The Decidability of the First-order Theory of Knuth-Bendix Order

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How to decide satisfi ability of order constraints?



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QFT: Quantifi er-free Theory.

UQT: Unary Quantifi ed Theory.



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QFT: Quantifi er-free Theory.

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QFT: Quantifi er-free Theory.

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QFT: Quantifi er-free Theory.

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Our approach: $Th(KBO) \rightarrow Th(PA)$



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Our approach: $Th(KBO) \rightarrow Th(PA)$

Reduce term constraints to integer constraints. [ZSM04a]



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Our approach: $Th(KBO) \rightarrow Th(PA)$

Reduce term constraints to integer constraints. [ZSM04a]

Reduce term quantifi ers to integer quantifi ers. ZSM04b]



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Our approach: $Th(KBO) \rightarrow Th(PA)$

- Reduce term constraints to integer constraints. [ZSM04a]
- Reduce term quantifi ers to integer quantifi ers. ZSM04b]

Integers rule!



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A term algebra \mathfrak{A}_{TA} : $\langle TA; C, \mathcal{A}, \mathcal{S}, \mathcal{T} \rangle$ consists of



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A term algebra \mathfrak{A}_{TA} : $\langle TA; C, \mathcal{A}, \mathcal{S}, \mathcal{T} \rangle$ consists of

1. TA: The term domain.



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A term algebra \mathfrak{A}_{TA} : $\langle TA; C, \mathcal{A}, \mathcal{S}, \mathcal{T} \rangle$ consists of

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2. *C*: A finite set of constructors: α , β , γ ,



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4. S: A finite set of selectors. $\alpha = (s_1^{\alpha}, \dots, s_k^{\alpha})$.



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4. S: A finite set of selectors. $\alpha = (s_1^{\alpha}, \dots, s_k^{\alpha})$.

5. \mathcal{T} : A finite set of testers. Is_{*a*} for $\alpha \in C$.



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rightarrow TA is generated exclusively using C.



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4. S: A finite set of selectors. $\alpha = (s_1^{\alpha}, \dots, s_k^{\alpha})$.

- 5. \mathcal{T} : A finite set of testers. Is_{*a*} for $\alpha \in C$.
- \sim TA is generated **exclusively** using C.
- Each element of TA is uniquely generated.



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 \langle list; {cons, nil}; {nil}; {car, cdr}; {ls_{nil}, ls_{cons}} \rangle



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 $\langle list; \{cons, nil\}; \{nil\}; \{car, cdr\}; \{ls_{nil}, ls_{cons}\} \rangle$

Axioms:

■ Signature:

$$\begin{split} \mathsf{Is}_{\mathsf{nil}}(x) &\leftrightarrow \neg \mathsf{Is}_{\mathsf{cons}}(x), \\ x &= \mathsf{car}(\mathsf{cons}(x, y)), \\ y &= \mathsf{cdr}(\mathsf{cons}(x, y)), \\ \mathsf{Is}_{\mathsf{nil}}(x) &\leftrightarrow \{\mathsf{car}, \mathsf{cdr}\}^+(x) = x, \\ \mathsf{Is}_{\mathsf{cons}}(x) &\leftrightarrow \mathsf{cons}(\mathsf{car}(x), \mathsf{cdr}(x)) = x. \end{split}$$



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We study KBO using selector language.



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We study KBO using selector language.

• For $L = s_1, \ldots, s_n$, Lx stands for

 $S_1(\ldots(S_n(x)\ldots)).$

|L| is called the depth of x in Lx.



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• depth\varphi(x) : the maximum depth of x in \varphi.
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• For $L = s_1, \ldots, s_n$, Lx stands for

 $S_1(\ldots(S_n(x)\ldots)).$

|L| is called the depth of x in Lx.

depth $\varphi(x)$: the maximum depth of x in φ .

Formulas are type-complete and selector terms are proper.
 For example,

 $car(x) \neq cdr(x)$

should be understood as

 $car(x) \neq cdr(x) \land Is_{cons}(x).$



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A Knuth-Bendix order (KBO) <^{kb} is parametrically defined with



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A Knuth-Bendix order (KBO) <^{kb} is parametrically defined with

• $W: TA \rightarrow \mathbb{N}: a \text{ weight function satisfying}$

$$W(\alpha(t_1,\ldots,t_k)) = W(\alpha) + \sum_{i=1}^k W(t_i).$$



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A Knuth-Bendix order (KBO) <^{kb} is parametrically defined with

• W : TA $\rightarrow \mathbb{N}$: a weight function satisfying

$$W(\alpha(t_1,\ldots,t_k)) = W(\alpha) + \sum_{i=1}^k W(t_i).$$

• $<^{\Sigma}$: a linear (precedence) order on *C* such that

$$\alpha_1 >^{\Sigma} \alpha_2 >^{\Sigma} \ldots >^{\Sigma} \alpha_{|C|}$$



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For $u, v \in TA$, $u <^{kb} v$ if one of the following holds:



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For $u, v \in TA$, $u <^{kb} v$ if one of the following holds:

• W(u) < W(v).



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For $u, v \in TA$, $u <^{kb} v$ if one of the following holds:

$$\mathbf{W}(u) < \mathbf{W}(v).$$

• W(u) = W(v) and type $(u) <^{\Sigma} type(v)$.



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Techinical Catches

For $u, v \in TA$, $u <^{kb} v$ if one of the following holds:

•
$$W(u) < W(v)$$
.

•
$$W(u) = W(v)$$
 and $type(u) <^{\Sigma} type(v)$.

•
$$W(u) = W(v), u \equiv \alpha(u_1, \dots, u_k), v \equiv \alpha(v_1, \dots, v_k)$$
, and

$$\exists i [1 \le i \le k \land u_i <^{\mathsf{kb}} v_i \land \forall j (1 \le j < i \to u_j = v_j)].$$



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■ Suffi ces to eliminate ∃-quantifi ers fromprimitive formulas

 $\exists \bar{x}(A_1(\bar{x}) \land \ldots \land A_n(\bar{x})),$

where $A_i(\bar{x})$ ($1 \le i \le n$) are literals.



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Suffices to eliminate \exists -quantifiers from primitive formulas $\exists \bar{x}(A_1(\bar{x}) \land \ldots \land A_n(\bar{x})),$

where $A_i(\bar{x})$ $(1 \le i \le n)$ are literals.

Suffices to assume $A_i \neq x = t$ if $x \notin t$, because

 $\exists x(x=t \land \varphi(x, \bar{y})) \leftrightarrow \varphi(t, \bar{y}).$



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Eliminating $\exists x \text{ from } (\exists x) \varphi(x, \bar{y}) \text{ is straightforward once}$

 $\varphi(x, \bar{y})$ is solved in *x*.



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Eliminating $\exists x \text{ from } (\exists x) \varphi(x, \bar{y}) \text{ is straightforward once}$

 $\varphi(x, \bar{y})$ is solved in x.

Depth Reduction.

Solved Form.

Depth reduction is to obtain solve forms as

 $\varphi(x, \bar{y})$ is solved in x iff $depth_{\varphi}(x) = 0$.



Solved Form

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• $\varphi(x, \bar{y})$ is solved in x if it is in the form

 $\bigwedge_{i\leq m} u_i \prec^{\mathsf{kb}} x \land \bigwedge_{j\leq n} x \prec^{\mathsf{kb}} v_j \land \varphi'(\bar{y}),$

where x does not appear in u_i , v_i and φ' .



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• $\varphi(x, \bar{y})$ is solved in x if it is in the form

$$\bigwedge_{\leq m} u_i \prec^{\mathsf{kb}} x \land \bigwedge_{j \leq n} x \prec^{\mathsf{kb}} v_j \land \varphi'(\bar{y}),$$

where x does not appear in u_i , v_i and φ' .

1

If $\varphi(x, \bar{y})$ is solved in *x*, then $(\exists x) \varphi(x, \bar{y})$ simplifies to

$$\bigwedge_{\leq m,j\leq n} u_i \prec_2^{\mathsf{kb}} v_j \land \varphi'(\bar{\boldsymbol{y}})$$

where $x \prec_n^{kb} y$, called gap order, states there is an increasing chain from x to y of length at least n.



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Case 1: All occurrences of *x* have depth greater than 0.



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Case 1: All occurrences of *x* have depth greater than 0. In this case, $\exists x \varphi(x, \bar{y})$ goes to $\exists x_1, \dots, \exists x_k \varphi'(x_1, \dots, x_k, \bar{y}),$

where

 $\varphi'(x_1,\ldots,x_k,\bar{y})\equiv\varphi(x,\bar{y})[x_i\leftarrow S_i^{\alpha}(x)].$



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Case 2: Some *x* have depth 0 and some do not.



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Case 2: Some *x* have depth 0 and some do not.

- Decompose 0-depth occurrences of x in terms of
 - $\mathbf{s}_1^{\alpha}(x),\ldots,\mathbf{s}_k^{\alpha}(x).$



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Case 2: Some *x* have depth 0 and some do not.

Decompose 0-depth occurrences of x in terms of

 $\mathsf{S}_1^{\alpha}(x),\ldots,\mathsf{S}_k^{\alpha}(x).$

This amounts to expressing $x <_n^{kb} t$ and $t <_n^{kb} x$ using $s_1^{\alpha}(x), \dots, s_k^{\alpha}(x).$



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Case 2: Some *x* have depth 0 and some do not.

Decompose 0-depth occurrences of x in terms of

 $\mathbf{S}_1^{\alpha}(x),\ldots,\mathbf{S}_k^{\alpha}(x).$

This amounts to expressing $x <_n^{kb} t$ and $t <_n^{kb} x$ using $s_1^{\alpha}(x), \dots, s_k^{\alpha}(x)$.

Then apply the reduction as in Case 1!



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Case 2: Some *x* have depth 0 and some do not.

Decompose 0-depth occurrences of x in terms of

 $\mathbf{S}_1^{\alpha}(x),\ldots,\mathbf{S}_k^{\alpha}(x).$

This amounts to expressing $x <_n^{kb} t$ and $t <_n^{kb} x$ using $s_1^{\alpha}(x), \dots, s_k^{\alpha}(x)$.

Then apply the reduction as in Case 1!

In order to do that, we need to extend the language.



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1. Decompose $<^{kb}$ into three disjoint suborders $<^{w}$, $<^{p}$ and $<^{I}$.



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1. Decompose $<^{kb}$ into three disjoint suborders $<^{w}$, $<^{p}$ and $<^{I}$.

2. Extend
$$\prec^w$$
, \prec^p and \prec^l to \prec^w_n , \prec^p_n and \prec^l_n , respectively.



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1. Decompose $<^{kb}$ into three disjoint suborders $<^{w}$, $<^{p}$ and $<^{I}$.

2. Extend $<^{w}$, $<^{p}$ and $<^{I}$ to $<^{w}_{n}$, $<^{p}_{n}$ and $<^{I}_{n}$, respectively.

3. Add Presburger arithmetic explicitly to represent weight.



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1. Decompose $<^{kb}$ into three disjoint suborders $<^{w}$, $<^{p}$ and $<^{I}$.

2. Extend \prec^w , \prec^p and \prec^l to \prec^w_n , \prec^p_n and \prec^l_n , respectively.

3. Add Presburger arithmetic explicitly to represent weight.

4. Definecounting constraints to count terms of certain weight.



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1. Decompose $<^{kb}$ into three disjoint suborders $<^{w}$, $<^{p}$ and $<^{I}$.

2. Extend \prec^w , \prec^p and \prec^l to \prec^w_n , \prec^p_n and \prec^l_n , respectively.

3. Add Presburger arithmetic explicitly to represent weight.

- 4. Definecounting constraints to count terms of certain weight.
- 5. Defineboundary functions to delineate gap orders.



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1. Decompose $<^{kb}$ into three disjoint suborders $<^{w}$, $<^{p}$ and $<^{I}$.

2. Extend \prec^w , \prec^p and \prec^l to \prec^w_n , \prec^p_n and \prec^l_n , respectively.

- 3. Add Presburger arithmetic explicitly to represent weight.
- 4. Definecounting constraints to count terms of certain weight.
- 5. Defi neboundary functions to delineate gap orders.
- 6. Extend all aforementioned notions to tuples of terms.



1. Weight Order <^w:

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$u \prec^{\mathsf{w}} v \Leftrightarrow \mathsf{W}(u) < \mathsf{W}(v).$



1. Weight Order <^w:

2.

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 $u <^{\mathsf{w}} v \Leftrightarrow \mathsf{W}(u) < \mathsf{W}(v).$

 $u <^{\mathsf{p}} v \Leftrightarrow \mathsf{W}(u) = \mathsf{W}(s) \& \mathsf{type}(u) <^{\Sigma} \mathsf{type}(v).$



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Weight Order <^w:

$$u \prec^{\mathsf{w}} v \Leftrightarrow \mathsf{W}(u) < \mathsf{W}(v).$$

2. Precedence Order <^p:

 $u <^{\mathsf{p}} v \Leftrightarrow \mathsf{W}(u) = \mathsf{W}(s) \& \mathsf{type}(u) <^{\Sigma} \mathsf{type}(v).$

3. Lexicographical Order <^I:

 $u \prec v \Leftrightarrow W(u) = W(v) \& type(u) = type(v) \& u \prec v.$



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1. Weight Order <^w:

$$u \prec^{\mathsf{w}} v \Leftrightarrow \mathsf{W}(u) < \mathsf{W}(v).$$

2. Precedence Order <^p:

 $u <^{\mathsf{p}} v \Leftrightarrow \mathsf{W}(u) = \mathsf{W}(s) \& \mathsf{type}(u) <^{\Sigma} \mathsf{type}(v).$

3. Lexicographical Order <^I:

 $u \prec v \Leftrightarrow W(u) = W(v) \& type(u) = type(v) \& u \prec v.$

Abbreviations:

$$u <^{\mathsf{pl}} v \Leftrightarrow u <^{\mathsf{p}} v \lor u <^{\mathsf{l}} v,$$
$$u <^{\mathsf{kb}} v \Leftrightarrow u <^{\mathsf{w}} v \lor u <^{\mathsf{p}} v \lor u <^{\mathsf{l}} v.$$



• Gap Order \prec_n^{kb} :

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 $u <_n^{\mathsf{kb}} v \leftrightarrow (\exists u_1, \dots, \exists u_n) \Big[u <^{\mathsf{kb}} u_1 <^{\mathsf{kb}} \dots <^{\mathsf{kb}} u_n \leq^{\mathsf{kb}} v \Big].$



• Gap Order $<_n^{kb}$:

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$$u <_n^{\mathsf{kb}} v \leftrightarrow (\exists u_1, \dots, \exists u_n) \Big[u <^{\mathsf{kb}} u_1 <^{\mathsf{kb}} \dots <^{\mathsf{kb}} u_n \leq^{\mathsf{kb}} v \Big]$$

• Weight Gap Order \prec_n^{w} :

$$u \prec^{\mathsf{w}}_{n} v \leftrightarrow u \prec^{\mathsf{kb}}_{n} v \land u \prec^{\mathsf{w}} v.$$



• Gap Order \prec_n^{kb} :

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$$u <_n^{\mathsf{kb}} v \leftrightarrow (\exists u_1, \dots, \exists u_n) \Big[u <^{\mathsf{kb}} u_1 <^{\mathsf{kb}} \dots <^{\mathsf{kb}} u_n \leq^{\mathsf{kb}} v \Big].$$

• Weight Gap Order $<_n^{w}$:

$$u <^{\mathsf{w}}_{n} v \leftrightarrow u <^{\mathsf{kb}}_{n} v \land u <^{\mathsf{w}} v.$$

• Precedence Gap Order \prec_n^p :

$$u \prec^{\mathsf{p}}_{n} v \leftrightarrow u \prec^{\mathsf{kb}}_{n} v \land u \prec^{\mathsf{p}} v.$$



• Gap Order $<_n^{kb}$:

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$$u \prec_n^{\mathsf{kb}} v \leftrightarrow (\exists u_1, \dots, \exists u_n) \Big[u \prec^{\mathsf{kb}} u_1 \prec^{\mathsf{kb}} \dots \prec^{\mathsf{kb}} u_n \preceq^{\mathsf{kb}} v \Big].$$

• Weight Gap Order $<_n^{W}$:

$$u <^{\mathsf{w}}_{n} v \leftrightarrow u <^{\mathsf{kb}}_{n} v \land u <^{\mathsf{w}} v.$$

• Precedence Gap Order $<_n^p$:

$$u \prec^{\mathsf{p}}_{n} v \leftrightarrow u \prec^{\mathsf{kb}}_{n} v \land u \prec^{\mathsf{p}} v.$$

• Lexicographical Gap Order \prec_n^{I} :

$$u \prec_n^{\mathsf{I}} v \leftrightarrow u \prec_n^{\mathsf{kb}} v \land u \prec^{\mathsf{I}} v.$$



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 $0^w, 1^w : \mathbb{N} \to TA; 0^p, 1^p : \mathbb{N}^2 \to TA$ such that

• $0^{w}(n)$: the smallest term of weight *n*.



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$0^w, 1^w : \mathbb{N} \to TA; 0^p, 1^p : \mathbb{N}^2 \to TA$ such that

- $0^{w}(n)$: the smallest term of weight *n*.
- $0^{p}(n,p)$: the smallest term of weight *n* and type α_{p} .



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- $1^{w}(n)$: the largest term of weight *n*.



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$0^w, 1^w: \mathbb{N} \to TA; 0^p, 1^p: \mathbb{N}^2 \to TA$ such that

- $0^{w}(n)$: the smallest term of weight *n*.
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- $1^{p}(n, p)$: the largest term of weight *n* and type α_{p} .



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 $0^w, 1^w: \mathbb{N} \to TA; 0^p, 1^p: \mathbb{N}^2 \to TA$ such that

- $0^{w}(n)$: the smallest term of weight n.
- $0^{p}(n,p)$: the smallest term of weight *n* and type α_{p} .
- $1^{w}(n)$: the largest term of weight *n*.
- $1^{p}(n, p)$: the largest term of weight *n* and type α_{p} .

Example of Use:

$$u \prec_5^{\mathsf{w}} v \leftrightarrow \bigvee_{n_1+n_2+n_3=5} u \prec_{n_1}^{\mathsf{pl}} \mathbf{1}_{(u^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_2}^{\mathsf{w}} \mathbf{0}_{(v^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_3}^{\mathsf{pl}} v.$$



• $CNT_n(x)$ states that

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"there are at least n + 1 distinct TA-terms of weight x."



• $CNT_n(x)$ states that

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"there are at least n + 1 distinct TA-terms of weight x."

• $CNT_0(x)$ (or Tree(x)) states that

x is a legitimate weight of a term.



• $CNT_n(x)$ states that

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"there are at least n + 1 distinct TA-terms of weight x."

• $CNT_0(x)$ (or Tree(x)) states that

x is a legitimate weight of a term.

• $CNT_n(x)$ is expressible in Presburger arithmetic.



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"there are at least n + 1 distinct TA-terms of weight x."

• $CNT_0(x)$ (or Tree(x)) states that

x is a legitimate weight of a term.

• $CNT_n(x)$ is expressible in Presburger arithmetic.

Example of Use:

• $CNT_n(x)$ states that

 $0_{(x)}^{\mathsf{w}} \prec_n^{\mathsf{pl}} 1_{(x)}^{\mathsf{w}} \leftrightarrow \mathsf{CNT}_n(x).$



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 $\begin{aligned} \mathfrak{A}_{\mathsf{k}\mathsf{b}^+}^{\mathbb{Z}} &= \langle & \mathfrak{A}_{\mathsf{T}\mathsf{A}}; \mathfrak{A}_{\mathbb{Z}}; (.)^{\mathsf{w}}; \\ & <_n^{\sharp}, \sharp \in \{\mathsf{k}\mathsf{b}, \mathsf{w}, \mathsf{p}, \mathsf{I}, \mathsf{p}\mathsf{I}\}, \end{aligned}$

Extended structure:

 $0^{*}(...), 1^{*}(...), * \in \{w, p\} >.$



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Extended structure:

$$\begin{split} \mathfrak{A}_{\mathsf{k}\mathsf{b}^+}^{\mathbb{Z}} &= \langle \quad \mathfrak{A}_{\mathsf{T}\mathsf{A}}; \mathfrak{A}_{\mathbb{Z}}; (.)^{\mathsf{w}}; \\ &\quad \langle_n^{\sharp}, \sharp \in \{\mathsf{k}\mathsf{b}, \mathsf{w}, \mathsf{p}, \mathsf{I}, \mathsf{p}\mathsf{I}\}, \\ &\quad \mathsf{0}^*(...), \mathsf{1}^*(...), * \in \{\mathsf{w}, \mathsf{p}\} \quad \rangle. \end{split}$$

• \mathfrak{A}_{TA} : Term algebras.



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$$\begin{split} \mathfrak{A}_{\mathsf{k}\mathsf{b}^+}^{\mathbb{Z}} &= \langle \quad \mathfrak{A}_{\mathsf{T}\mathsf{A}}; \mathfrak{A}_{\mathbb{Z}}; (.)^{\mathsf{w}}; \\ &\quad \langle_n^{\sharp}, \sharp \in \{\mathsf{k}\mathsf{b}, \mathsf{w}, \mathsf{p}, \mathsf{l}, \mathsf{p}\mathsf{l}\}, \\ &\quad \mathsf{0}^*(\ldots), \mathsf{1}^*(\ldots), * \in \{\mathsf{w}, \mathsf{p}\} \quad \rangle. \end{split}$$

- \mathfrak{A}_{TA} : Term algebras.
- $\mathfrak{A}_{\mathbb{Z}}$: Presburger arithmetic.



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$$\begin{split} \mathfrak{A}_{\mathsf{k}\mathsf{b}^+}^{\mathbb{Z}} &= \langle \quad \mathfrak{A}_{\mathsf{T}\mathsf{A}}; \mathfrak{A}_{\mathbb{Z}}; (.)^{\mathsf{w}}; \\ &\quad \langle_n^{\sharp}, \sharp \in \{\mathsf{k}\mathsf{b}, \mathsf{w}, \mathsf{p}, \mathsf{l}, \mathsf{p}\mathsf{l}\}, \\ &\quad \mathsf{0}^*(\ldots), \mathsf{1}^*(\ldots), * \in \{\mathsf{w}, \mathsf{p}\} \quad \rangle. \end{split}$$

- \mathfrak{A}_{TA} : Term algebras.
- $\mathfrak{A}_{\mathbb{Z}}$: Presburger arithmetic.
- (.)^w : Weight function.



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$$\begin{split} \mathfrak{A}_{\mathsf{k}\mathsf{b}^+}^{\mathbb{Z}} &= \langle \quad \mathfrak{A}_{\mathsf{T}\mathsf{A}}; \mathfrak{A}_{\mathbb{Z}}; (.)^{\mathsf{w}}; \\ &\quad \langle_n^{\sharp}, \sharp \in \{\mathsf{k}\mathsf{b}, \mathsf{w}, \mathsf{p}, \mathsf{l}, \mathsf{p}\mathsf{l}\}, \\ &\quad \mathsf{0}^*(\ldots), \mathsf{1}^*(\ldots), * \in \{\mathsf{w}, \mathsf{p}\} \quad \rangle. \end{split}$$

- $\mathfrak{A}_{\mathsf{TA}}$: Term algebras.
- $\mathfrak{A}_{\mathbb{Z}}$: Presburger arithmetic.
- (.)^w : Weight function.





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$$\begin{split} \mathfrak{A}_{\mathsf{k}\mathsf{b}^+}^{\mathbb{Z}} &= \langle \quad \mathfrak{A}_{\mathsf{T}\mathsf{A}}; \mathfrak{A}_{\mathbb{Z}}; (.)^{\mathsf{w}}; \\ &\quad \langle_n^{\sharp}, \sharp \in \{\mathsf{k}\mathsf{b}, \mathsf{w}, \mathsf{p}, \mathsf{l}, \mathsf{p}\mathsf{l}\}, \\ &\quad \mathsf{0}^*(\ldots), \mathsf{1}^*(\ldots), * \in \{\mathsf{w}, \mathsf{p}\} \quad \rangle. \end{split}$$

- $\mathfrak{A}_{\mathsf{TA}}$: Term algebras.
- $\mathfrak{A}_{\mathbb{Z}}$: Presburger arithmetic.
- (.)^w : Weight function.
- $<_n^{\sharp}$: Gap orders.
- 0^{*}(...), 1^{*}(...) : Boundary terms.



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$$\begin{split} \mathfrak{A}_{\mathsf{k}\mathsf{b}^+}^{\mathbb{Z}} &= \langle \quad \mathfrak{A}_{\mathsf{T}\mathsf{A}}; \mathfrak{A}_{\mathbb{Z}}; (.)^{\mathsf{w}}; \\ &\quad \langle_n^{\sharp}, \sharp \in \{\mathsf{k}\mathsf{b}, \mathsf{w}, \mathsf{p}, \mathsf{l}, \mathsf{p}\mathsf{l}\}, \\ &\quad \mathsf{0}^*(\ldots), \mathsf{1}^*(\ldots), * \in \{\mathsf{w}, \mathsf{p}\} \quad \rangle. \end{split}$$

- $\mathfrak{A}_{\mathsf{TA}}$: Term algebras.
- $\mathfrak{A}_{\mathbb{Z}}$: Presburger arithmetic.
- (.)^w : Weight function.
- $<_n^{\sharp}$: Gap orders.
- 0^{*}(...), 1^{*}(...) : Boundary terms.
- Example of Use:

$$(\exists x: \mathsf{TA}) \left[\mathsf{0}_{(x^{\mathsf{w}})}^{\mathsf{w}} \prec_{2}^{\mathsf{l}} x \prec_{3}^{\mathsf{l}} \mathsf{1}_{(x^{\mathsf{w}})}^{\mathsf{w}} \right]$$



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Input: $(\exists \bar{x}) \varphi(\bar{x}, \bar{y}).$

<u>While</u> $\bar{x} \neq \emptyset$.

• While $(\forall x \in \bar{x}) depth_{\varphi}(x) > 0.$

Depth Reduction.

- ♦ VARIABLE SELECTION.
- DECOMPOSITION.
- SIMPLIFICATION.

Done.

- While $(\exists x \in \bar{x}) depth_{\varphi}(x) = 0.$
 - Elimination.

Done.

Done.



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Select a variable $x \in \bar{x}$ such that $s_i^{\alpha}(x)$ appears in $\varphi(\bar{x}, \bar{y})$.



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Select a variable $x \in \bar{x}$ such that $s_i^{\alpha}(x)$ appears in $\varphi(\bar{x}, \bar{y})$.

The variable selection is done in depth-fi rstmanner.



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Select a variable $x \in \bar{x}$ such that $s_i^{\alpha}(x)$ appears in $\varphi(\bar{x}, \bar{y})$.

The variable selection is done in depth-fi rstmanner.

I.e., choose variables generated in the previous round.



Decomposition

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Rewrite $(\exists \bar{x}) \varphi(\bar{x}, \bar{y})$ to

$$\exists x_1 \dots \exists x_k \exists \bar{x} \left[\mathsf{Is}_{\alpha}(x) \land \bigwedge_{1 \leq i \leq k} \mathsf{s}_i^{\alpha}(x) = x_i \land \varphi(\bar{x}, \bar{y}) \right].$$



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• Apply the following rules to each occurrence of x.



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- Apply the following rules to each occurrence of x.
 - 1. Replace $x <_{n}^{\sharp} t$ (or $t <_{n}^{\sharp} x$) by a quantifier-free formula $\varphi'(\mathbf{s}_{1}^{\alpha}(x), \dots, \mathbf{s}_{k}^{\alpha}(x), \mathbf{s}_{1}^{\alpha}(t), \dots, \mathbf{s}_{k}^{\alpha}(t)).$



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Apply the following rules to each occurrence of *x*.

1. Replace $x <_{n}^{\sharp} t$ (or $t <_{n}^{\sharp} x$) by a quantifier-free formula $\varphi'(\mathbf{s}_{1}^{\alpha}(x), \dots, \mathbf{s}_{k}^{\alpha}(x), \mathbf{s}_{1}^{\alpha}(t), \dots, \mathbf{s}_{k}^{\alpha}(t)).$

2. Replace $\mathbf{s}_i^{\alpha}(x)$ in $\varphi(\bar{\mathbf{x}}, \bar{\mathbf{y}})$ by x_i $(1 \le i \le k)$.



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- Apply the following rules to each occurrence of x.
 - 1. Replace $x <_{n}^{\sharp} t$ (or $t <_{n}^{\sharp} x$) by a quantifier-free formula $\varphi'(s_{1}^{\alpha}(x), \dots, s_{k}^{\alpha}(x), s_{1}^{\alpha}(t), \dots, s_{k}^{\alpha}(t)).$
 - 2. Replace $s_i^{\alpha}(x)$ in $\varphi(\bar{x}, \bar{y})$ by x_i $(1 \le i \le k)$.
- Denote the result of this simplifi cation as
 - $\exists x_1 \ldots \exists x_k \exists (\bar{x} \setminus x) \left[\varphi'(\bar{x} \setminus x, x_1, \ldots, x_k, \bar{y}) \right].$



Elimination

Now we must have

 $\exists x \Big[\bigwedge u_i <^{\mathsf{kb}} x \land \bigwedge x <^{\mathsf{kb}} v_j \land \varphi'(\bar{y}) \Big],$

 $i \le m$ $j \le n$

where x appears none of u_i , v_j and φ' .

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Now we must have

 $\exists x \Big[\bigwedge u_i <^{\mathsf{kb}} x \land \bigwedge x <^{\mathsf{kb}} v_j \land \varphi'(\bar{y}) \Big],$ i≤n i<m

where x appears none of u_i , v_j and φ' .

Guess a gap order completion, we rewrite it to

 $u_{i'} \prec_2^{\mathsf{kb}} v_{j'} \land \varphi'(\bar{y})$

 \wedge " $u_{i'}$ is the greatest of $\{u_i | i \leq m\}$ "

 \wedge " $v_{j'}$ is the smallest of $\{v_j \mid j \leq n\}$ ".



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Termination is subtle as some complexity measures increase.
■ Depth reduction increases the depth of other variables.
For example, *x* ≠ *t* becomes

$$\bigvee_{1\leq i\leq k} \mathbf{S}_i^{\alpha}(t) \neq x_i \vee \neg \mathbf{IS}_{\alpha}(t).$$



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Termination is subtle as some complexity measures increase.

Depth reduction increases the depth of other variables.
For example, $x \neq t$ becomes

$$\bigvee_{1 \le i \le k} \mathbf{S}_i^{\alpha}(t) \neq x_i \lor \neg \mathbf{IS}_{\alpha}(t).$$

Depth reduction introduces more existential quantifiers.
For example, $(\exists \bar{x}) \varphi(\bar{x}, \bar{y})$ becomes

$$\exists x_1 \dots \exists x_k \exists \bar{x} \left[\mathsf{Is}_{\alpha}(x) \land \bigwedge_{1 \leq i \leq k} \mathsf{s}_i^{\alpha}(x) = x_i \land \varphi(\bar{x}, \bar{y}) \right].$$



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Depth reduction introduces more order literals.
 For example, u <^w₅ v becomes

$$\bigvee_{n_1+n_2+n_3=5} u \prec_{n_1}^{\mathsf{pl}} \mathbf{1}_{(u^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_2}^{\mathsf{w}} \mathbf{0}_{(v^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_3}^{\mathsf{pl}} v.$$



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Depth reduction introduces more order literals. For example, $u <_{5}^{w} v$ becomes

$$\bigvee_{n_1+n_2+n_3=5} \mathcal{U} \prec_{n_1}^{\mathsf{pl}} \mathbf{1}_{(u^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_2}^{\mathsf{w}} \mathbf{0}_{(v^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_3}^{\mathsf{pl}} v.$$

Depth reduction introduces more equalities.
For example, $x <^{l} t$ could produce

 $\operatorname{car}(x) = \operatorname{car}(t) \wedge \operatorname{cdr}(x) <^{\mathsf{kb}} \operatorname{cdr}(t).$



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Depth reduction introduces more order literals. For example, $u \prec_5^w v$ becomes

$$\bigvee_{n_1+n_2+n_3=5} \mathcal{U} \prec_{n_1}^{\mathsf{pl}} \mathbf{1}_{(\mathcal{U}^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_2}^{\mathsf{w}} \mathbf{0}_{(\mathcal{v}^{\mathsf{w}})}^{\mathsf{w}} \prec_{n_3}^{\mathsf{pl}} \mathcal{v}.$$

Depth reduction introduces more equalities.
For example, x <^I t could produce

 $\operatorname{car}(x) = \operatorname{car}(t) \wedge \operatorname{cdr}(x) <^{\mathsf{kb}} \operatorname{cdr}(t).$

Does it terminate???



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Real measure: # of open gap order literals (OGOL).

OGOL: a gap order relation between ordinary terms.



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Real measure: # of open gap order literals (OGOL).

OGOL: a gap order relation between ordinary terms.

No transformation generates new OGOLs.



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Real measure: # of open gap order literals (OGOL).

OGOL: a gap order relation between ordinary terms.

No transformation generates new OGOLs.

The final elimination step removes at least one OGOL.



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Real measure: # of open gap order literals (OGOL).

OGOL: a gap order relation between ordinary terms.

No transformation generates new OGOLs.

The final elimination step removes at least one OGOL.

Without OGOLs, the depths of terms strictly decrease!



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Consider the KBO on LISP list structure parameterized with

 $W(cons) = W(nil) = 1 \& nil <^{\Sigma} cons.$



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• Consider the KBO on LISP list structure parameterized with W(cons) = W(nil) = 1 & nil $<^{\Sigma}$ cons.

Consider the formula

 $(\exists x) \left[\operatorname{car}(x) \prec_2^{\mathsf{l}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{l}} y \right],$ where depth(x) = 3.



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 $(\exists x) \left[\operatorname{car}(x) \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{I}} y \right],$



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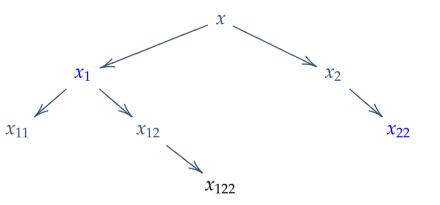
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 $(\exists x) \left[\operatorname{car}(x) \prec_2^{\mathsf{l}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{l}} y \right],$





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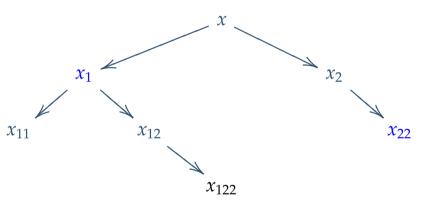
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 $(\exists x) \left[\operatorname{car}(x) \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{I}} y \right],$



 $x_1 : car(x),$ $x_2 : cdr(x),$ $x_{11} : car(car(x)),$ $x_{12} : cdr(car(x)),$ $x_{22} : cdr(cdr(x)),$ $x_{122} : cdr(cdr(car(x)))$



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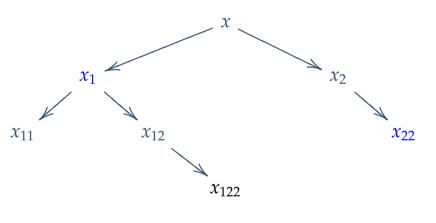
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 $(\exists x) \left[\operatorname{car}(x) \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{I}} y \right],$



Solution: x = ?

 $x_1 : car(x),$ $x_2 : cdr(x),$ $x_{11} : car(car(x)),$ $x_{12} : cdr(car(x)),$ $x_{22} : cdr(cdr(x)),$ $x_{122} : cdr(cdr(car(x)))$



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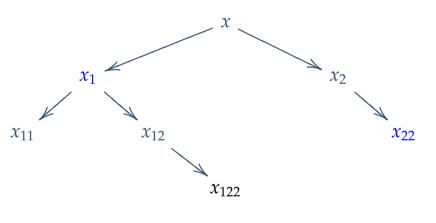
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 $(\exists x) \left[\operatorname{car}(x) \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{I}} y \right],$



Solution: x = ?

 $x_1 : car(x),$ $x_2 : cdr(x),$ $x_{11} : car(car(x)),$ $x_{12} : cdr(car(x)),$ $x_{22} : cdr(cdr(x)),$ $x_{122} : cdr(cdr(car(x)))$



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Select variable x.



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Select variable x.

• Decompose x in terms of car(x) and cdr(x). We have

$$(\exists x \exists x_1 \exists x_2) \left[\operatorname{car}(x) = x_1 \land \operatorname{cdr}(x) = x_2 \\ \land \operatorname{car}(x) \prec_2^{\mathsf{l}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{l}} y \right].$$



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Select variable x.

• Decompose x in terms of car(x) and cdr(x). We have

$$\exists x \exists x_1 \exists x_2) \Big[\operatorname{car}(x) = x_1 \wedge \operatorname{cdr}(x) = x_2 \\ \wedge \operatorname{car}(x) \prec_2^{\mathsf{l}} \operatorname{cdr}(\operatorname{cdr}(x)) \wedge \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec_3^{\mathsf{l}} y \Big].$$

Simplifi cation.

 $(\exists x_1 \exists x_2) \Big[x_1 \prec_2^{\mathsf{I}} \mathsf{cdr}(x_2) \land \mathsf{cdr}(\mathsf{cdr}(x_1)) \prec_3^{\mathsf{I}} y \Big],$ where $depth(x_1) = 2$ and $depth(x_2) = 1,$



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 $(\exists x_1 \exists x_2) \left[x_1 \prec_2^{\mathsf{I}} \mathsf{cdr}(x_2) \land \mathsf{cdr}(\mathsf{cdr}(x_1)) \prec_3^{\mathsf{I}} y \right].$



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 $(\exists x_1 \exists x_2) \left[x_1 \prec_2^{\mathsf{I}} \mathsf{cdr}(x_2) \land \mathsf{cdr}(\mathsf{cdr}(x_1)) \prec_3^{\mathsf{I}} y \right].$

• Select variable x_1 .



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```
(\exists x_1 \exists x_2) \Big[ x_1 \prec_2^{\mathsf{I}} \mathsf{cdr}(x_2) \land \mathsf{cdr}(\mathsf{cdr}(x_1)) \prec_3^{\mathsf{I}} y \Big].
```

• Select variable x_1 .

• Decompose x_1 . Replace $x_1 \prec_2^{\mathsf{I}} \mathsf{cdr}(x_2)$ by

 $\operatorname{car}(x_1) = \operatorname{car}(\operatorname{cdr}(x_2)) \wedge \operatorname{cdr}(x_1) \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x_2)).$



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 $(\exists x_1 \exists x_2) \Big[x_1 \prec_2^{\mathsf{I}} \mathsf{cdr}(x_2) \land \mathsf{cdr}(\mathsf{cdr}(x_1)) \prec_3^{\mathsf{I}} y \Big].$

■ Select variable *x*₁.

• Decompose x_1 . Replace $x_1 \prec_2^{\mathsf{I}} \mathsf{cdr}(x_2)$ by

 $\operatorname{car}(x_1) = \operatorname{car}(\operatorname{cdr}(x_2)) \land \operatorname{cdr}(x_1) \prec_2^{\mathsf{l}} \operatorname{cdr}(\operatorname{cdr}(x_2)).$

Simplifi cation.

$$(\exists x_2 \exists x_{11} \exists x_{12}) \left[x_{11} = \operatorname{car}(\operatorname{cdr}(x_2)) \land x_{12} \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x_2)) \land \operatorname{cdr}(x_{12}) \prec_3^{\mathsf{I}} y \right],$$



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 $(\exists x_2 \exists x_{11} \exists x_{12}) \left[\begin{array}{c} x_{11} = \operatorname{car}(\operatorname{cdr}(x_2)) \land x_{12} \prec_2^{\mathsf{l}} \operatorname{cdr}(\operatorname{cdr}(x_2)) \\ \land \operatorname{cdr}(x_{12}) \prec_3^{\mathsf{l}} y \end{array} \right],$



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$(\exists x_2 \exists x_{11} \exists x_{12}) \left[\begin{array}{c} x_{11} = \mathsf{car}(\mathsf{cdr}(x_2)) \land x_{12} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \\ \land \ \mathsf{cdr}(x_{12}) \prec_3^{\mathsf{I}} y \end{array} \right],$

• Elimination. Since depth $(x_{11}) = 0$, we have

$$(\exists x_2 \exists x_{12}) \left[x_{12} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land \mathsf{cdr}(x_{12}) \prec_3^{\mathsf{I}} y \right],$$



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 $(\exists x_2 \exists x_{12}) \left[x_{12} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land \mathsf{cdr}(x_{12}) \prec_3^{\mathsf{I}} y \right].$



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 $(\exists x_2 \exists x_{12}) \left[x_{12} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land \mathsf{cdr}(x_{12}) \prec_3^{\mathsf{I}} y \right].$

• Select variable x_{12} .



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 $(\exists x_2 \exists x_{12}) \begin{bmatrix} x_{12} < \\ 2 \end{bmatrix} cdr(cdr(x_2))$

$$_{2})\left[x_{12} \prec_{2}^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_{2})) \land \mathsf{cdr}(x_{12}) \prec_{3}^{\mathsf{I}} y \right].$$

• Select variable x_{12} .

• Decompose x_{12} . Replace $x_{12} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2))$ by

 $x_{121} = car(cdr(cdr(x_2))) \land x_{122} \prec_2^{\mathsf{I}} cdr(cdr(x_2)).$



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 $(\exists x_2 \exists x_{12}) \left[x_{12} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land \mathsf{cdr}(x_{12}) \prec_3^{\mathsf{I}} y \right].$

• Select variable x_{12} .

• Decompose x_{12} . Replace $x_{12} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2))$ by

 $x_{121} = car(cdr(cdr(x_2))) \land x_{122} \prec_2^{l} cdr(cdr(x_2)).$

Simplifi cation.

 $(\exists x_2 \exists x_{121} \exists x_{122}) \left[x_{121} = \operatorname{car}(\operatorname{cdr}(\operatorname{cdr}(x_2))) \land x_{122} \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right],$



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$(\exists x_2 \exists x_{121} \exists x_{122}) \left[x_{121} = \operatorname{car}(\operatorname{cdr}(\operatorname{cdr}(x_2))) \land x_{122} \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right],$



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$(\exists x_2 \exists x_{121} \exists x_{122}) \left[x_{121} = \operatorname{car}(\operatorname{cdr}(\operatorname{cdr}(x_2))) \land x_{122} \prec_2^{\mathsf{I}} \operatorname{cdr}(\operatorname{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right],$

Elimination. Since depth $(x_{121}) = 0$, we have

 $(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right].$



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 $(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right].$



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 $(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right].$

Elimination.



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 $(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right].$

Elimination.

Take a gap order completion

$$x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y$$



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 $(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right].$

Elimination.

Continue with

Take a gap order completion

$$x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y$$

We have

 $(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y \right],$



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$$(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right].$$

Elimination.

Take a gap order completion

$$x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y$$

We have

$$(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y \right],$$

which simplifies to

$$(\exists x_2) \left[\mathsf{O}^{\mathsf{w}}_{((\mathsf{cdr}(\mathsf{cdr}(x_2)))^{\mathsf{w}})} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y \right]$$



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$$(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \land x_{122} \prec_3^{\mathsf{I}} y \right].$$

Elimination.

Take a gap order completion

$$x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y$$

We have

$$(\exists x_2 \exists x_{122}) \left[x_{122} \prec_2^{\mathsf{I}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{I}} y \right],$$

which simplifies to

$$(\exists x_2) \left[\mathsf{O}^{\mathsf{w}}_{((\mathsf{cdr}(\mathsf{cdr}(x_2)))^{\mathsf{w}})} \prec_2^{\mathsf{l}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{l}} y \right]$$

The number of OGOLs reduced to 1!



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 $(\exists x_2) \left[0^{\mathsf{w}}_{((\mathsf{cdr}(\mathsf{cdr}(x_2)))^{\mathsf{w}})} \prec_2^{\mathsf{l}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{l}} y \right].$



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$(\exists x_2) \left[0^{\mathsf{w}}_{((\mathsf{cdr}(\mathsf{cdr}(x_2)))^{\mathsf{w}})} \prec_2^{\mathsf{l}} \mathsf{cdr}(\mathsf{cdr}(x_2)) \prec_1^{\mathsf{l}} y \right].$ Depth Reduction. Repeating twice the DEPTH-REDUCTION

subprocedure, we have

$$(\exists x_{222}) \left[0^{\mathsf{w}}_{(x^{\mathsf{w}}_{222})} \prec^{\mathsf{l}}_{2} x_{222} \prec^{\mathsf{l}}_{1} y \right].$$



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 $(\exists x_{222}) \left[0^{\mathsf{w}}_{(x^{\mathsf{w}}_{222})} \prec^{\mathsf{l}}_{2} x_{222} \prec^{\mathsf{l}}_{1} y \right].$



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 $(\exists x_{222}) \left[0^{\mathsf{w}}_{(x^{\mathsf{w}}_{222})} \prec_{2}^{\mathsf{l}} x_{222} \prec_{1}^{\mathsf{l}} y \right].$

Reduce term quantifi ers to integer quantifi ers.

$$(\exists z) \left[\mathsf{0}_{(z)}^{\mathsf{w}} \prec_{3}^{\mathsf{l}} y \land \mathsf{Tree}^{\mathsf{cons}}(z) \right].$$



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 $(\exists x_{222}) \left[\begin{array}{c} \mathbf{0}_{(x_{222}^{\mathsf{w}})}^{\mathsf{w}} \prec_{2}^{\mathsf{l}} x_{222} \prec_{1}^{\mathsf{l}} y \end{array} \right].$

Reduce term quantifi ers to integer quantifi ers.

$$(\exists z) \left[0_{(z)}^{\mathsf{w}} \prec_{3}^{\mathsf{l}} y \wedge \mathsf{Tree}^{\mathsf{cons}}(z) \right].$$

Eliminate integer quantifi ers.

 $0^{\mathsf{w}}_{(y^{\mathsf{w}})} \prec^{\mathsf{l}}_{3} y \wedge \operatorname{Tree}^{\operatorname{cons}}(y^{\mathsf{w}}).$



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$$(\exists x_{222}) \left[0^{\mathsf{w}}_{(x^{\mathsf{w}}_{222})} \prec^{\mathsf{l}}_{2} x_{222} \prec^{\mathsf{l}}_{1} y \right].$$

■ Reduce term quantifiers to integer quantifiers. $(\exists z) \left[0_{(z)}^{w} \prec_{3}^{l} y \land \text{Tree}^{\text{cons}}(z) \right].$

Eliminate integer quantifi ers.

$$0^{\mathsf{w}}_{(y^{\mathsf{w}})} \prec^{\mathsf{l}}_{3} y \wedge \mathsf{Tree}^{\mathsf{cons}}(y^{\mathsf{w}}).$$

• As
$$0_{(y^w)}^w \prec_3^l y$$
 implies Tree^{cons} (y^w) , we have

$$\mathsf{O}^{\mathsf{w}}_{(y^{\mathsf{w}})} \prec^{\mathsf{l}}_{3} y.$$



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 $\begin{array}{l} \mathbf{0}_{(y^{w})}^{w} <_{3}^{\mathsf{l}} y \implies \\ (\exists x) \left[\operatorname{car}(x) <_{2}^{\mathsf{l}} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) <_{3}^{\mathsf{l}} y \right], \end{array}$



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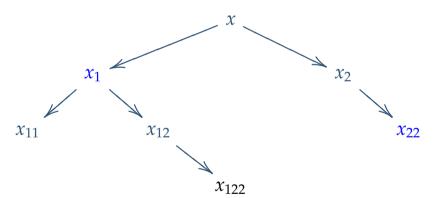
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 $0^{\mathsf{w}}_{(y^{\mathsf{w}})} \prec^{\mathsf{l}}_{3} y \implies (\exists x) \left[\operatorname{car}(x) \prec^{\mathsf{l}}_{2} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec^{\mathsf{l}}_{3} y \right],$





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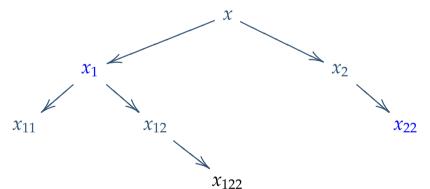
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 $0^{\mathsf{w}}_{(y^{\mathsf{w}})} \prec^{\mathsf{l}}_{3} y \implies (\exists x) \left[\operatorname{car}(x) \prec^{\mathsf{l}}_{2} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec^{\mathsf{l}}_{3} y \right],$



 $x_1 : car(x),$ $x_2 : cdr(x),$ $x_{11} : car(car(x)),$ $x_{12} : cdr(car(x)),$ $x_{22} : cdr(cdr(x)),$ $x_{122} : cdr(cdr(car(x)))$



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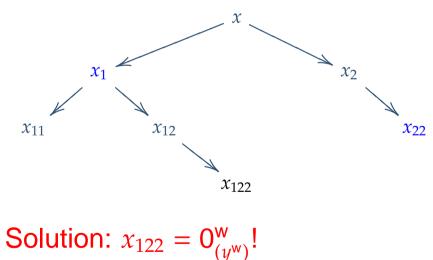
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 $0^{\mathsf{w}}_{(y^{\mathsf{w}})} \prec^{\mathsf{l}}_{3} y \implies (\exists x) \left[\operatorname{car}(x) \prec^{\mathsf{l}}_{2} \operatorname{cdr}(\operatorname{cdr}(x)) \land \operatorname{cdr}(\operatorname{cdr}(\operatorname{car}(x))) \prec^{\mathsf{l}}_{3} y \right],$



 $x_1 : car(x),$ $x_2 : cdr(x),$ $x_{11} : car(car(x)),$ $x_{12} : cdr(car(x)),$ $x_{22} : cdr(cdr(x)),$ $x_{122} : cdr(cdr(x)))$



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Smallest extensions for quantifi er elimination.

More expressive power induces higher complexity.



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Smallest extensions for quantifi er elimination.

More expressive power induces higher complexity.

Block-wise quantifi er elimination.

Small quantifi er alternations in real life.



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Smallest extensions for quantifi er elimination.

More expressive power induces higher complexity.

Block-wise quantifi er elimination.

Small quantifi er alternations in real life.

Decidability of KBO on term domain with variables.



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- Simplification of Selector Terms
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Elimination of Equalities.

 $\exists x \left[x = \mathbf{0}_{((\operatorname{car}(x))^{w})}^{\mathsf{w}} \land \operatorname{car}(x) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(x) \right].$



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Elimination of Equalities.

$$\exists x \left[x = \mathbf{0}_{((\operatorname{car}(x))^{w})}^{\mathsf{w}} \land \operatorname{car}(x) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(x) \right].$$

Simplifi cation of Selector Terms.

 $\operatorname{car}(0^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}).$



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Elimination of Equalities.

$$\exists x \left[x = \mathbf{0}_{((\operatorname{car}(x))^{w})}^{\mathsf{w}} \land \operatorname{car}(x) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(x) \right].$$

Simplifi cation of Selector Terms.

 $\operatorname{car}(0^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}).$

Elimination of Negations.

$$\neg (\operatorname{car}(x) \prec^{\mathsf{w}}_{3} \operatorname{cdr}(x)).$$



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Continue with

$$\exists x \left[x = \mathbf{0}_{((\operatorname{car}(x))^{w})}^{w} \land \operatorname{car}(x) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(x) \right].$$



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Continue with

$$\exists x \left[x = \mathbf{0}_{((\operatorname{car}(x))^{\mathsf{w}})}^{\mathsf{w}} \land \operatorname{car}(x) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(x) \right].$$

Substitution.

$$\exists x \left[x = \mathsf{O}^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})} \land \operatorname{car}(\mathsf{O}^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(\mathsf{O}^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}) \right].$$



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$$\exists x \left[x = \mathbf{0}_{((\operatorname{car}(x))^{w})}^{w} \land \operatorname{car}(x) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(x) \right].$$

Substitution.

$$\exists x \left[x = \mathsf{O}^{\mathsf{w}}_{((\mathsf{car}(x))^{\mathsf{w}})} \land \mathsf{car}(\mathsf{O}^{\mathsf{w}}_{((\mathsf{car}(x))^{\mathsf{w}})}) \prec_{4}^{\mathsf{p}} \mathsf{cdr}(\mathsf{O}^{\mathsf{w}}_{((\mathsf{car}(x))^{\mathsf{w}})}) \right].$$

Reduction to Integer Quantifi ers.

 $\exists (\operatorname{car}(x))^{\mathsf{w}} \Big[\operatorname{Tree}((\operatorname{car}(x))^{\mathsf{w}}) \land \operatorname{car}(\mathsf{O}^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(\mathsf{O}^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}) \Big].$



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Continue with

$$\exists x \left[x = \mathbf{0}_{((\operatorname{car}(x))^{\mathsf{w}})}^{\mathsf{w}} \land \operatorname{car}(x) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(x) \right].$$

Substitution.

$$\exists x \left[x = \mathsf{O}^{\mathsf{w}}_{((\mathsf{car}(x))^{\mathsf{w}})} \land \mathsf{car}(\mathsf{O}^{\mathsf{w}}_{((\mathsf{car}(x))^{\mathsf{w}})}) \prec_{4}^{\mathsf{p}} \mathsf{cdr}(\mathsf{O}^{\mathsf{w}}_{((\mathsf{car}(x))^{\mathsf{w}})}) \right].$$

Reduction to Integer Quantifiers.

 $\exists (\operatorname{car}(x))^{\mathsf{w}} \Big[\operatorname{Tree}((\operatorname{car}(x))^{\mathsf{w}}) \land \operatorname{car}(\mathsf{0}^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}) \prec_{4}^{\mathsf{p}} \operatorname{cdr}(\mathsf{0}^{\mathsf{w}}_{((\operatorname{car}(x))^{\mathsf{w}})}) \Big].$

Renaming.

 $\exists z \left[\operatorname{Tree}(z) \land \operatorname{car}(0^{\mathsf{w}}_{(z)}) \prec^{\mathsf{p}}_{4} \operatorname{cdr}(0^{\mathsf{w}}_{(z)}) \right].$



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 $car(0^{w}_{((car(x))^{w})})$



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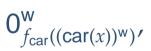
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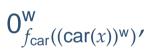
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 $car(0_{((car(x))^w)}^w)$



 $f_{car}(\cdot)$

is an integer function from Presburger arithmetic.



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 $\neg (\operatorname{car}(x) \prec_3^{\mathsf{w}} \operatorname{cdr}(x)).$



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$$\neg (\operatorname{car}(x) \prec_3^{\mathsf{w}} \operatorname{cdr}(x)).$$

 $\operatorname{cdr}(x) <^{\mathsf{w}} \operatorname{car}(x) \lor (\operatorname{cdr}(x))^{\mathsf{w}} = (\operatorname{car}(x))^{\mathsf{w}} \lor$ $\operatorname{car}(x) \leq_{1}^{\mathsf{w}} \operatorname{cdr}(x) \lor \operatorname{car}(x) \leq_{2}^{\mathsf{w}} \operatorname{cdr}(x).$

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