -Port Terminal 4 3 Cell Pin Pin Shape -Site Row -Bound Plan Group R Placement Blockage -Routing Blockage Power Plan Region -RP Group RP Keepou Voltage Area Pin Guide 125 Boute Guide -Routing Corrido -Global Route oute -Shape 5 Net True - Route Type Track -Fill Cell Via Region Preset + Default 128 pecify zoom in box (Click opposite corners or drag -39.0650.93.2650

Compliant Layout Case

HSPICE Simulation Model Standard Compliance Check Example:

U To show a <u>non-compliant</u> simulation model case

- Result device leakage current 2% above standard specification
- **U** To show a <u>compliant</u> simulation model case
 - Result device leakage current within the standard specification

Non-Compliant Result Case

Compliant Result Case

	Terminal	_ D X	💷 Terminal _ n x
File Edit View S	Search Terminal Help		File Edit View Search Terminal Help
	549102e-08 -4.214924e-08 0.000000e+00 -1.006883e-07 000000e+00 2.107452e-08 0.000000e+00 5.034414e-08	×.	current 4.488824e-08 -3.831359e-08 0.000000e+00 -9.909708e-08 power 0.000000e+00 1.915679e-08 0.000000e+00 4.954854e-08
volts 0.0 current 1.8	0:V9Men 0:V5 0000000+00 0.000000+00 8865556-08 7.7649932-08 000000+00 0.000000+00		subckt element 0.0000000e+00 0.00000e+00 volts 0.000000e+00 0.00000e+00 <u>current</u> 1.869709e-08 7.392534e-08 power 0.000000e+00 0.000000e+00
total vol	ltage source power dissipation= 7.141876e-08 watt	5	total voltage source power dissipation# 6.870533e-08 watts
1****** HSPICE	** job concluded E K-2015.06-2 linux64 (Jun 25 2015) ****** ation for public bsim-soi models		***** job concluded f 1****** HSPICE K-2015.06-2 linux64 (Jun 25 2015) ****** * test simulation for public bsim-soi models
· test simula	ación for public bsim-sor models		test simulation for public bain-sol models
****** job st	tatistics summary tnom= 25.000 temp= 0.000 ******		****** job statistics summary tnom= 25.000 temp= 0.000 ******
****** Machi	ine Information ******		****** Machine Information ******
CPU: model name cpu MHz	: Intel(R) Core(TM) i7-4790 CPU @ 3.60GHz : 3591.601	+	CPU: model name : Intel(R) Core(TM) 17-4790 CPU @ 3.60GHz cpu MHz : 3591.601
	n 2.6.32-696.3.1.el6.x86_64 (mockbuild@x86-042.build. rsion 4.4.7 20120313 (Red Hat 4.4.7-18) (GCC)) #1 SM 017		05: Linux version 2.6.32-696.3.1.el6.x86_64 (mockbuild@x86-042.build.eng.bos.redhat .com) (gcc version 4.4.7 20120313 (Red Hat 4.4.7-18) (GCC)) #1_SMP Thu Apr 20 1 1:30:02 EDT 2017
***** HSP]	ICE Threads Information ******		****** HSPICE Threads Information ******
Command Line Available CF Actual Three			Command Line Threads Count : 1 Available CPU Count : 8 Actual Threads Count : 1
<pre># nodes # resistors # mutual_ind # cccs # curr_srcs</pre>	ds = 0 # vccs = 0 # vccs = = 0 # ccvs = 0 # volt_srcs = = 0 # diodes = 0 # bjts = = 0 # mosfets = 0 # b elements = 5 = 0 # W elements = 0 # B elements =	0 6 0 0 5	# nodes 30 # elements 11 # resistors 0 # capacitors 0 # inductors 0 # mexistors 0 # capacitors 0 # inductors 0 # mutual_inds 0 # vccs 0 # vccs 0 # cccs 0 # vccs 0 # volt_srcs 6 # cccs 0 # volt_srcs 6 # ccrr srcs 0 # diodes 0 # blts 0 # jfots 0 # mosfets 0 # U elements 0 # clements 0 # W elements 0 # w device 5



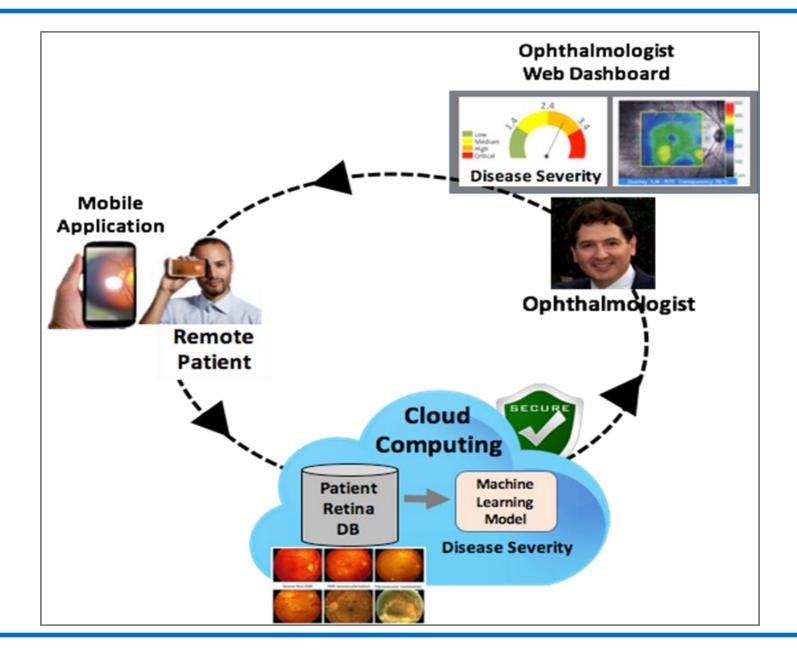


Workflow Compliance Example-2

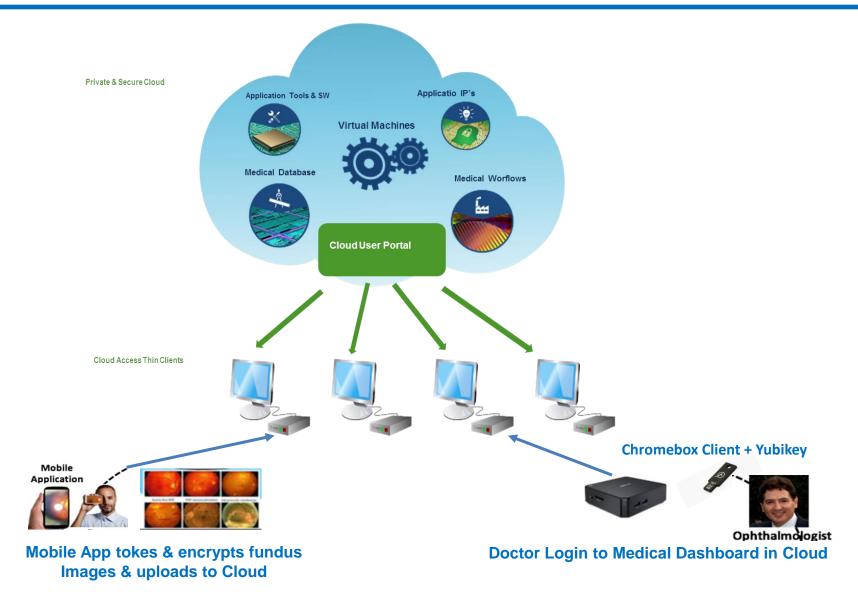
UTSA Cloud-Based Secure Mobile Teleophthalmology Healthcare Application

Professor Paul Rad (UTSA)

Cloud-Based Secure Mobile Teleophthalmology Workflow



Silicon Cloud Integration of the Teleophthalmology Workflow



Step 1: Taking Picture from Patient's Fundus in different locations

Step 2: The mobile-cloud App will encrypt the Taken Fundus Image and will send the encrypted image to the secure multitenant cloud storage

Step 3: The Machine Learning Algorithm that is designed to identify patient disease severity level runs on all collected fundus images stored in the secure multitenant cloud storage.

Step 4: System will classifies patients based on their disease severity level and results will be send to the Ophthalmologist dashboard.

Step 5: When Ophthalmologist login to his cloud dashboard, list of patients and their disease severities as well as their fundus image is presented.

Step 6: Ophthalmologist can follow up with patients on follow ups and diagnoses

